

COMPUTATIONAL TOPOLOGY

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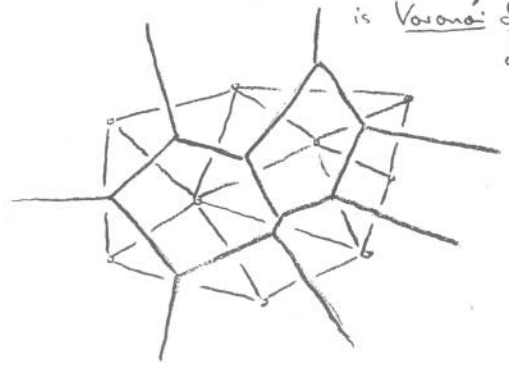
A. DELAUNAY COMPLEXES

A.1 Voronoi tessellations in \mathbb{R}^2
(Descartes 16..)

$X \subseteq \mathbb{R}^2$, finite

$$\text{Vor}(x) = \{a \in \mathbb{R}^2 \mid \|a-x\| \leq \|a-y\| \forall y \in X\}$$

is Voronoi domain of $x \in X$.



1. $\text{Vor}(x)$ is intersection of $|X|-1$ closed half-spaces.
2. Voronoi domains have disjoint interiors and intersect in closed faces.

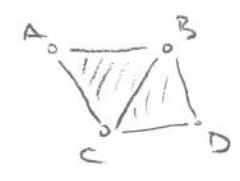
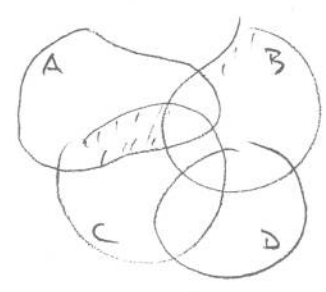
Voronoi tessellation is $\{\text{Vor}(x) \mid x \in X\}$.

A.3 Nerve and Nerve Theorem.

S is collection of sets.

The nerve of S is the system of subcollections with non-empty intersection (Alexander 19..)

$$\text{Nrv } S = \{Q \subseteq S \mid \bigcap Q \neq \emptyset\}$$



NERVE THEM (Leray 1945, Borsuk 1945).

If S is finite and every non-empty common intersection is contractible, then $\text{Nrv } S \simeq U S$.

E.g. if sets in S are convex, then the non-empty common intersections are convex and therefore contractible.

A.2 Delaunay triangulations in \mathbb{R}^2
(Delaunay 1937)

Draw edge xy iff $\text{Vor}(x) \cap \text{Vor}(y) \neq \emptyset$

Draw triangle xyz iff $\text{Vor}(x) \cap \text{Vor}(y) \cap \text{Vor}(z) \neq \emptyset$

1. Edges form a planar graph.

$$\Rightarrow \# \text{edges} < 3|X|$$

$$\Rightarrow \# \text{triangles} < 2|X|$$

Proof: $v - e + f = 2$ (Euler)

Suppose we add edges outside to triangulate also the outside. Then

$$3f = 2e$$

Hence

$$v - e + \frac{2e}{3} = 2 \Rightarrow e = 3v - 6$$

$$v - \frac{3f}{2} + f = 2 \Rightarrow f = 2v - 4 \quad \square$$

2. $Q \subseteq \text{Nrv } S \neq \emptyset$ has an empty circumcircle.

A.4 Abstract simplicial complexes.

- set = set
- collection = set of sets
- system = set of sets of sets

Let U be collection of sets.

A subcollection of U is an abstract simplex

A system of abstract simplices σ is an abstract simplicial complex if $\alpha \in \sigma$ and $\beta \subseteq \alpha$ implies $\beta \in \sigma$.

1. The nerve of a collection of sets is an abstract simplicial complex.
2. The Delaunay triangulation is the nerve of its Voronoi tessellation (or is it not?)

