The Newton-Puisuex Polygon for q—difference equations. Applications.

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First things first

1.-Thanks to the Scientific Committee

2.-Thanks to the Organizing Committee (Prof. Walser, Greg, Galina, Sławomir)

3.-Enjoy

Prior Art

- Maillet-Malgrange (1903-1989): Any power series solution of an analytic ODE is of Gevrey type.
- J. Cano (1993): Any power series solution of a Gevrey ODE is Gevrey, explicit computation (Newton Polygon).
- Zhang 1998 (& di Vizio 2007 ultrametric): Any power series solution of an analytic q-difference equation is of q-Gevrey type.

This work is the natural next step. Techniques like Cano's.

$$P(x, y, \sigma(y)) = \sigma(y)y^{2} + \sigma(y)^{2} + x^{2}y + x = 0$$

Given a general q-difference equation like (no independent term)

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Compute solutions

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- Behaviour of solutions
 - Convergence, *q*—Gevrey asymptotics
 - Rational rank, "size" of the exponent's semigroup

Framework

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$$P(x, y_0, \ldots, y_n) \in \mathbb{C}[[x^{\mathbb{R}_{>0}}]][[y_0, \ldots, y_n]]$$

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Solution: A solution of P is a generalized power series

$$f(x) = \sum_{\gamma \in \Gamma} f_{\gamma} x^{\gamma}, \text{ with } \Gamma \subset \mathbb{R}_{>0}$$

where Γ is well ordered, such that

$$P(x, f(x), f(qx), \dots, f(q^n x)) = 0$$

(a determination of the logarithm is fixed).

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$$y - x^{10} - x^{15} + [\dots] + 10x^9y^2 + [\dots] + y^5 = 0$$

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Go on. This one was easy.

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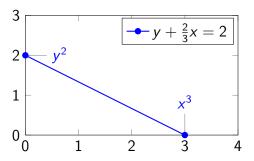
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What with ...

$$y^6 + x^5y^5 + xy^4 + x^3y^2 + x^{10} = 0?$$

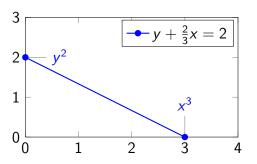
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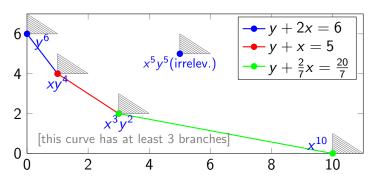


The coefficient $y = 1 \cdot x^{3/2}$ comes from substituting $y = y + cx^{3/2}$, which gives at y = 0:

$$x^3 + c^2 x^3 = 0 \Rightarrow c = \pm 1.$$

The Polygon: positive convex hull of cloud

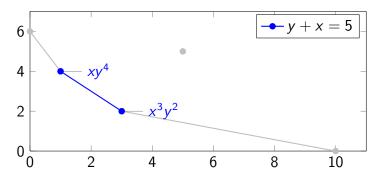
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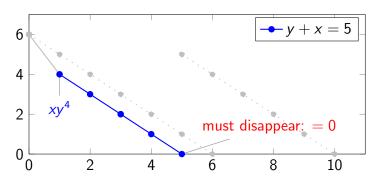
The inner point (5,5) is irrelevant for starting a solution. Solutions $y = cx^{\alpha} + \ldots$ admit $\alpha = \frac{1}{2}, 1, \frac{7}{2}$, the *inclinations*.

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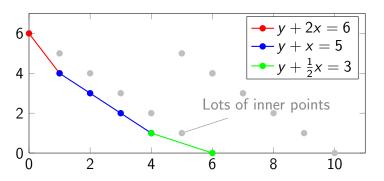
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Lowest x—term for that side must be 0:

$$P_{\alpha}(C) = x(y+cx)^4 + x^3(y+cx)^2\Big]_{y=0} = 0 \simeq c^4 + c^2 = 0.$$

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Setting c = i, a new Newton Polygon appears. The green side is *new*, and has inclination $\alpha = 1$.

Recurse with new equation and Polygon.

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But: in the algebraic case, μ is always rational and any slope and root give rise to a solution.

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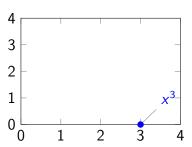
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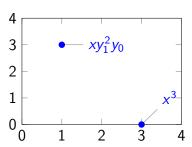
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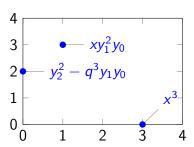
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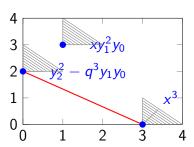


Newton-Puisuex for q-difference equations

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Any c will make $Q_\mu(c) = 0$: free slopes.

