

Helly graphs and Helly groups

9/5/2019 Lecture 2:

Continuation of Isbell's proof:

(8) Lemma 3: If $e(x) \subseteq Y \subseteq E(x)$, and γ is injective, then
 $Y = E(x)$.

Proof: $\alpha: E(x) \rightarrow E(x)$ 1-Lip. s.t. α is ^{identity} idempotent on $e(x)$.

We claim that α is trivial.

Let $f \in E(x)$ and $g := \alpha(f)$.

$$g(x) \stackrel{(by c)}{=} d_\infty(g, e(x)) \leq d_\infty(f, e(x)) \stackrel{(by c)}{=} f(x) \Rightarrow g = f \quad \square$$

1-Lip ∴ injective $Y = E(x)$

Suppose $|X| = n$. Then $e: X \rightarrow \mathbb{R}^n$

$\Delta(x) = \{f: f(x) + f(y) \geq d(x, y)\}$ is a convex unbounded polyhedron.

$E(x) \rightarrow$ union of compact faces of $\Delta(x)$.

$$\text{i.e., } E(x) = \bigcap_{\substack{e(x) \\ x \in X}} \left(\bigcup_{\substack{e(y) \\ y \in X}} I(e(x), e(y)) \right) \rightarrow \text{octahedron}$$

3. Helly graphs

Thm A: A graph H is Helly iff for any graph G , the pair (G, H) has the extension property.

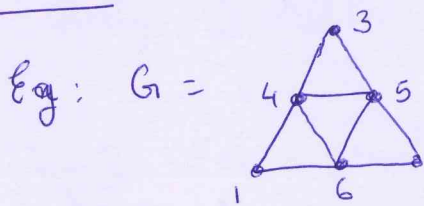
Graphs are always simplicial, and viewed as metric spaces on the vertices.

Thm B: For any graph G , \exists the smallest Helly graph $\text{Helly}(G)$ into which G isometrically embeds.

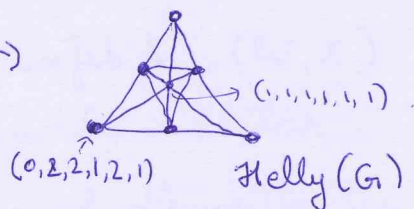
Helly (G) - Hellyfication of G .

$$V(\text{Helly}(G)) = \{ f \in E(G, d_G) \text{ integer valued} \}$$

$$E(\text{Helly}(G)) = \{ f, g : d_\infty(f, g) = 1 \}$$



Kuratowski image of 1: $(0, 2; 2, 1, 2, 1)$.



Cubulation

A wall-space is a set X with a collection of walls

$$\mathcal{W} = \{ (A_i, A'_i) : A_i \cap A'_i = \emptyset, A_i \cup A'_i = X \}$$
 and \mathcal{W} such that

$\forall x, y \in X, \exists$ finitely many walls separating them.

For any wall-space, one associates canonically a CAT(0) cube complex.

Half-spaces of CAT(0) cube complexes are components of hyperplane complements.

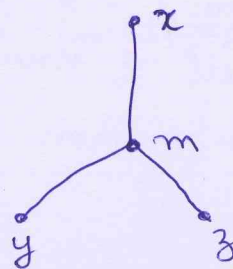
Half-spaces are gated sets in the 1-skeleton (with graph metric).

Median graph: Modular graphs with uniqueness of medians

i.e. $m: V^3 \rightarrow V$

$$(x, y, z) \mapsto \text{median of } \{x, y, z\}$$

median: point in $[x, y] \cap [y, z] \cap [z, x]$



$(V, m) \rightarrow$ median algebra

A median algebra is a set along with a ternary operation m

s.t. 1) $m(u, v, v) = v$.

2) $m(u, v, w) = m(u, w, v) = m(v, u, w)$

3) $m(u, v, m(u, w, x)) = m(u, m(u, v, w), x)$

Given a wall-space (X, \mathcal{W}) , define $d_{\mathcal{W}}(x, y) = \#$ walls separating x & y .

$(X, d_{\mathcal{W}}) \xrightarrow{\alpha} \text{Boolean cube}$

Fix $x_0 \in X$. $\alpha(x) =$ all walls separating x_0 from x .

$$d_{\mathcal{W}}(x, y) = |\alpha(x) \Delta \alpha(y)|$$

Given a graph G , $X_{\text{clique}}(G)$, the clique complex of G is obtained by filling each clique with a simplex (whose 1-skeleton is the clique).

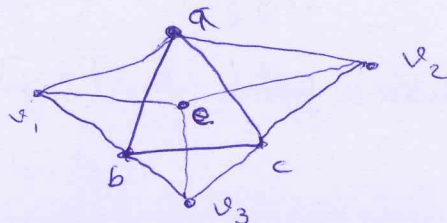
$$X_{\Delta}(G) = X_{\text{clique}}^{(2)}(G)$$

$X_{\text{cube}}(G) =$ fill in cubes for each 1-skeleton of cube in G .

$$X_{\square}(G) = X_{\text{cube}}^{(2)}(G)$$

Helly graphs and clique-Helly graphs.

(Berge-Duchet) Lemma 1: G is clique-Helly iff \forall triangle $T = \{a, b, c\}$, the vertices v having at least 2 neighbours in T have at least one common neighbour e .