A.5 Geometric realization

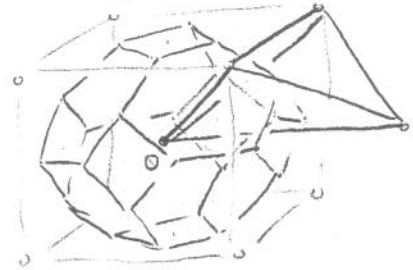
Mapping every set in U to a point in \mathbb{R}^n , we draw every abstract simplex as the convex hull of its points. If we avoid "improper intersections" we get a

(geometric) simplicial complex: a set of simplices, K , in \mathbb{R}^n such that
(i) $Q \in K$ and $P \in Q$ implies $P \in K$,
(ii) $P \cap Q$ is either empty or a common face.

1. Delaunay triangulation is geometric realization of nerve of Voronoi tessellation, mapping $\text{Vor}(x)$ to x (or is it not?)
2. Every a.s.c. of dimension k has a geom. realization in \mathbb{R}^{2k+1} .
3. Underlying space, $|K|$, is union of simplices with induced topology.

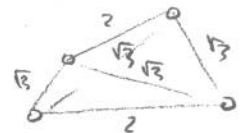
A.7 BCC Lattice

The body centered cubic lattice is set of points $x = (x_1, x_2, x_3)$ with either three even or three odd integer coordinates

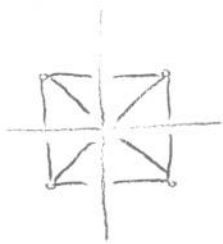


All Voronoi domains are truncates of $\text{Vor}(0) =$ truncated octahedron

All Delaunay tetrahedra are congruent copies of



A.6 General position



Now $\text{Vor}(X)$ is a tetrahedron $\not\subset \mathbb{R}^2$.

- Cope:
1. Don't insist on geom. real.
 2. Draw quadrangle instead.
 3. Assume general position.

Def. $X \subseteq \mathbb{R}^n$ finite is in general position if no $n+2$ points belong to common n -plane and no $n+3$ points belong to common n -sphere.

(custom made for Del. tri.)

A.8 Alpha shapes

$$B_r(x) = \{a \in \mathbb{R}^n \mid \|a-x\| \leq r\}$$

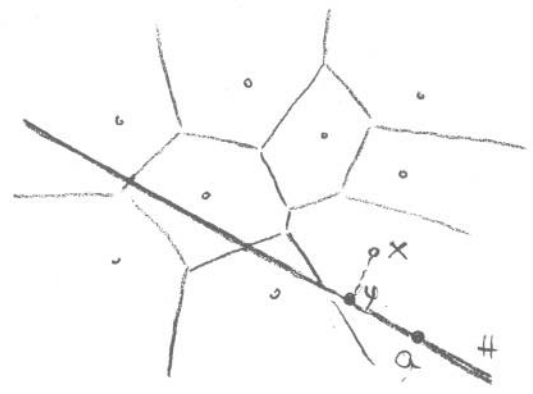
$\text{Del}_r X =$ geometric realization of Nerve $\{ \text{Vor}(x) \cap B_r(x) \mid x \in X \}$
(E. Edelsbrunner, Seidel 1983)



is Delaunay complex of X for r ,
or alpha complex of X for $x=r$,
or alpha shape = underlying space of $\text{Del}_r X$.

A.9 Slices

$X \subseteq \mathbb{R}^n$ finite, $H = k$ -plane in \mathbb{R}^n
Consider $\{Vor(x) \mid x \in X\}$

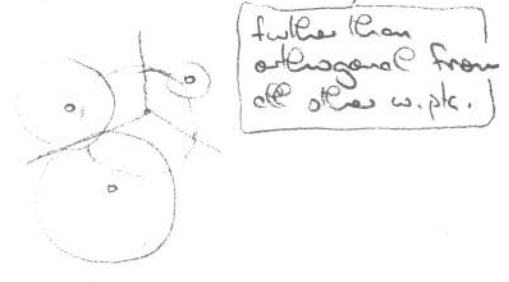


y is orthogonal proj. of x onto H .
 $\|a-x\|^2 = \|a-y\|^2 + \|y-x\|^2$

A.11 Weighted Delaunay triangulation

$D(Y) =$ geometric realization of nerve of $Vor Y$

1. Vertices of $D(Y)$ form subset of weighted points in Y
2. $Q \in D(Y)$ iff Q has an "empty" orthosphere:



A.10 Weighted Voronoi tessellation

$Y \subseteq \mathbb{R}^k \times \mathbb{R}$ finite
 y w is weight

$\pi_y(a) = \|y-a\|^2 - w$ is power distance of a from weighted point $(y, w) \in Y$.

