

## A rigid Boolean algebra that admits the elimination of $Q_1^2$

by

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**Abstract.** Using  $\diamond$ , we construct a rigid atomless Boolean algebra that has no uncountable antichain and that admits the elimination of the Malitz quantifier  $Q_1^2$ .

**1. Introduction.** Malitz quantifiers are introduced in [Mag-Mal]. Let us recall the semantics of  $Q_\alpha^n$ ,  $n \geq 1$ ,  $\alpha \in \text{ORD}$ :  $\mathfrak{A} \models Q_\alpha^n \bar{x} \phi(\bar{a}, \bar{x})$  iff there is a subset  $H$  of  $A$  such that  $\text{card}(H) \geq \aleph_\alpha$  and  $\mathfrak{A} \models \phi(\bar{a}, h)$  for all pairwise different  $h_0, h_1, \dots, h_{n-1} \in H$ . Such a set  $H$  is called a *homogeneous set* for  $\phi(\bar{a}, \bar{x})$ . Baldwin and Kueker [Bal-Ku], Rothmaler and Tuschik [Ro-Tu], Bürger [Bü] and Koepke [Ko] consider the question of elimination of some of these quantifiers in certain theories or structures. [Ro-Tu] shows that any saturated model allows the elimination of all  $Q_\alpha^n$ ,  $\alpha \in \text{ORD}$ ,  $n \geq 1$ .

Saturated models with two elements of the same type are not rigid. On the other hand, there are  $\mathcal{L}_{\omega\omega}(Q_1^2)$ -sentences  $\phi$  that have only rigid models and that are satisfiable under CH (see [Ot], [Mil]). We consider

$$\begin{aligned} \phi := & \text{“the structure is a Boolean algebra with } 0 \neq 1\text{”} \\ & \wedge \forall x(x \neq 0 \rightarrow Q_1 y y \subseteq x) \wedge \neg Q_1^2 x y x \not\subseteq y. \end{aligned}$$

[Ba-Ko, Theorem 5(a)] shows that all models of  $\phi$  are rigid. The search for a model of  $\phi$  that contains two different elements of the same  $\mathcal{L}_{\omega\omega}(Q_1^2)$ -type leads, under  $\diamond$ , to a model of  $\phi$  that admits the elimination of  $Q_1^2$  and in which therefore any two elements  $\neq 0, 1$  have the same  $\mathcal{L}_{\omega\omega}(Q_1^2)$ -type.

In ZFC +  $\diamond$  and even in ZFC + CH there are various constructions of uncountable Boolean algebras with no uncountable antichains and with some other algebraic properties (see [Ba-Ko], [Sh], [Ru], but also [Ba]). In the course of showing that additional tasks may be fulfilled along the way given in [Ba-Ko], we get a partition of all formulas  $\phi(z, x, y) \in \mathcal{L}_{\omega\omega}(Q_1^2)$ ,  $r \in \omega$ , into two classes  $\Phi_1$  and  $\Phi_2$  such that

1. The methods of [Ba-Ko] are applicable to any  $\phi(\overset{r}{z}, x, y) \in \Phi_1$ . They will allow us to show that the homogeneous sets for any  $\phi(\overset{r}{z}, x, y) \in \Phi_1$  will grow only during countably many steps in the chain which we build in the next section.

2. For any Boolean algebra  $\mathfrak{A}$  with  $\mathfrak{A} \models \forall x \neq 0 Q_1 y y \subseteq x$  and any  $\phi(\overset{r}{z}, x, y) \in \Phi_2$ :  $\mathfrak{A} \models \exists \overset{r}{z} Q_1^2 xy \phi(\overset{r}{z}, x, y)$ .

“ $\phi(\overset{r}{z}, x, y) \in \Phi_1$ ” will be shown to be equivalent under the first order theory of atomless Boolean algebras to a first order formula with its free variables among  $z_0, z_1, \dots, z_{r-1}$ . The consideration of the possible quantifierfree types of the  $\overset{r}{z}$  leads to a procedure for eliminating  $Q_1^2$ .

## 2. The construction

*Notation.* We will use  $\mathfrak{A}, \mathfrak{B}, \mathfrak{B}_\alpha$  to denote Boolean algebras. Boolean algebras are considered as  $\tau_{BA}$ -structures with  $\tau_{BA} = \{\cap, \cup, -, 0, 1\}$ .  $x \subseteq y$  is written for  $x \cap y = x$ ,  $\subset$  means strict inclusion,  $x \setminus y$  is used for  $x \cap (-y)$ .  $\mathcal{P}(\omega)$  denotes the powerset algebra of  $\omega$ . For  $\mathfrak{A} \subseteq \mathcal{P}(\omega)$  we often write  $A$  for  $\mathfrak{A}$ . The interpretations of the  $\tau_{BA}$ -symbols in  $\mathcal{P}(\omega)$  are denoted by the symbols themselves.

