

Subsequences of binary recursive sequences

by

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1. For any sequence of rational integers $u_0, u_1, \dots, u_m, \dots$ satisfying $u_m = ru_{m-1} + su_{m-2}$ for integers r, s with $r^2 + 4s \neq 0$, we have

$$u_m = a\alpha^m + b\beta^m, \quad m = 0, 1, 2, \dots$$

where α and β are the roots of the polynomial $x^2 - rx - s$ associated to the sequence $\{u_m\}$ and

$$a = \frac{u_0\beta - u_1}{\beta - \alpha}, \quad b = \frac{u_1 - u_0\alpha}{\beta - \alpha}.$$

The sequence $\{u_m\}$ is said to be *non-degenerate binary recursive sequence* if a, b, α and β are non-zero and a/β is not a root of unity. We prove:

THEOREM 1. Let $u_0, u_1, \dots, u_m, \dots$ be a non-degenerate binary recursive sequence. Then there exists an effectively computable number $N > 0$ depending only on the sequence $\{u_m\}$ such that the equation

$$(1) \quad u_m = u_n$$

has no solution in non-negative integers m, n with $\max(m, n) > N$ and $m \neq n$.

In particular, if $d > \max_{0 \leq m \leq N} u_m$, then the sequence $\{u_m\}$ can assume the value d at most once only. Compare this with a theorem of Kubota [3]. The proofs of Theorem 1 and all other results of this paper depend on the theory of linear forms in logarithms.

Proof. Put $K = Q(\alpha)$ and denote by c_1, c_2, \dots effectively computable positive constants depending only on the sequence $\{u_m\}$. Suppose that non-negative integers m, n with $m \neq n$ satisfy (1). It is no loss of generality to assume that $m > n$. We assume that $m > c_1$ with c_1 sufficiently large. If $|\alpha| < |\beta|$ or $|\alpha| > |\beta|$, the theorem follows trivially. Thus we assume that $|\alpha| = |\beta|$. Since α/β is not a root of unity, the numbers α/β and β/α are

not integers in K . Thus there exists a prime ideal p in K such that $\text{ord}_p(\alpha/\beta) > 0$. Now from (1), we have

$$\alpha^n(a^{m-n}-1) = b\beta^n(1-\beta^{m-n}).$$

Further

$$c_2 + c_3 n = \text{ord}_p(b(1-\beta^{m-n})) < c_4 + c_5 \log m.$$

For the last inequality, see Hasse [2], p. 168. Hence, by (1),

$$(2) \quad |aa^m + b\beta^m| < m^{c_6}.$$

The left-hand side of the above inequality does not vanish if c_1 is large enough. Further, by a theorem of Baker [1],

$$(3) \quad |aa^m + b\beta^m| > |a|^m m^{-c_7}.$$

Combining (2), (3) and observing $|a| > 1$, we find that $m < c_8$. This completes the proof of Theorem 1.

2. Now we consider the equation $u_m = v_n$ where

$$u_m = aa^m + b\beta^m, \quad v_n = c\gamma^n + d\delta^n \quad (m, n = 0, 1, 2, \dots)$$

are non-degenerate binary recursive sequences whose associated polynomials have real roots. Assume that

$$(4) \quad |\alpha| < |\beta|, \quad |\gamma| < |\delta|, \quad |\beta| > 1, \quad |\delta| > 1.$$

Then we have:

LEMMA. There exists an effectively computable positive number N_0 depending only on $a, b, c, d, \alpha, \beta, \gamma$ and δ such that if m, n is a solution of $u_m = v_n$ in non-negative integers, then m, n is either a solution of the system of equations

$$aa^m = c\gamma^n, \quad b\beta^m = d\delta^n$$

or

$$\max(m, n) \leq N_0.$$

This is a direct consequence of a theorem of Baker [1]. Now we give two applications of the lemma.

THEOREM 2. Let $u_0, u_1, \dots, u_m, \dots$ and $v_0, v_1, \dots, v_n, \dots$ be non-degenerate binary recursive sequences whose associated polynomials have real roots. Suppose that $\{v_n\}$ is a subsequence of $\{u_m\}$. Then there exist integers $p > 0$ and f depending only on these sequences such that

$$v_n = u_{pn+f}$$

for all integers $n > |f|$.

Similar result was proved by Mignotte [4].

Proof. Since $\{v_n\}$ is a subsequence of $\{u_m\}$, there exist positive integers m_1 and m_2 such that

$$u_{m_1} = v_{N_0+1}, \quad u_{m_2} = v_{N_0+2}.$$

Now the lemma gives

$$\begin{aligned} aa^{m_1} &= c\gamma^{N_0+1}, & b\beta^{m_1} &= d\delta^{N_0+1} \\ aa^{m_2} &= c\gamma^{N_0+2}, & b\beta^{m_2} &= d\delta^{N_0+2}. \end{aligned}$$

These equations give

$$\gamma = a^p, \quad \delta = \beta^p$$

with $p = m_2 - m_1$. From (4), we find that $p > 0$. Put

$$f = m_1 - p(N_0 + 1).$$

Now, for $n > |f|$, we have

$$v_n = c\gamma^n + d\delta^n = aa^{pn+f} + b\beta^{pn+f} = u_{pn+f}.$$

This completes the proof of Theorem 2.