$Vor(y) = \{a \in \mathbb{R}^k \mid \pi_y(a) \leq \pi_x(a) \forall (x, w_x) \in Y\}$



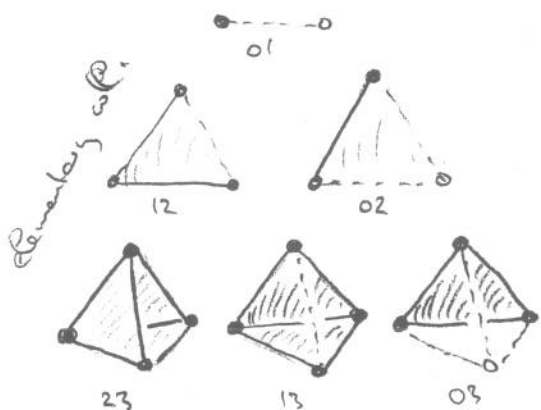
radius = $\sqrt{\text{weight}}$

B. DISCRETE MORSE THEORY

B.1 Collapses

Free simplex in K is proper face of exactly one maximal simplex. Collapse removes:

1. free simplex
2. max. simplex
3. everybody in between



B.3 Deformation retraction

Continuous $D: X \times [0,1] \rightarrow X$ is deformation retraction from X to $Y \subseteq X$ if

- (i) $D(x,0) = x \quad \forall x \in X$
- (ii) $D(x,1) \in Y \quad \forall x \in X$
- (iii) $D(y,t) = y \quad \forall y \in Y, t \in [0,1]$

Then X and Y are hom. equivalent.

$X = \text{triangle}$
 $x = \alpha A + \beta B + \gamma C$

$$\begin{bmatrix} \alpha \\ \beta \\ \gamma \end{bmatrix} \rightarrow \begin{bmatrix} \alpha \\ \beta \\ \gamma \end{bmatrix} + c \begin{bmatrix} 1 \\ -1/2 \\ -1/2 \end{bmatrix}$$

$$c = \min \{1 - \alpha, 2\beta, 2\gamma\}$$

$$= \min \{2\beta, 2\gamma\}$$

$$D(x,t) = (1-t) \begin{bmatrix} \alpha \\ \beta \\ \gamma \end{bmatrix} + t \begin{bmatrix} \alpha + c \\ \beta - c/2 \\ \gamma - c/2 \end{bmatrix}$$

B.2 Barycentric coordinates

$x = \alpha A + \beta B + \gamma C$
 with
 $1 = \alpha + \beta + \gamma$
 $0 \leq \alpha, \beta, \gamma$

$x = (\frac{1}{3}, \frac{1}{3}, \frac{1}{3})$

17

$\alpha + \beta + \gamma = 1$
 $\alpha, \beta, \gamma \geq 0$

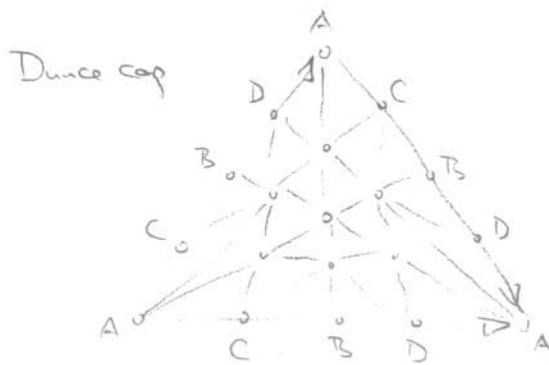
standard 2-simplex

B.4 Collapsibility and contractibility

Simplicial complex K is collapsible if $K \rightarrow pt.$

K is contractible if hom. eq. to point.