^{Chalopin}
^{Chepoi}
^{Hirai}
^{exida}
Thm 1 (CCHO) Let G be \bullet (finitely) clique-Helly and let \tilde{G} be the 1-skeleton of the universal cover \tilde{X} of $X = X_{\Delta}(G)$. Then \tilde{G} is Helly. Consequently, a graph is Helly iff $X_{\Delta}(G)$ is simply-connected and G clique-Helly.

Thm 2: For a graph G , TFAE

- (i) G is Helly.
- (ii) G is 1-Helly and weakly-modular.
- (iii) G is clique-Helly and dismantlable. (G loc-finite).
- (iv) G is clique-Helly and $X_{\Delta}(G)$ is simply connected.
- (v) G is clique-Helly ~~and~~ ^{with} contractible $X_{\text{clique}}(G)$ in the case of finite dimensional $X_{\text{clique}}(G)$.

1-Helly: Helly property for 1-balls.

Properties of Helly graphs

1. Helly \Rightarrow clique-Helly.

Let C be maximal clique. Then $C = \bigcap_{w \in C} B_1(w)$.

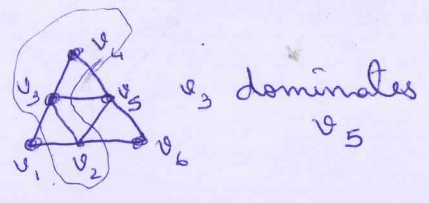
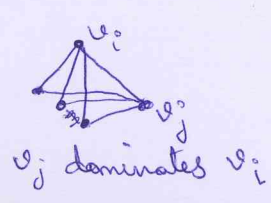
Since Helly property holds for families produced by intersection, we are done.

2. Helly \Rightarrow dismantlability. (assume locally finite G).

A dismantling order on G is a total order \prec on vertices.

v_1, v_2, \dots s.t. $\forall v_i, \exists v_j \prec v_i, j < i$ so that if ~~$v_k \prec v_i, v_k \sim v_i$~~

~~$v_k \prec v_i, v_k \sim v_i \Rightarrow v_k \prec v_j \text{ or } k=j, v_k \sim v_j$~~ (\sim : adjacent).



Lemma 2: If G is dismantlable, then $X_{\text{clique}}(G)$ is contractible. (\therefore also simply-connected).

Lemma 3: If G is Helly, then any base-point order is a dismantling.

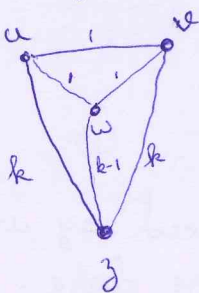
Let $v \in G$ be a vertex. A base-point ordering on G at v is any labelling of the vertices of G such that for any vertices u, w ,
 $d(u, v) < d(w, v) \Rightarrow u < w$.

There is no condition on the order of u, w if $d(u, v) = d(w, v)$.

3. Helly \Rightarrow weak (pseudo) - modularity:

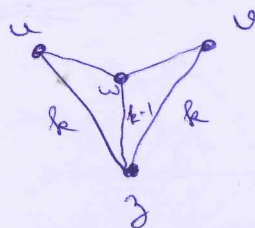
A graph is weakly-modular if it satisfies

1) (TC) triangle condition: $\forall u, v, z$ s.t. $u \sim v$, $d(u, z) = d(v, z) = k$,
 $\exists w \sim u, v$ s.t. $d(w, z) = k - 1$



2) (QC) Quadrangle condition: $\exists w \sim v, v'$

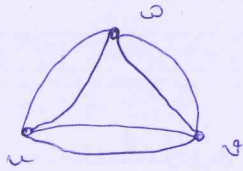
Pseudo-modularity: u, v , $d(u, v) \leq 2$, $d(u, z) = d(v, z) = k \Rightarrow \exists w$
s.t. $w \sim u, v$, $d(w, z) = k - 1$.



(u, v, w) is a metric triangle if $I(u, v) \cap I(u, w) = \{u\}$

$$I(v, u) \cap I(v, w) = \{v\}$$

$$I(w, u) \cap I(w, v) = \{w\}$$



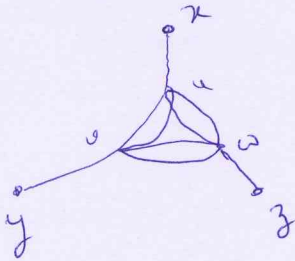
Quasi-median of x, y, z : Metric triangle (u, v, w) such that

$$1) d(x, y) = d(x, u) + d(u, v) + d(v, y)$$

$$2) d(y, z) = d(y, v) + d(v, w) + d(w, z)$$

$$3) d(z, x) = d(z, w) + d(w, u) + d(u, x)$$

$$4) d(u, v) = d(u, w) = d(v, w)$$



Lemma 4: \forall graph G , \forall triplet x, y, z , \exists at least one quasi-median.

Prop: G is weakly-modular $(\Rightarrow) \forall$ metric triangle (u, v, w) ,

$$\forall x, x' \in I(v, w), d(u, x) = d(u, x')$$

In weakly modular graphs, metric triangles are strongly equilateral.

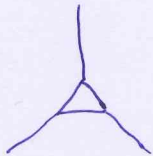


Modular graphs \equiv graphs where all metric triangles are vertices.

Lemma 5: Pseudomodular \Rightarrow metric Δ s have size 0 or 1.



or



~~Proof~~:

Lemma 6: If G is pseudo-modular \Rightarrow ~~any~~ ^{any} three pairwise intersecting balls intersect.

Lemma 7: Helly \Rightarrow pseudo-modular \Rightarrow weakly modular.

Sketch of proof of Thm 1