$a, b \in A$  are *comparable (in  $\mathfrak{A}$ )* iff  $a \subseteq^{\mathfrak{A}} b$  or  $b \subseteq^{\mathfrak{A}} a$ .  $C \subseteq \mathfrak{A}$  is a *chain* (an *antichain*) iff any two distinct elements of  $C$  are comparable (not comparable). For  $a \subset^{\mathfrak{A}} b \in A$  let  $(a, b)_A := \{c \in A \mid a \subset^{\mathfrak{A}} c \subset^{\mathfrak{A}} b\}$ .

Using  $\diamond$ , we shall construct a Boolean algebra  $\mathfrak{B}$  such that  $\mathfrak{B}$  is a model of the sentence  $\phi$  from the introduction and  $\mathfrak{B}$  admits the elimination of  $Q_1^2$ . As the construction of our Boolean algebra  $\mathfrak{B}$  follows the pattern of [Ba-Ko], we restrict ourselves to a short description, heavily referring to [Ba-Ko].

Inductively on  $\alpha \in \omega_1$ , we shall build a chain  $(\mathfrak{B}_\alpha, M_\alpha)_{\alpha \in \omega_1}$ , where the  $\mathfrak{B}_\alpha$  are countable atomless subalgebras of  $\mathcal{P}(\omega)$  and each  $M_{\alpha+1}$  is a countable collection of pairs  $(M, \phi(\bar{c}, x, y))$ , where  $M \subseteq B_\alpha$  and  $\phi(\bar{c}, x, y)$  is a quantifierfree (qf)  $\mathcal{L}_{\omega\omega}[\tau_{BA}]$ -formula with a property that will be defined later on, and  $\bar{c}$  are elements of  $B_\alpha$ . At limit steps we take unions.  $\mathfrak{B}_{\alpha+1}$  will be the Boolean algebra that is generated by  $B_\alpha \cup \{x_\alpha\}$  in  $\mathcal{P}(\omega)$ , where the  $x_\alpha$  is chosen by the same forcing  $P(B_\alpha)$  as in [Ba-Ko], namely:  $P(B_\alpha) = \{(a, b)_{B_\alpha} \mid a \subset b \in B_\alpha\}$ ,  $(a', b')_{B_\alpha} \leq^{P(B_\alpha)} (a, b)_{B_\alpha}$  iff  $a \subseteq a' \subset b' \subseteq b$ .

We shall define  $D_A(M, \phi(\bar{c}, x, y), e, f)$  and  $M_{\alpha+1}$ . Then we take a  $\{D_A(M, \phi(\bar{c}, x, y), e, f) \mid e, f \in B_\alpha, (M, \phi(\bar{c}, x, y)) \in M_{\alpha+1}\}$ -generic subset  $\{(a_n, b_n) \mid n \in \omega\}$  of  $P(B_\alpha)$  such that  $\{(a_n, b_n) \mid n \in \omega\}$  additionally satisfies the properties described in [Ba-Ko] and set  $x_\alpha = \bigcup \{a_n \mid n \in \omega\}$ . In [Ba-Ko],  $M_{\alpha+1}$  is chosen so that chains and antichains are countable. Our  $M_{\alpha+1}$  differs from that of [Ba-Ko], because we also want all homogeneous sets for

any  $\phi(\bar{z}, x, y) \in \Phi_1$  to be countable. The next items are the generalizations of the corresponding points of [Ba-Ko].

DEFINITION 2.1. Let  $A \subseteq \mathcal{P}(\omega)$  and  $\bar{c}, e, f \in A$ . Let  $\phi(\bar{c}, x, y)$  be qf.

(i)  $D_A(M, \phi(\bar{c}, x, y), e, f) := \{(a, b)_A \in P(A) \mid \text{for any } u \in (a, b)_{\mathcal{P}(\omega)} \text{ one of the following points is true:}$

1.  $(u \cap e) \cup (f \setminus u) \in M$ .
2. There is some  $y \in M$  such that

$$\mathcal{P}(\omega) \models \neg\phi(\bar{c}, (u \cap e) \cup (f \setminus u), y) \vee \neg\phi(\bar{c}, y, (u \cap e) \cup (f \setminus u)).$$

(ii)  $M$  is called *maximally homogeneous for  $\phi(\bar{c}, x, y)$  in  $\mathfrak{A}$*  iff  $M \subseteq A$  is homogeneous for  $\phi(\bar{c}, x, y)$  and for all  $a \in A \setminus M$  there is some  $b \in M$  such that  $\mathfrak{A} \models \neg\phi(\bar{c}, a, b) \vee \neg\phi(\bar{c}, b, a)$ .

(iii)  $\phi(\bar{c}, x, y)$  is *small in  $\mathfrak{A}$*  iff for any  $\emptyset \neq M \subseteq A$  that is maximally homogeneous for  $\phi(\bar{c}, x, y)$  in  $\mathfrak{A}$ ,  $D_A(M, \phi(\bar{c}, x, y), 1, 0)$  is dense in  $P(A)$ .

LEMMA 2.2. Let  $\mathfrak{A} \subseteq \mathcal{P}(\omega)$  be atomless,  $\bar{c} \in A^{<\omega}$ ,  $\phi(\bar{c}, x, y)$  qf and small in  $\mathfrak{A}$ ,  $e, f \in A$  and  $M \neq \emptyset$  be maximally homogeneous for  $\phi(\bar{c}, x, y)$  in  $\mathfrak{A}$ . Then  $D_A(M, \phi(\bar{c}, x, y), e, f)$  is dense in  $P(A)$  for any  $e, f$  in  $A$ .

