

On the study of the operator norms of some random matrices.

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Random operators.

Local theory of Banach spaces : how to produce subspaces with **nice properties** in a Banach space ?

The Euclidean case

$g_j \sim \mathcal{N}(0, 1)$, $g = (g_j)_{j \leq k}$

Take $x \in (\mathbb{R}^k, |\cdot|_2)$ then $\sum x_j g_j \sim (\sum x_j^2)^{1/2} g_1$. That is

$$G(x) = \sum x_j g_j \in L_1$$

$\ell_2^k \xrightarrow{1} L_1$ isometrically.

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More generally, $g_{ij} \sim \mathcal{N}(0, 1)$

$$G = (g_{ij}) : \ell_2^k \rightarrow (\mathbb{R}^n, \|\cdot\|) = X$$

Find condition on k such that $\ell_2^k \xrightarrow{1+\varepsilon} X$

Dvoretzky-Milman in the '70's, Gordon in the '80's

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The non Euclidean case. $1 \leq p \leq 2$

θ_j p-stable random variables, $\theta = (\theta_j)_{j \leq k}$

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θ_j p-stable random variables, $\theta = (\theta_j)_{j \leq k}$. Take $x \in (\mathbb{R}^k, |\cdot|_p)$ then $\sum x_j \theta_j \sim (\sum |x_j|^p)^{1/p} \theta_1$. That is

$$\Theta(x) = \sum x_j \theta_j \in L_1$$

$\ell_p^k \xrightarrow{1} L_1$ isometrically.

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Hum, not so easy....

Not good concentration properties. You have to perturb the random operator,

Pisier '83,

Friedland-G '2013

Sparsity and compressed sensing

Reconstruction of a signal.

Let Φ be an $n \times N$ matrix. You receive $\Phi x := y$ where $x \in \mathbb{R}^N$ is the unknown signal.

How to recover x ? **Assumption** on the signal : it is an **m -sparse signal** with $m \ll n$.

Restricted Isometry Property with constant $\delta \in (0, 1)$: an $n \times N$ matrix Φ such that for all m -sparse vectors $x \in \mathbb{R}^N$,

$$(1 - \delta)|x|_2 \leq |\Phi x|_2 \leq (1 + \delta)|x|_2.$$

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Candès, Romberg, Tao [’06]

For such matrices, if δ is small enough, any m -sparse vectors is uniquely defined by

$$\min\{|t|_1 \text{ subject to } \Phi t = y\}$$

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Random matrices are good !

$\Phi = (g_{ij})/\sqrt{N}$ satisfies **RIP** with $m \approx c(\delta) \frac{n}{\log(eN/n)}$.

Question : give an **explicit construction** of such a matrix ?

Graph theory

Random graph

Classical Erdős-Rényi model : n vertices, edges between i and j are taken with probability p

Inhomogeneous Erdős-Rényi model : n vertices, edges between i and j with probability p_{ij}

Adjacency matrix

$$Adj = (x_{ij}) \quad \text{and} \quad \mathbb{E}Adj = (p_{ij})$$

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Want to evaluate $\mathbb{E}\|Adj - \mathbb{E}Adj\|_{2 \rightarrow 2}$? You start with a classical trick

$$\begin{aligned} \mathbb{E}\|(x_{ij} - \mathbb{E}x_{ij})\| &= \mathbb{E}\|(x_{ij} - \mathbb{E}'x'_{ij})\| \leq \mathbb{E}\mathbb{E}'\|(x_{ij} - x'_{ij})\| \\ &= \mathbb{E}\mathbb{E}'\mathbb{E}_\varepsilon\|(\varepsilon_{ij}(x_{ij} - x'_{ij}))\| \leq 2\mathbb{E}\mathbb{E}_\varepsilon\|(\varepsilon_{ij}x_{ij})\| \end{aligned}$$