Collapsibility $\not\Rightarrow$ Contractibility



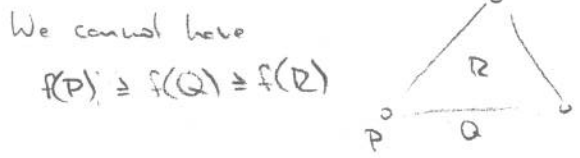
contractible but not collapsible.

B.5 Discrete Morse functions

K is simplicial complex.
 $f: K \rightarrow \mathbb{R}$ is discrete Morse function if

- $f(P) \neq f(Q)$ for at most one face P of Q ,
- $f(P) \leq f(Q)$ for at most one coface P of Q .

(Forman 1998)



Matching formed by exceptional pairs (collapses), and unmatched simplices are critical.

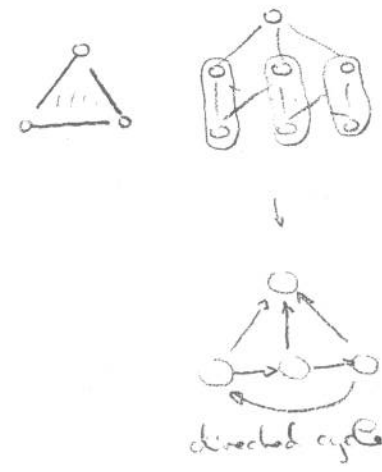
Then: $f: K \rightarrow \mathbb{R}$ is d.M. function, and $K_0 \in K$ s.t. $K \setminus K_0$ contains no critical simplices. Then $K \approx K_0$.

B.5

B.7 Gen. discrete vector fields

K is simplicial complex
Gen. discrete vector field, V , is partition of K into intervals.
 Gray (200.), Freij (201.)

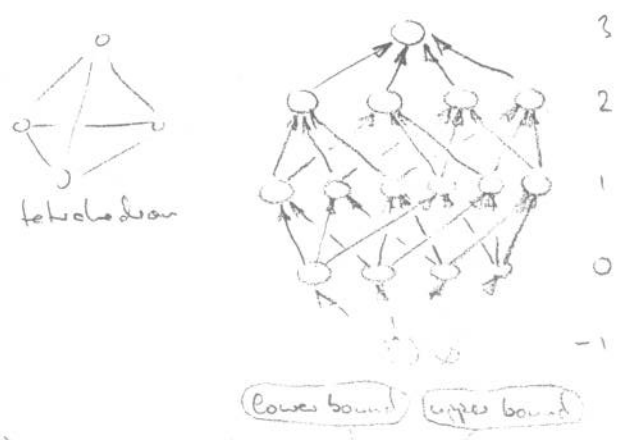
Contract intervals to nodes and keep all directed edges in Hasse diagram.



B.7

B.6 Hasse Diagrams and intervals

Hasse diagram of K is digraph of simplices with containment:



for simplices $L \in U$,
 $[L, U] = \{Q \in K \mid L \subseteq Q \subseteq U\}$ is interval.

- $|[L, U]| = 2^{\dim(L)-\dim(U)}$
- Interval can be partitioned into pairs.

B.6

B.8 Gen. discrete gradient

... is acyclic g.d.v.f.

To prove acyclicity of V find function $f: K \rightarrow \mathbb{R}$ such that $f(P) \leq f(Q)$ whenever $P \subseteq Q$, with equality iff P, Q in same interval.

Then f is generalized discrete Morse function and V is its generalized discrete gradient.

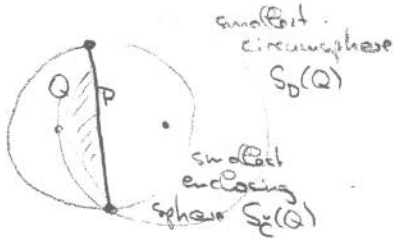
- V can be refined to a matching
- f can be perturbed to a discrete Morse function.