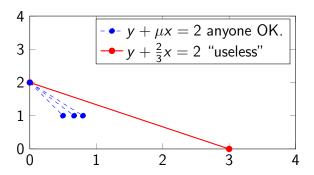
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Any c will make $Q_\mu(c) = 0$: free slopes.
Constant Q : For $\mu = \frac{2}{3} \to ln_\mu = y_2^2 - q^3 y_1 y_0 + x^3$, so $Q_\mu(T) = (q^2 T)^2 - q^3 (qT) T + 1 \equiv 1$
No c will make $Q_\mu(c) = 0$. Useless slope.

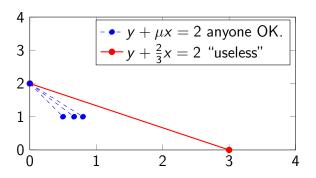
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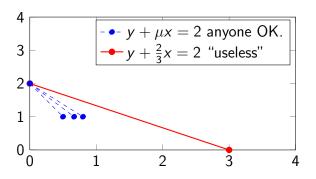
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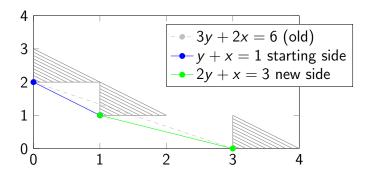
- $Q_{\mu}(T) = 0$ for $\mu > 2/3$: $y = cx^{1/\mu} + ...$ possible (may have $\mu \notin \mathbb{Q}$: not so in algebraic curves).
- $Q_{2/3}(T) = 1$: no solution starting with $cx^{3/2}$.

Substitution step

After substituting

$$y = y + cx^{\alpha}$$

new Polygon, continue with greater inclination (to the right). Say y = y + x in $y_2^2 - q^3y_1y_0 + xy_1^2y_0 + x^3 = 0$,



Results: Solution construction

Lemma

Generically, the Newton-Puiseux algorithm gives rise to a solution. And solutions correspond to the algorithm.

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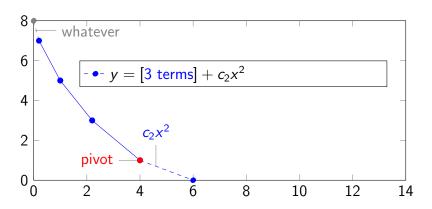
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Theorem

If $P = A(x, y_0) + B(x, y_0)y_1$ ("order and degree 1"), then rational rank(Γ) ≤ 2 . ("Only one irrational exponent").

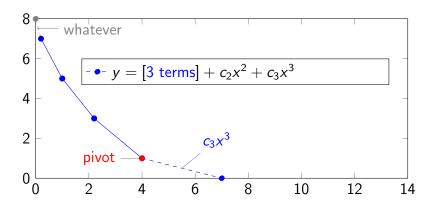
Pivot point

In a finite number of steps, the topmost vertex of the *interesting* side is fixed (usually at height 1).



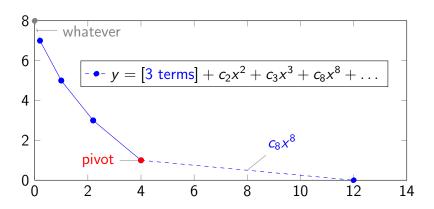
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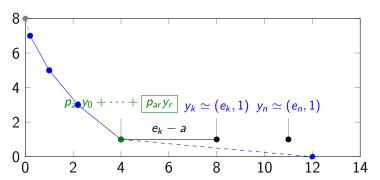
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Results: *q*—Gevrey bounds (intro)

$$t+1-$$
 Gevrey $\Leftrightarrow \sum \frac{A_n}{|q|^{tn(n+1)/2}} x^n$ convergent

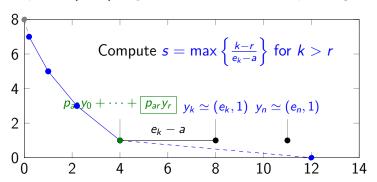
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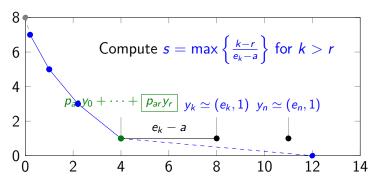
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For pivot at (a, b), analogous computation.

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Formal power series solution $f(x) \in \mathbb{C}[[x]]$ (can be generalized to rational exponents easily). Recall $s = \frac{\Delta \text{ord}}{\text{distance}}$.

Theorem (q-Gevrey Malgrange-Maillet)

If P has q-Gevrey order t+1, then f(x) has q-Gevrey order $\leq s+t+1$.

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Corollary

If the pivot point of f(x) contains order n, then f(x) has the same q—Gevrey order as P.

Results: closed-eyed bound

Corollary

If $P(x, y_0, ..., y_n)$ has q-Gevrey order t + 1, then any solution has at most q-Gevrey order n + t + 1.

No information required on the solution, whereas the results above assume the pivot point is known.

Thanks, etc

Questions?



And many thanks