PROOF. [Ba-Ko, Lemmas 2.3 and 2.4].

Also the proof of the next lemma can be carried out as in [Ba-Ko]: just take a  $u$  for  $\mathfrak{A}$  and  $\bar{M}$  in the same way as they take  $x_\alpha$  for  $\mathfrak{B}_\alpha$  and  $M_{\alpha+1}$ .

LEMMA 2.3. Let  $\mathfrak{A} \subseteq \mathcal{P}(\omega)$  be atomless and countable and let  $\bar{M}$  be a countable subset of

$$\{(M, \phi(\bar{c}, x, y)) \mid \bar{c} \in A^{<\omega}, \phi(\bar{c}, x, y) \in \mathcal{L}_{\omega\omega}[\tau_{BA}] \text{ qf, } \phi(\bar{c}, x, y) \text{ small in } A \\ \text{and } M \text{ is maximally homogeneous for } \phi(\bar{c}, x, y) \text{ in } A\}.$$

Then for any  $(a, b)_A \in P(A)$  there is a  $u \in (a, b)_{\mathcal{P}(\omega)}$  such that:

1.  $u \notin A$ .
2.  $[A \cup \{u\}]^{\mathcal{P}(\omega)}$ , the subalgebra generated by  $A \cup \{u\}$  in  $\mathcal{P}(\omega)$ , is atomless.
3. For any  $(M, \phi(\bar{c}, x, y)) \in \bar{M}$  the set  $M$  is maximally homogeneous for  $\phi(\bar{c}, x, y)$  also in  $[A \cup \{u\}]^{\mathcal{P}(\omega)}$ .

Now using Lemma 2.3 and  $\diamond$ , we can construct our  $\mathfrak{B}$ . Let  $\langle S_\alpha \mid \alpha \in \omega_1 \rangle$  be a  $\diamond$ -sequence. Let  $\langle a_\xi \mid \xi \in \omega_1 \rangle$  be an enumeration of  $\mathcal{P}(\omega)$  in which each element of  $\mathcal{P}(\omega)$  appears  $\omega_1$  times.

In step  $\alpha + 1$ , let  $M_{\alpha+1} = M_\alpha \cup \{(\{a_\xi \mid \xi \in S_\alpha\}, \phi(\bar{c}, x, y)) \mid \{a_\xi \mid \xi \in S_\alpha\} \text{ is a maximally homogeneous set for } \phi(\bar{c}, x, y) \text{ in } \mathfrak{B}_\alpha \text{ and } \phi(\bar{c}, x, y) \text{ is small in } \mathfrak{B}_\alpha \text{ and } \bar{c} \in B_\alpha\}$ . Apply Lemma 2.3 with  $\mathfrak{A} = \mathfrak{B}_\alpha$  and  $\bar{M} = M_{\alpha+1}$  to get an  $x_\alpha$ . Define  $B_{\alpha+1}$  as  $[B_\alpha \cup \{x_\alpha\}]^{\mathcal{P}(\omega)}$ . Let  $\mathfrak{B} = \bigcup \{\mathfrak{B}_\alpha \mid \alpha \in \omega_1\}$ . Take the  $x_\alpha$  so that  $\mathfrak{B} \models \forall x (x \neq 0 \rightarrow Q_1 y y \subseteq x)$ . Then it is easy to see that for any  $\phi(\bar{c}, x, y)$  which is small in every  $\mathfrak{B}_\alpha$  with  $\bar{c} \in B_\alpha$ , we have

$\mathfrak{B} \models \neg Q_1^2 xy \phi(\bar{c}, x, y)$ . In particular,  $\mathfrak{B}$  is a model of  $\phi$  from the introduction (because “ $x \not\subseteq y$ ” is small), hence  $\mathfrak{B}$  is rigid.

**3. Large homogeneous sets.** The aim of this section is to define a mapping

$$\begin{aligned} \text{big} : \bigcup_{r \in \omega} \mathcal{L}_{\omega\omega}[\tau_{BA}]^r(\bar{z}, x, y) &\rightarrow \bigcup_{r \in \omega} \mathcal{L}_{\omega\omega}[\tau_{BA}]^r(\bar{z}), \\ \phi(\bar{z}, x, y) &\mapsto \text{big}(\phi(\bar{z}, x, y))(\bar{z}), \end{aligned}$$

such that for every  $\phi(\bar{z}, x, y) \in \mathcal{L}_{\omega\omega}[\tau_{BA}]$

$$(*) \quad \mathfrak{B} \models \forall \bar{z}^r (Q_1^2 xy \phi(\bar{z}, x, y) \leftrightarrow \text{big}(\phi(\bar{z}, x, y))(\bar{z})).$$

Then  $\Phi_2$  will be

$$\{\phi(\bar{z}, x, y) \mid \text{big}(\phi(\bar{z}, x, y))(\bar{z}) \text{ is valid in any atomless Boolean algebra}\}.$$

In order to simplify the notation we tacitly assume that always the variables  $x$  and  $y$  are intended to be quantified by  $Q_1^2$ .