B.8

B.9 Radius function

$f: D \subseteq X \rightarrow \mathbb{R}$ defined by
 $f(Q) = \min \{ r \mid Q \in D_{\leq r} X \}$.

$\Rightarrow D_{\leq r} X = f^{-1}([0, r])$.



- $Q \in D \subseteq X$ is critical simplex of f if $S_D(Q) = S_E(Q)$.
- $Q \in D \subseteq X$ is upper bound of interval if $S_D(Q)$ is empty.
- $P \in Q$ belongs to some interval if all facets $P \in F \subseteq Q$ are visible from center of $S_D(Q)$.

Hence, f is a generalized discrete Morse function.

B.11 Delaunay Čech complexes

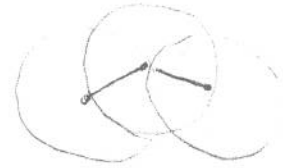
$D_{\leq r} \check{C}ech_r X = D_{\leq r} X \cap \check{C}ech_r X$.

1. $D_{\leq r} X \subseteq D_{\leq r} \check{C}ech_r X$ but they are not necessarily the same.

$Q \in D_{\leq r} X$ if $\bigcap_{x \in Q} [V(x) \cap B_r(x)] \neq \emptyset$

$Q \in D_{\leq r} \check{C}ech_r X$ if

$\bigcap_{x \in Q} V(x) \neq \emptyset$ and $\bigcap_{x \in Q} B_r(x) \neq \emptyset$



2. Delaunay-Čech complex does not grow in intervals.

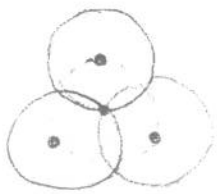
?

B.10 Čech complexes

$X \subseteq \mathbb{R}^n$ finite, $B_r(x)$ is ball of radius $r \geq 0$ centered at $x \in X$.

$\check{C}ech_r X = \{ Q \subseteq X \mid \bigcap_{x \in Q} B_r(x) \neq \emptyset \}$

is Čech complex of X for radius r .



1. $Q \in \check{C}ech_r(X)$ iff radius of $S_{\check{C}}(Q)$ is at most r .

$f_{\check{C}}: \Delta(X) \rightarrow \mathbb{R}$ defined by
 $f_{\check{C}}(Q) = \text{radius of } S_{\check{C}}(Q)$.

$\Rightarrow \check{C}ech_r X = f_{\check{C}}^{-1}([0, r])$.

- $Q \in \Delta(X)$ is lower bound of interval if $S_{\check{C}}(Q) = S_D(Q)$. Conv. upper bound is points inside and on $S_{\check{C}}(Q)$.
- Q is critical simplex of $f_{\check{C}}$ if $S_{\check{C}}(Q) = S_D(Q)$ is empty.

B.12 Collapsing Hierarchy

$X \subseteq \mathbb{R}^n$ finite, $r \geq 0$. Then

$\check{C}ech_r X \hookrightarrow D_{\leq r} \check{C}ech_r X \hookrightarrow D_{\leq r} X \hookrightarrow \text{Wrap}_r X$
 (Bauer, E 2015)

The Wrap complex (for $r = \infty$) was used by Geometric for surface reconstruction.

Hasse diagram \rightarrow poset of intervals

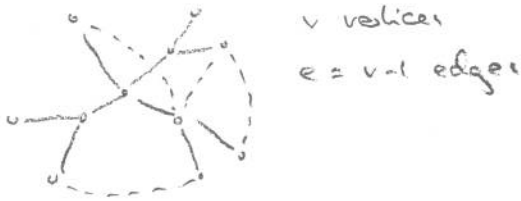


$\text{crit}_r(f) = \text{set of crit. simplices } Q \text{ of } f \text{ with } f(Q) \leq r$.

$\text{Wrap}_r X$ is lower set of $\text{crit}_r(f)$.

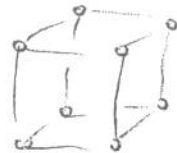
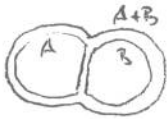
C.1 Cycleomatic numbers

G is simple finite graph,
 G is tree if connected and acyclic.