Operator norms of random matrix

Take a deterministic matrix

$$X = (x_{ij})$$

Let R_i be its rows and C_j be its columns. Then

$$\|X\|_{2 \rightarrow 2} \geq \max \left(\max_i |R_i|_2, \max_j |C_j|_2 \right).$$

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Theorem [Seginer 2000]

If x_{ij} are independent **identically** distributed random variables then

$$\mathbb{E} \|X\|_{2 \rightarrow 2} \approx \mathbb{E} \max_i |R_i|_2 + \mathbb{E} \max_j |C_j|_2.$$

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More is true : for any $p \leq \log n$,

$$\mathbb{E} \|X\|_{2 \rightarrow 2}^p \approx \mathbb{E} \max_i |R_i|_2^p + \mathbb{E} \max_j |C_j|_2^p.$$

Operator norms of random matrix

The moment method

$$\|X\|_{2 \rightarrow 2} = \max_{1 \leq i \leq n} |s_i(X)| \leq \left(\sum_{i=1}^n |s_i(X)|^p \right)^{1/p} \leq n^{1/p} \max_{1 \leq i \leq n} |s_i(X)|$$

Hence for $p \approx \log n$,

$$\|X\|_{2 \rightarrow 2} \approx (\operatorname{tr}(X^*X)^{p/2})^{1/p}$$

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We have proven that for $p \approx \log n$,

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And you use combinatorics.

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And you use combinatorics.

Not enough. You need to use the assumption **identically** distributed. Hence you add an additional randomness by allowing random permutation on each coordinates.

Operator norms of random matrix

A simple case of independent **not identically** distributed r.v. :

Theorem [Seginer 2000]

If $x_{ij} = \varepsilon_{ij}a_{ij}$ where ε_{ij} are independent random variables then

$$\mathbb{E}\|X\|_{2 \rightarrow 2} \leq (\log n)^{1/4} \left(\max_i \sqrt{\sum_j a_{ij}^2} + \max_j \sqrt{\sum_i a_{ij}^2} \right).$$

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Conjecture

$$\mathbb{E}\|X\|_{2 \rightarrow 2} \approx \mathbb{E} \max_i |R_i|_2 + \mathbb{E} \max_j |C_j|_2.$$

is **not** valid in general.

The Gaussian case

We have

$$\mathbb{E}\|(\varepsilon_{ij}a_{ij})\|_{2 \rightarrow 2} \leq \sqrt{\frac{2}{\pi}} \mathbb{E}\|(g_{ij}a_{ij})\|_{2 \rightarrow 2}$$

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Theorem [Latała 2005]

$$\mathbb{E}\|(g_{ij}a_{ij})\|_{2 \rightarrow 2} \lesssim \max_i \sqrt{\sum_j a_{ij}^2} + \max_j \sqrt{\sum_i a_{ij}^2} + \left(\sum_{ij} a_{ij}^4\right)^{1/4}$$

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Theorem [Bandeira - Van Handel 2016]

$$\mathbb{E}\|(g_{ij}a_{ij})\|_{2 \rightarrow 2} \leq \max_i \sqrt{\sum_j a_{ij}^2} + \max_j \sqrt{\sum_i a_{ij}^2} + C \sqrt{\log n} \max |a_{ij}|$$

where C is a universal constant

The Gaussian case

Conjecture. In the Gaussian case, $X = (g_{ij}a_{ij})$, we have

$$\mathbb{E}\|X\|_{2 \rightarrow 2} \approx \mathbb{E} \max_i |R_i|_2 + \mathbb{E} \max_j |C_j|_2.$$

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Lemma [Van Handel 2017]

$$\begin{aligned} & \mathbb{E} \max_i \sqrt{\sum_j a_{ij}^2 g_{ij}^2} + \mathbb{E} \max_j \sqrt{\sum_i a_{ij}^2 g_{ij}^2} \\ & \leq \max_i \sqrt{\sum_j a_{ij}^2} + \max_j \sqrt{\sum_i a_{ij}^2} + C \mathbb{E} \max |a_{ij} g_{ij}| \end{aligned}$$

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Lemma [Van Handel 2017]

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This would fit with Seginer's example.

