

# Which problems can be solved by randomized algorithms?

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Let  $F$  and  $G$  be separable Banach spaces. A linear problem is given by a continuous linear operator  $S : F \rightarrow G$ . The aim is to approximate  $S(f)$  for every  $f \in F$  by means of an algorithm. Approximation error of an algorithm  $A$  is given by

$$e^{det}(A, S) = \sup_{f \in F, \|f\|=1} \|S(f) - A(f)\|_G \quad \text{if } A \text{ is a deterministic algorithm,}$$

and

$$e^{ran}(A, S) = \sup_{f \in F, \|f\|=1} \mathbb{E} [\|S(f) - A(f)\|_G] \quad \text{if } A \text{ is a deterministic algorithm.}$$

Denote by  $e^{det}(n, S)$  the minimal error of a deterministic algorithm using at most  $n$  information evaluation. We say that the problem is solvable in a deterministic setting when

$$\lim_{n \rightarrow \infty} e^{det}(n, S) = 0.$$

We define  $e^{ran}(S, n)$  analogously and say that the problem  $S$  is solvable in the randomized setting if

$$\lim_{n \rightarrow \infty} e^{ran}(n, S) = 0.$$

In this talk we focus on the situation when we are allowed to use any continuous linear functional as an information.

The main question is: which problems are solvable? It is well-known that in the deterministic setting the problem given by  $S$  is solvable exactly when  $S$  is compact. In this talk we give a partial answer to the question in the randomized setting. The two main results are the following:

1. If the operator  $S$  is not finitely strictly singular then the problem given by  $S$  is not solvable in the randomized setting.
2. If  $G$  is additionally assumed to be a Hilbert space then there is a large class of Banach spaces  $F$  (characterized by the behaviour of Gaussian measures) for which  $S : F \rightarrow G$  is solvable in the randomized setting if and only if  $S$  is compact.