Let  $\mathfrak{A}$  be any atomless Boolean algebra. Since  $\mathfrak{A}$  admits the elimination of  $\exists$  it is enough to define  $\text{big}$  for quantifierfree  $\phi(\bar{z}, x, y) \in \mathcal{L}_{\omega\omega}[\tau_{BA}]$ .

For any  $\bar{c} \in A$  and qf  $\phi(\bar{c}, x, y)$  there is a qf  $\psi(\bar{c}', x, y)$  such that  $\bar{c}'$  is an (injective) enumeration of the atoms of the subalgebra generated by  $\bar{c}$ , and  $\mathfrak{A} \models \forall xy (\psi(\bar{c}', x, y) \leftrightarrow \phi(\bar{c}, x, y))$ . Also if  $\phi(\bar{z}, x, y)$  is a disjunction  $\bigvee_i (\phi(\bar{z}, x, y) \wedge \psi_i(\bar{z}))$  then knowing  $\chi_i = \text{big}(\phi(\bar{z}, x, y) \wedge \psi_i(\bar{z}))(\bar{z})$  we can define  $\text{big}(\phi(\bar{z}, x, y))(\bar{z})$  to be  $\bigvee_i \chi_i$ . Hence it suffices to define  $\text{big}(\phi(\bar{z}, x, y))(\bar{z})$  only for those qf  $\phi(\bar{z}, x, y)$  that imply that  $\{z_0, \dots, z_{r-1}\}$  is the set of atoms in the subalgebra generated by  $\{z_0, \dots, z_{r-1}\}$ .

If  $H$  is an uncountable homogeneous set for  $\phi(\bar{c}, x, y)$ , then there is an  $\mathcal{L}_{\omega\omega}$ -1-type  $t(\bar{c}, x)$  over  $\bar{c}$  and an uncountable  $H_1 \subseteq H$  such that every element of  $H_1$  has the  $\mathcal{L}_{\omega\omega}$ -1-type  $\text{tp}(x/\bar{c}) = t(\bar{c}, x)$  over  $\bar{c}$ . Hence it is enough to define  $\text{big}$  for the  $\phi(\bar{z}, x, y)$  with the above mentioned property and the additional property that there is an  $\mathcal{L}_{\omega\omega}$ -1-type  $t(\bar{z}, x)$  over  $\bar{z}$  (independent of the assignment  $\bar{c}$  of  $\bar{z}$ , because we consider only  $\bar{c}$  that are atoms in the subalgebra generated by  $\bar{z}$ ) such that

$$\mathfrak{A} \models \forall xy \bar{z}^r (\phi(\bar{z}, x, y) \leftrightarrow (\phi(\bar{z}, x, y) \wedge t(\bar{z}, x) = \text{tp}(x/\bar{z}) \wedge t(\bar{z}, y) = \text{tp}(y/\bar{z}))).$$

We will call such formulas *special*. Finally, note that any  $\mathcal{L}_{\omega\omega}$ -2-type  $t(\bar{c}, x, y)$  over  $\bar{c}$  is determined by the corresponding  $r$ -tuple of the quantifierfree types of  $x \cap c_i, y \cap c_i$  in  $\{a \in A \mid a \subseteq c_i\}$ ,  $i < r$ . For any such type there are 15 possibilities, and under the condition  $\text{tp}(x/\bar{z}) = \text{tp}(y/\bar{z})$  there remain the 9 possibilities not marked with an  $\bullet$  in the table below.

The possibilities for the quantifierfree types of  $x \cap c_i, y \cap c_i, i < r$ , in  $\{a \in A \mid a \subseteq c_i\}$

No.	$x \cap y \cap z_i$	$(-x) \cap (-y) \cap z_i$	$x \cap (-y) \cap z_i$	$(-x) \cap y \cap z_i$	Remarks
0	$\neq 0$	$\neq 0$	$\neq 0$	$\neq 0$	
1	$\neq 0$	$\neq 0$	$\neq 0$	0	
2	$\neq 0$	$\neq 0$	0	$\neq 0$	
3	$\neq 0$	$\neq 0$	0	0	$x \cap z_i =$ $y \cap z_i \neq 0, z_i$
4	$\neq 0$	0	$\neq 0$	$\neq 0$	
•5	$\neq 0$	0	$\neq 0$	0	$x \cap z_i = z_i$ $y \cap z_i \neq z_i$
•6	$\neq 0$	0	0	$\neq 0$	$y \cap z_i = z_i$ $x \cap z_i \neq z_i$
7	$\neq 0$	0	0	0	$x \cap z_i =$ $y \cap z_i = z_i$
8	0	$\neq 0$	$\neq 0$	$\neq 0$	
•9	0	$\neq 0$	$\neq 0$	0	$x \cap z_i \neq 0$ $y \cap z_i = 0$
•10	0	$\neq 0$	0	$\neq 0$	$x \cap z_i = 0$ $y \cap z_i \neq 0$
11	0	$\neq 0$	0	0	$x \cap z_i =$ $y \cap z_i = 0$
12	0	0	$\neq 0$	$\neq 0$	$x \cap z_i \neq 0, z_i$ $y \cap z_i = (-x) \cap z_i$
•13	0	0	$\neq 0$	0	$x \cap z_i = z_i$ $y \cap z_i = 0$
•14	0	0	0	$\neq 0$	$x \cap z_i = 0$ $y \cap z_i = z_i$