Now we apply the lemma to prove:

THEOREM 3. Suppose that $u_0, u_1, \dots, u_m, \dots$ and $v_0, v_1, \dots, v_n, \dots$ are non-degenerate binary recursive sequences with infinitely many terms in common. Assume that the roots of their associated polynomials are positive. Then they are subsequences of a sequence $W_0, W_1, \dots, W_i, \dots$ where $\{r'W_i\}$ is a non-degenerate binary recursive sequence for a fixed integer r' depending only on the sequences $\{u_m\}$ and $\{v_n\}$. Further the polynomial associated to $\{r'W_i\}$ has positive roots.

Proof. Proceed similarly as in Theorem 2 to conclude that there exist relatively coprime positive integers p and q such that $\alpha^p = \gamma^q, \beta^p = \delta^q$. Since $(p, q) = 1$, there exist integers u, v such that $pu + qv = 1$. Therefore

$$\alpha = \alpha^{pu+qv} = (\gamma^u \alpha^v)^q = \eta^q, \quad \text{where } \eta = \gamma^u \alpha^v.$$

Now

$$\gamma^q = \alpha^p = \eta^{pq},$$

which implies that

$$\gamma = \eta^p.$$

Similarly

$$\beta = \delta^q, \quad \delta = \eta^p, \quad \text{where } \eta = \delta^u \beta^v.$$

There exist positive integers m_0, n_0 such that $aa^{m_0} = c\gamma^{n_0}$ and $b\beta^{m_0} = d\delta^{n_0}$. This follows from the lemma. Put $D = m_0q - n_0p$. Then

$$c/a = \eta^D, \quad d/b = \eta^D.$$

For $l = 0, 1, 2, \dots$, put

$$W_l = \begin{cases} a\eta^l + b\varrho^l & \text{if } D \geq 0, \\ c\eta^l + d\varrho^l & \text{if } D < 0. \end{cases}$$

Then

$$u_m = \begin{cases} W_{mq}, & D \geq 0, \\ W_{mq-D}; & D < 0, \end{cases} \quad v_n = \begin{cases} W_{np+D}, & D \geq 0, \\ W_{np}, & D < 0, \end{cases}$$

for $m = 0, 1, 2, \dots$ and $n = 0, 1, 2, \dots$

Clearly $\{u_m\}$ and $\{v_n\}$ are subsequences of $\{W_l\}$. Put $r' = (\beta - \alpha) \times (\delta - \gamma)$. It is easy to check that the sequence $\{r'W_l\}$ is a non-degenerate binary recursive sequence and its associated polynomial has positive roots.

Remarks. (i) In fact the lemma can be strengthened as follows:

Let $\{u_m\}$ and $\{v_n\}$ be non-degenerate binary recursive sequences. Suppose that their associated polynomials have real roots. Then the equation $u_m = v_n$ has finitely many solutions in non-negative integers m, n if and only if the system

$$aa^m = c\gamma^n, \quad b\beta^m = d\delta^n$$

has at most one solution in non-negative integers m, n . Moreover the result is effective.

(ii) It will be very interesting to prove the lemma when the associated polynomials of the sequences $\{u_m\}$ and $\{v_n\}$ have complex roots.

References

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Meilleures approximations d'une forme linéaire cubique

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I. Introduction, notations. Le développement d'un nombre réel a en fraction continue permet de bien connaître les approximations rationnelles de a ou de la forme linéaire $qa - p$. Si p_n/q_n est une réduite de a on a les propriétés

$$(i) |q_n a - p_n| < 1/q_n, n \geq 0,$$

$$(ii) |qa - p| < |q_n a - p_n| \Rightarrow |q| > q_n,$$

(iii) Le développement est périodique pour $a = \sqrt[D]{D}$ (D entier non carré).

Beaucoup d'auteurs (Jacobi, O. Perron, N. Pipping, V. Brun, G. Szekeres, ...) ont tenté de généraliser cette théorie à plusieurs nombres réels.

Nous renvoyons à G. Szekeres [5], p. 113-117, pour la discussion des propriétés que l'on peut demander à de tels algorithmes.

Dans cet article nous proposons une nouvelle définition de la notion de meilleure approximation de zéro par une forme linéaire cubique, $p_0 + p_1 a_1 + p_2 a_2$. Nous montrons au paragraphe II que l'algorithme fournitant ces approximations peut être considéré comme une généralisation

des fractions continues. Le développement de $a_1 = \sqrt[3]{m}$, $a_2 = \sqrt[3]{m^2}$, où m est un entier naturel distinct d'un cube, est périodique (théorème 1). Au paragraphe IV on étudie les propriétés générales de cet algorithme appliqué à deux nombres réels a_1, a_2 linéairement indépendants avec 1 et on montre que les approximations de zéro par la forme linéaire $p_0 + p_1 a_1 + p_2 a_2$ et les approximations simultanées de a_1 et a_2 qui en résultent vérifient le meilleur degré d'approximation possible.

Soient a_1, a_2 deux nombres réels supérieurs à 1 (cette restriction n'est pas fondamentale), p_0, p_1, p_2 trois entiers. Posons:

$$\Omega = (1, p_1, a_2), \quad P = (p_0, p_1, p_2),$$

$$(I.1) \quad \psi(P) = P \cdot \Omega = p_0 + p_1 a_1 + p_2 a_2,$$

$$\mathcal{E}(P) = \frac{1}{2}((p_0 - p_1 a_1)^2 + (p_1 a_1 - p_2 a_2)^2 + (p_2 a_2 - p_0)^2).$$