The Gaussian case

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$$\mathbb{E} \| (g_{ij} a_{ij}) \|_{2 \rightarrow 2} \leq \max_i \sqrt{\sum_j a_{ij}^2} + \max_j \sqrt{\sum_i a_{ij}^2} + C \sqrt{\log n} \max |a_{ij}|$$

where C is a universal constant

They use the moment method.

The Gaussian case

Maybe not the good strategy.

$$\mathbb{E} \|X\|_{2 \rightarrow 2} \leq \mathbb{E} (\operatorname{tr}(X^* X)^{p/2})^{1/p} \leq (\mathbb{E} \operatorname{tr}(X^* X)^{p/2})^{1/p}$$

and you take $p \approx \log n$.

The Gaussian case

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$$\mathbb{E} \|X\|_{2 \rightarrow 2} \leq \mathbb{E} (\operatorname{tr}(X^* X)^{p/2})^{1/p} \leq (\mathbb{E} \operatorname{tr}(X^* X)^{p/2})^{1/p}$$

and you take $p \approx \log n$. Example :

$$A = (a_{ij}) = \operatorname{diag}(a_1, \dots, a_n)$$

Hence

$$(\mathbb{E} \operatorname{tr}(X^* X)^{p/2})^{1/p} = \left(\mathbb{E} \sum_{i=1}^n |a_i|^p |g_i|^p \right)^{1/p} \approx \sqrt{p} \left(\sum_{i=1}^n |a_i|^p \right)^{1/p}$$

And then, you take $p \approx \log n$, which gives $\sqrt{\log n} \max |a_{ij}|$

While obviously

$$\mathbb{E} \|\operatorname{diag}(a_1 g_1, \dots, a_n g_n)\| = \mathbb{E} \max |a_i g_i|.$$

Other strategy.

[Van Handel 2017] A strategy based on some comparison of Gaussian processes.

Conjecture is proved when (a_{ij}^2) is a symmetric positive matrix !

But not in general.

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Other operator norms.

$X = (a_{ij}g_{ij}) : \ell_{p^*}^n \rightarrow \ell_q^n$ where $1 < p^* \leq 2 \leq q$.

$$\|X\|_{p^* \rightarrow q} \geq \max \left(\max_i |R_i|_p, \max_j |C_j|_q \right)$$

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$$\|X\|_{p^* \rightarrow q} \geq \max \left(\max_i |R_i|_p, \max_j |C_j|_q \right)$$

Conjecture. Let p be the conjugate of p^* then

$$\mathbb{E} \|X\|_{p^* \rightarrow q} \approx_{p,q} \mathbb{E} \max_i |R_i|_p + \mathbb{E} \max_j |C_j|_q$$

$$\approx_{p,q} \max_i \left(\sum_j |a_{ij}|^p \right)^{1/p} + \max_j \left(\sum_i |a_{ij}|^q \right)^{1/q} + \mathbb{E} \max |a_{ij}g_{ij}|$$

Other strategy - Other operator norms

Theorem [G-Hinrichs-Litvak-Prochno]. We have

$$\mathbb{E} \|X\|_{p^* \rightarrow q} \lesssim_{p,q} (\log n)^{1/q} \left(\max_i \left(\sum_j |a_{ij}|^p \right)^{1/p} + \mathbb{E} \max |a_{ij} g_{ij}| \right) \\ + \max_j \left(\sum_i |a_{ij}|^q \right)^{1/q} .$$

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Observe that

$$\begin{aligned} \mathbb{E} \|X\|_{p^* \rightarrow q} &= \mathbb{E} \sup_{y \in B_{p^*}^n} \left(\sum_i |\langle R_i, y \rangle|^q \right)^{1/q} \leq \left(\mathbb{E} \sup_{y \in B_{p^*}^n} \sum_i |\langle R_i, y \rangle|^q \right)^{1/q} \\ &\leq \left(\mathbb{E} \sup_{y \in B_{p^*}^n} \sum_i |\langle R_i, y \rangle|^q - \mathbb{E} |\langle R_i, y \rangle|^q \right)^{1/q} + \sup_{y \in B_{p^*}^n} (\mathbb{E} |\langle R_i, y \rangle|^q)^{1/q}. \end{aligned}$$