Every additional edge forms a new cycle

Assuming G is connected,
cycleomatic number is $e - (v - 1)$.



cycle # = 5 because
 every cycle is lin. comb. of 5 basis cycles.

C.3 Homology groups

K is simplicial complex

$C_p(K)$ = group of p -chains

set of p -simplices
 $Q + Q = 0$

$$\partial_p c = \sum_{Q \in c} \partial_p Q$$



$Z_p(K)$ = group of p -cycles

p -chains with
 $\partial_p c = 0$

$B_p(K)$ = group of p -boundaries

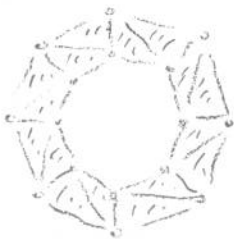
$\exists (p+1)$ -chain d with
 $\partial_{p+1} d = c$

$H_p(K) = Z_p(K) / B_p(K)$ is p -th homology group

$\beta_p(K) = \text{rank } H_p(K)$ is p -th Betti number

C.2 First Betti numbers

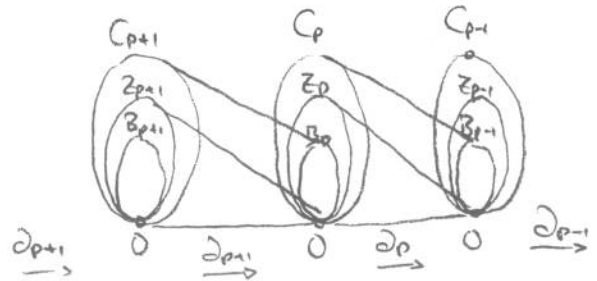
K is simplicial complex,
 $G = K^{(1)}$ is edge skeleton.



2-chain is set of triangles, d .
boundary, ∂d , is edges shared by odd # triangles.

1-cycles c, c' are homologous (same) if $\exists d$ s.t. $c + c' = \partial d$.

$\beta_1(K) = \#$ classes of non-trivial 1-cycles.



$Z_p = \ker \partial_p$ Lemma:
 $B_p = \text{im } \partial_{p+1}$ $\partial_p \circ \partial_{p+1} = 0$

Euler characteristic is

$$\chi(K) = \sum_{p \geq 0} (-1)^p \#p\text{-simplices}$$

Euler-Poincaré Formula:

$$\chi(K) = \sum_{p \geq 0} (-1)^p \beta_p(K)$$



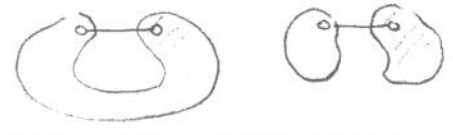
$\beta_0 = 1$
 $\beta_1 = 2$
 $\beta_2 = 1$
 $\chi = 1 - 2 + 1 = 0$

C.4 Incremental Birth # algorithm

K is simplicial complex.

Sort simplices s.t. faces precede
 each simplex: Q_1, Q_2, \dots, Q_m ;
 for $p \geq 0$ do $\beta_p = 0$ endfor; $K_0 = \emptyset$;
 for $i = 1$ to m do
 $K_i = K_{i-1} \cup \{Q_i\}$; $p = \dim Q_i$;
 if Q_i belongs to p -cycle in K_i
 then $\beta_p = \beta_p + 1$
 else $\beta_{p-1} = \beta_{p-1} - 1$
endit

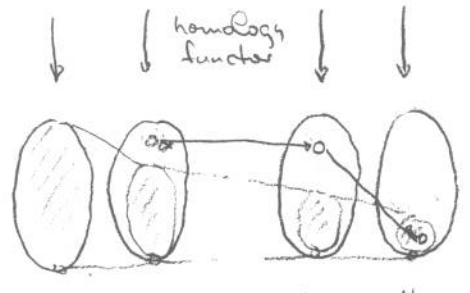
endfor. (E. DeChinada 1995)
 time = $m \log m$ in \mathbb{R}^2 and \mathbb{R}^3 .
 In \mathbb{R}^2 only choice is \mathbb{R}^2 and \mathbb{R}^3 .