Let  $\phi^k(z_i, x \cap z_i, y \cap z_i)$  say “the  $\mathcal{L}_{\omega\omega}$ -type of  $x \cap c_i, y \cap c_i$  over  $c_i$  has number  $k$ ”,  $k = 0, \dots, 14$ . The disjunction  $\phi^{012}(u, v, w) := \phi^0(u, v, w) \vee \phi^1(u, v, w) \vee \phi^2(u, v, w)$  will play an important role in the following.

DEFINITION 3.1. Let  $\phi(z, x, y) \in \mathcal{L}_{\omega\omega}[\tau_{BA}]$  be quantifierfree and be of the special form as described above.

$\text{big}(\phi(z, x, y))(z) =$

$$\exists a \subset b \forall xy \left( \left( a \subseteq x, y \subseteq b \wedge \bigwedge_{i < r} ((b \setminus a) \cap z_i \neq 0 \rightarrow \phi^{012}(z_i, x \cap z_i, y \cap z_i)) \right) \rightarrow \phi(z, x, y) \right).$$

Equivalent to  $\text{big}(\phi(z, x, y))(z)$  is the formula

$$\bigvee_{I_0 \dot{\cup} I_1 \dot{\cup} I_2 \dot{\cup} I_3 = \{0, \dots, r-1\}, I_0 \neq \emptyset} \forall xy \left( \left( \bigwedge_{i \in I_0} \phi^{012}(z_i, x \cap z_i, y \cap z_i) \right) \wedge \bigwedge_{i \in I_1} x \cap z_i = y \cap z_i \neq 0, z_i \right)$$

$$\begin{aligned} & \wedge \bigwedge_{i \in I_2} x \cap z_i = y \cap z_i = 0 \\ & \wedge \bigwedge_{i \in I_3} x \cap z_i = y \cap z_i = z_i \rightarrow \phi(z, x, y)^r, \end{aligned}$$

( $\dot{\cup}$  denotes the disjoint union) which will be useful for the easy direction of (\*):

LEMMA 3.2. *Let  $\mathfrak{A}$  be an atomless Boolean algebra. Let  $\mathfrak{A} \models \forall x \neq 0 Q_1 y y \subseteq x$ , and  $\phi(z, x, y)^r$  be as above. Then  $\mathfrak{A} \models \forall z^r (\text{big}(\phi(z, x, y))^r(z) \rightarrow Q_1^2 xy \phi(z, x, y)^r)$ .*

PROOF. Let  $\mathfrak{A} \models \text{big}(\phi(z, x, y))^r(\dot{c})$ . For  $i \in I_0$  take an uncountable set  $H_i \subseteq (0, c_i)_{\mathfrak{A}}$  such that for any  $x \in H_i$  the relative complement  $c_i \setminus x \notin H_i$ . Let  $\langle h_{i, \alpha} \mid \alpha \in \omega_1 \rangle$  be an injective enumeration of a subset of  $H_i$ . Finally, for  $i \in I_1$  let  $H_i = \{d_i\}$  for some  $d_i$  with  $0 \subset d_i \subset c_i$ , for  $i \in I_2$  let  $H_i = \{0\}$ , and for  $i \in I_3$  let  $H_i = \{c_i\}$ . Then

$$H := \left\{ \bigcup \{h_{i, \alpha} \mid i \in I_0\} \cup \bigcup \{d_i \mid i \in I_1\} \cup \bigcup \{c_i \mid i \in I_3\} \mid \alpha \in \omega_1 \right\}$$

is an uncountable homogeneous set for  $\phi(\dot{c}, x, y)^r$ .

Now for  $\mathfrak{B}$  as in Section 2, we shall prove the other direction of (\*). By the construction, it would suffice to show:

(\*\*) For any enumeration  $\dot{c}$  of the atoms in the subalgebra of  $\mathfrak{B}$  generated by  $\dot{z}$ , if  $\mathfrak{B} \models \neg \text{big}(\phi(\dot{z}, x, y))^r(\dot{c})$ , then  $\phi(\dot{c}, x, y)^r$  is small in every  $\mathfrak{B}_\alpha$  with  $\dot{c} \in B_\alpha$ .

Unfortunately, this is true only for  $\phi(\dot{c}, x, y)^r$  that do not forbid certain equalities of Boolean terms. We introduce some notation and then give a sketch of our proof of the hard direction of (\*).

We say briefly “ $\phi(z, x, y)^r$  is valid” or just “ $\phi$ ” for “ $\phi(z, x, y)^r$  is valid in all atomless Boolean algebras if the assignment of  $\dot{z}$  is an enumeration of the atoms in the subalgebra generated by  $\dot{z}$ ”.  $\phi(z, x, y)^r$  is *satisfiable* or *consistent* if  $\neg \phi(z, x, y)^r$  is not valid.