Other strategy - Other operator norms

$$\begin{aligned} & \mathbb{E} \|X\|_{p^* \rightarrow q} \leq \\ & \leq \left(\mathbb{E} \sup_{y \in B_{p^*}^n} \sum_i |\langle R_i, y \rangle|^q - \mathbb{E}' |\langle R'_i, y \rangle|^q \right)^{1/q} + \sup_{y \in B_{p^*}^n} (\mathbb{E} |\langle R_i, y \rangle|^q)^{1/q}. \\ & \leq \left(\mathbb{E} \mathbb{E}' \sup_{y \in B_{p^*}^n} \sum_i |\langle R_i, y \rangle|^q - |\langle R'_i, y \rangle|^q \right)^{1/q} + \sup_{y \in B_{p^*}^n} (\mathbb{E} |\langle R_i, y \rangle|^q)^{1/q}. \\ & \leq \left(2\mathbb{E} \mathbb{E}_\varepsilon \sup_{y \in B_{p^*}^n} \sum_i \varepsilon_i |\langle R_i, y \rangle|^q \right)^{1/q} + \sup_{y \in B_{p^*}^n} (\mathbb{E} |\langle R_i, y \rangle|^q)^{1/q}. \end{aligned}$$

How to evaluate this first term ?

Non commutative Khinchine inequality

The case $p=q=2$.

$$\begin{aligned}\mathbb{E}_\varepsilon \sup_{y \in B_2^n} \sum_{i=1}^n \varepsilon_i |\langle R_i, y \rangle|^2 &= \mathbb{E}_\varepsilon \left\| \sum_{i=1}^n \varepsilon_i R_i \otimes R_i \right\|_{2 \rightarrow 2} \\ &\lesssim \sqrt{\log n} \left(\left\| \sum_{i=1}^n |R_i|_2^2 R_i \otimes R_i \right\|_{2 \rightarrow 2} \right)^{1/2} \\ &\lesssim \sqrt{\log n} \max_i |R_i|_2 \left(\left\| \sum_{i=1}^n R_i \otimes R_i \right\|_{2 \rightarrow 2} \right)^{1/2}\end{aligned}$$

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You have bounded a quantity by its square root !

Some generalization.

Theorem [G-Rudelson '07, G-Mendelson-Pajor-Tomczak '08]

$$\mathbb{E} \sup_{y \in B_{p^*}^n} \left| \sum_{i=1}^n \varepsilon_i |\langle R_i, y \rangle|^q \right| \lesssim_p$$
$$\sqrt{\log n} \max_{1 \leq i \leq n} \|R_i\|_p \sup_{y \in B_{p^*}^n} \left(\sum_{i=1}^n |\langle R_i, y \rangle|^{2(q-1)} \right)^{1/2} .$$

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$$\mathbb{E} \sup_{y \in B_{p^*}^n} \left| \sum_{i=1}^n \varepsilon_i |\langle R_i, y \rangle|^q \right| \lesssim_p \sqrt{\log n} \max_{1 \leq i \leq n} \|R_i\|_p \sup_{y \in B_{p^*}^n} \left(\sum_{i=1}^n |\langle R_i, y \rangle|^{2(q-1)} \right)^{1/2}.$$

Valid because $q \geq 2$, $1 < p^* \leq 2$ and $\ell_{p^*}^n$ has modulus of convexity of power type 2.

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Valid because $q \geq 2$, $1 < p^* \leq 2$ and $\ell_{p^*}^n$ has modulus of convexity of power type 2.

Use concentration of Lipschitz function of a Gaussian vector, and compute all parameters