C.6 Filtrations and persistence modules

Filtration is nested seq. of s.c.

$$\dots \subseteq K_{i-1} \subseteq K_i \subseteq \dots \subseteq K_{j-1} \subseteq K_j \subseteq \dots$$



$$\dots \rightarrow H_{i-1} \rightarrow H_i \rightarrow \dots \rightarrow H_{j-1} \rightarrow H_j \rightarrow \dots$$

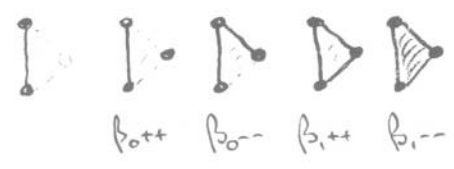
$$\bigoplus_{p \geq 0} H_p(K_{i-1})$$

with maps
 $f_i^j : H_i \rightarrow H_j$
 induced by $K_i \subseteq K_j$,
 is persistence module.

C.5 Almost collapses

$K = \mathbb{D} \circ P X$,
 $p(Q_1) \subseteq p(Q_2) \subseteq \dots \subseteq p(Q_m)$
 ties broken by base dim. first.

E.g. anti-collapse (oblate triangle):



at 4 steps at same radius
 cancelling effects on β the

Similar for slightly acute triangle
 except that steps for slightly diff.
 radius.

C.7 Birth and death

Let $\alpha \in H_i$ be a homology class.

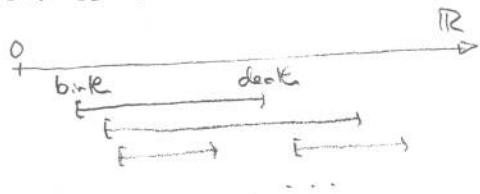
α is born at H_i if $\alpha \notin \text{im } f_{i-1}^i$
 and it dies entering H_j if

$$f_{i-1}^{j-1}(\alpha) \notin \text{im } f_{i-1}^{j-1} \text{ but } f_i^j(\alpha) \in \text{im } f_{i-1}^{j-1}.$$

The persistence of α is $j-i$
 or $p(Q_j) - p(Q_i)$.

C.8 Barcode and persistence diagrams

The barcode of the filtration is a multi-set of intervals:



Equivalently, the persistence diagram is a multi-set of points:



(E. Edelsbrunner, Zomorodian 2000)

Add all diagonal points to diagram for good measure.

C.10 Stability

Functions $f, g: X \rightarrow \mathbb{R}$ are close if all sublevel sets have finite symmetric difference, and there are finitely many hom. crit. values.

Then,

$$W_{\infty}(F, G) \leq \|f - g\|_{\infty}$$

(Cohen-Steiner, Edelsbrunner 2007)

C.9 Bottleneck distance

Given persistence diagrams $F = \text{Dgm}(f)$ and $G = \text{Dgm}(g)$, the bottleneck distance is the length of the longest edge in the minimizing matching:

$$W_{\infty}(F, G) = \inf_{\text{bij: } F \rightarrow G \cup \Delta} \max \|a - \text{bij}(a)\|_{\infty}$$



C.11 Vietoris-Rips complexes

$X \subseteq \mathbb{R}^n$ finite

$\text{Rips}_r X = \{Q \subseteq X \mid \text{diam}(Q) \leq 2r\}$
is Vietoris-Rips complex of X for radius $r \geq 0$.

1. 1-skeleta of $\text{Rips}_r X$ and $\check{C}ech_r X$ are the same.
2. $\text{Rips}_r X$ is clique complex of its 1-skeleton $\Rightarrow \check{C}ech_r X \subseteq \text{Rips}_r X$
3. $\text{Rips}_r X \subseteq \check{C}ech_{r/2} X$.

Map radius to $f(r) = \log_2(r+1)$

$$\log_2(\sqrt{2}r+1) < \log_2(r+1) + \frac{1}{2}$$

Hence

$$W_{\infty}(\check{C}ech, \text{Rips}) < \frac{1}{2}$$