For a given special  $\phi(\dot{z}, x, y)^r$  set

$$R(\phi) := \{i < r \mid \phi \rightarrow x \cap z_i = y \cap z_i \text{ is not valid}\}.$$

We will define two mappings  $s$  and  $\text{enl}$  from the set of all special  $\phi(\dot{z}, x, y)^r$  into itself. The mapping  $s$  is a technical means used to prove  $\text{enl}(\text{enl}(s(\phi))) \rightarrow \text{enl}(s(\phi))$  (Lemma 3.7) and  $\neg \text{big}(s(\phi)) \rightarrow \neg \text{big}(\text{enl}(s(\phi)))$  (Lemma 3.8). Lemma 3.9 says that (\*\*) is true for formulas of the form  $\text{enl}(s(\phi))$  for some

special  $\phi$ . Hence we get from the construction and from 3.8

$$\mathfrak{B} \models \neg \text{big}(s(\phi))(\vec{c}) \rightarrow \neg Q_1^2 xy \text{ enl}(s(\phi))(\vec{c}, x, y),$$

whence  $s(\phi) \rightarrow \text{enl}(s(\phi))$  and the monotonicity of the quantifier  $Q_1^2$  imply

$$\mathfrak{B} \models \neg \text{big}(s(\phi))(\vec{c}) \rightarrow \neg Q_1^2 xy s(\phi)(\vec{c}, x, y)$$

(Theorem 3.10). Using this result we prove by induction on  $\text{card}(R(\phi))$ , simultaneously for all special formulas  $\phi$ ,

$$\mathfrak{B} \models \neg \text{big}(\phi)(\vec{c}) \rightarrow \neg Q_1^2 xy \phi(\vec{c}, x, y),$$

which will finish the proof of (\*).

In order to simplify the notation, we often suppress the free variables  $(\vec{z}, x, y)$  or  $(z_i, x \cap z_i, y \cap z_i)$ .

**DEFINITION 3.3** (The mapping  $s$ ). For  $R \subseteq r = \{0, 1, \dots, r-1\}$  and for  $\chi(z_i, x \cap z_i, y \cap z_i) \in \mathcal{L}_{\omega\omega}[\tau_{BA}]$  we define

$$s_R(\chi(z_i, x \cap z_i, y \cap z_i)) := \begin{cases} \chi(z_i, x \cap z_i, y \cap z_i) & \text{if } i \notin R \text{ or} \\ \phi^{012}(z_i, x \cap z_i, y \cap z_i) \rightarrow \chi(z_i, x \cap z_i, y \cap z_i) & \\ \text{is valid;} & \\ \chi(z_i, x \cap z_i, y \cap z_i) \wedge x \cap z_i \neq y \cap z_i & \\ \text{else.} & \end{cases}$$

Let  $S = \{\bigwedge_{i < r} \chi_{w,i}(z_i, x \cap z_i, y \cap z_i) \mid w \in W\}$  be a finite set such that for all  $w \in W$  the conjunction  $\bigwedge_{i < r} \chi_{w,i}(z_i, x \cap z_i, y \cap z_i)$  is satisfiable and  $\bigwedge_{i < r} \chi_{w,i}(z_i, x \cap z_i, y \cap z_i) \rightarrow \phi(\vec{z}, x, y)$  is valid, and such that for any satisfiable conjunction  $\delta = \bigwedge_{i < r} \chi'_i(z_i, x \cap z_i, y \cap z_i)$  such that  $\delta \rightarrow \phi(\vec{z}, x, y)$  is valid there is a  $w \in W$  with  $\bigwedge_{i < r} \chi'_i(z_i, x \cap z_i, y \cap z_i) \rightarrow \bigwedge_{i < r} \chi_{w,i}(z_i, x \cap z_i, y \cap z_i)$ . We will call such a set  $S$  a *set of representatives* for  $\phi$ . Given such a set, let  $R = R(\phi)$  and define

$$s(\phi(\vec{z}, x, y)) = \bigvee_{w \in W} \bigwedge_{i < r} s_R(\chi_{w,i}(z_i, x \cap z_i, y \cap z_i)).$$

If  $\models \neg \exists x y z \phi(\vec{z}, x, y)$ , then let  $s(\phi(\vec{z}, x, y))$  be any inconsistent formula.

A brief reflection shows that  $s(\phi)$  is well defined up to logical equivalence: Let  $S' = \{\bigwedge_{i < r} \chi'_{w',i}(z_i, x \cap z_i, y \cap z_i) \mid w' \in W'\}$  be another set of representatives for  $\phi$ .

For  $\bigvee_{w' \in W'} \bigwedge_{i < r} s_R(\chi'_{w',i}) \rightarrow \bigvee_{w \in W} \bigwedge_{i < r} s_R(\chi_{w,i})$ , it suffices to show that for each  $w' \in W'$  there is some  $w \in W$  such that  $\bigwedge_{i < r} s_R(\chi'_{w',i}) \rightarrow \bigwedge_{i < r} s_R(\chi_{w,i})$ . Let  $w' \in W'$  be given. Since  $S$  is a set of representatives for  $\phi$  there is a  $w \in W$  such that  $\bigwedge_{i < r} \chi'_{w',i} \rightarrow \bigwedge_{i < r} \chi_{w,i}$ , which is equivalent to

$\chi'_{w',i} \rightarrow \chi_{w,i}$  for  $i < r$ . Immediately from the definition of  $s_R$ , if  $\chi'_{w',i} \rightarrow \chi_{w,i}$ , then  $s_R(\chi'_{w',i}) \rightarrow s_R(\chi_{w,i})$ . Hence  $\bigwedge_{i < r} s_R(\chi'_{w',i}) \rightarrow \bigwedge_{i < r} s_R(\chi_{w,i})$ .

The other direction follows by symmetry.

**Remark.**  $s(\phi)$  may be unsatisfiable, e.g. for  $\phi = (x \cap z_0 = y \cap z_0 \wedge x \cap z_1 \subset y \cap z_1) \vee (x \cap z_0 \subset y \cap z_0 \wedge x \cap z_1 = y \cap z_1) \wedge \bigwedge_{i=0,1} x \cap z_i \neq z_i, 0 \wedge \bigwedge_{i=0,1} y \cap z_i \neq z_i, 0 \wedge z_0 \cap z_1 = 0 \wedge z_0 \cup z_1 = 1$ .

**DEFINITION 3.4** (The mapping  $\text{enl}$ ). For  $\chi(z_i, x \cap z_i, y \cap z_i) \in \mathcal{L}_{\omega\omega}[\tau_{BA}]$  we define

$$\text{enl}(\chi(z_i, x \cap z_i, y \cap z_i)) := \begin{cases} \chi(z_i, x \cap z_i, y \cap z_i) \\ \vee (x \cap z_i = (-y) \cap z_i \wedge \exists x \chi(z_i, x \cap z_i, y \cap z_i) \\ \wedge \exists y \chi(z_i, x \cap z_i, y \cap z_i)) \\ \text{if } \phi^{012}(z_i, x \cap z_i, y \cap z_i) \rightarrow \chi(z_i, x \cap z_i, y \cap z_i) \\ \text{is not valid;} \\ \chi(z_i, x \cap z_i, y \cap z_i) \vee ((x \cap z_i = (-y) \cap z_i \\ \vee x \cap z_i = y \cap z_i) \wedge \exists x \chi(z_i, x \cap z_i, y \cap z_i) \\ \wedge \exists y \chi(z_i, x \cap z_i, y \cap z_i)) \\ \text{otherwise.} \end{cases}$$

Let  $\{\bigwedge_{i < r} \chi_{w,i}(z_i, x \cap z_i, y \cap z_i) \mid w \in W\}$  be a set of representatives for  $\phi$ . Then set

$$\text{enl}(\phi^r(z, x, y)) = \bigvee_{w \in W} \bigwedge_{i < r} \text{enl}(\chi_{w,i}(z_i, x \cap z_i, y \cap z_i)).$$

If  $\models \neg \exists x y z \phi^r(z, x, y)$ , then let  $\text{enl}(\phi^r(z, x, y))$  be any inconsistent formula.

From the fact that  $\chi'_{w',i} \rightarrow \chi_{w,i}$  implies  $\text{enl}(\chi'_{w',i}) \rightarrow \text{enl}(\chi_{w,i})$ , we conclude by an analogous consideration as above that  $\text{enl}(\phi)$  is well-defined.

In order to apply Lemmas 2.2 and 2.3 we may replace  $\text{enl}(\phi^r(z, x, y))$  by an equivalent (with respect to the theory of atomless Boolean algebras) qf formula.

The next two lemmas collect some properties of  $s$  and  $\text{enl}$  that will be useful in the proofs of 3.7 and of 3.8.

**LEMMA 3.5.** *Let  $\chi_s(z_i, x \cap z_i, y \cap z_i)$ ,  $s = 0, 1$ , be qf and  $R \subseteq r$ .*

- (i)  $(\text{enl}(\chi_0) \vee \text{enl}(\chi_1)) \rightarrow \text{enl}(\chi_0 \vee \chi_1)$ .
- (ii)  $(s_R(\chi_0) \vee s_R(\chi_1)) \rightarrow s_R(\chi_0 \vee \chi_1)$ .

For (iii), (iv) and (v), assume additionally that  $\chi_s(z_i, x \cap z_i, y \cap z_i)$ ,  $s = 0, 1$ , determine the same 1-type  $t(z_i, x \cap z_i)$  of  $x \cap z_i$  over  $z_i$  and of  $y \cap z_i$  over  $z_i$ .

(iii) Assume that, for  $s = 0, 1$ , if not  $\phi^{012}(z_i, x \cap z_i, y \cap z_i) \rightarrow \chi_s(z_i, x \cap z_i, y \cap z_i)$ , then  $\chi_s(z_i, x \cap z_i, y \cap z_i) \rightarrow x \cap z_i \neq y \cap z_i$ . Then  $(\text{enl}(\chi_0) \wedge \text{enl}(\chi_1)) \rightarrow \text{enl}(\chi_0 \wedge \chi_1)$ .

(iv)  $(s_R(\chi_0) \wedge s_R(\chi_1)) \rightarrow s_R(\chi_0 \wedge \chi_1)$ .

(v) Assume that  $\chi_s \rightarrow x \cap z_i = y \cap z_i$  for  $s = 0, 1$  if  $i \notin R$ . Then for any  $i < r$  the formula

$$\begin{aligned} (\text{enl}(s_R(\chi_0))(z_i, x \cap z_i, y \cap z_i) \wedge \text{enl}(s_R(\chi_1))(z_i, x \cap z_i, y \cap z_i)) \\ \rightarrow \text{enl}(s_R(\chi_0 \wedge \chi_1))(z_i, x \cap z_i, y \cap z_i) \end{aligned}$$

is valid.

Proof. (i), (ii)  $\chi_s \rightarrow \chi_0 \vee \chi_1$  implies  $\text{enl}(\chi_s) \rightarrow \text{enl}(\chi_0 \vee \chi_1)$  and  $s_R(\chi_s) \rightarrow s_R(\chi_0 \vee \chi_1)$ .

(iii) Define

$$\phi_=(z_i, x \cap z_i, y \cap z_i) := x \cap z_i = y \cap z_i \wedge t(z_i, x \cap z_i) \quad \text{and}$$

$$\phi_-(z_i, x \cap z_i, y \cap z_i) := x \cap z_i = (-y) \cap z_i \wedge t(z_i, x \cap z_i) \wedge t(z_i, y \cap z_i).$$

Case 1:  $\phi^{012} \rightarrow \chi_s$  for  $s = 0, 1$ . Then  $\phi^{012} \rightarrow \chi_0 \wedge \chi_1$  and  $\text{enl}(\chi_0) \wedge \text{enl}(\chi_1) = (\chi_0 \vee \phi_- \vee \phi_-) \wedge (\chi_1 \vee \phi_- \vee \phi_-) \leftrightarrow (\chi_0 \wedge \chi_1) \vee \phi_- \vee \phi_- = \text{enl}(\chi_0 \wedge \chi_1)$ .

Case 2: Not  $\phi^{012} \rightarrow \chi_s$  for  $s = 0, 1$ . Then not  $\phi^{012} \rightarrow \chi_0 \wedge \chi_1$  and  $\text{enl}(\chi_0) \wedge \text{enl}(\chi_1) = (\chi_0 \vee \phi_-) \wedge (\chi_1 \vee \phi_-) \leftrightarrow (\chi_0 \wedge \chi_1) \vee \phi_- = \text{enl}(\chi_0 \wedge \chi_1)$ .

Case 3:  $\phi^{012} \rightarrow \chi_0$  and not  $\phi^{012} \rightarrow \chi_1$ . Then not  $\phi^{012} \rightarrow \chi_0 \wedge \chi_1$  and  $\text{enl}(\chi_0) \wedge \text{enl}(\chi_1) = (\chi_0 \vee \phi_- \vee \phi_-) \wedge (\chi_1 \vee \phi_-) \leftrightarrow (\chi_0 \wedge \chi_1) \vee \phi_- \vee (\phi_- \wedge \chi_1)$ . Since by the assumption of (iii),  $\phi_- \wedge \chi_1$  is not satisfiable, the latter formula is equivalent to  $(\chi_0 \wedge \chi_1) \vee \phi_- = \text{enl}(\chi_0 \wedge \chi_1)$ .

(iv) Assume  $i \in R$ , otherwise  $s_R$  does not change  $\chi_0, \chi_1, \chi_0 \wedge \chi_1$ .

Case 1:  $\phi^{012} \rightarrow \chi_s$  for  $s = 0, 1$ . Then  $\phi^{012} \rightarrow \chi_0 \wedge \chi_1$  and  $s_R(\chi_0) \wedge s_R(\chi_1) = \chi_0 \wedge \chi_1 = s_R(\chi_0 \wedge \chi_1)$ .

Case 2: E.g. not  $\phi^{012} \rightarrow \chi_0$ . Then not  $\phi^{012} \rightarrow \chi_0 \wedge \chi_1$  and  $s_R(\chi_0) \wedge s_R(\chi_1) = (\chi_0 \wedge x \cap z_i \neq y \cap z_i) \wedge s_R(\chi_1) \leftrightarrow (\chi_0 \wedge \chi_1) \wedge x \cap z_i \neq y \cap z_i = s_R(\chi_0 \wedge \chi_1)$ .

(v) For  $i \in R$ , the assumptions for (iii) are true for  $\psi_s = s_R(\chi_s)$ . Hence by (iii) and (iv),

$$\begin{aligned} (\text{enl}(s_R(\chi_0))(z_i, x \cap z_i, y \cap z_i) \wedge \text{enl}(s_R(\chi_1))(z_i, x \cap z_i, y \cap z_i)) \\ \rightarrow \text{enl}(s_R(\chi_0 \wedge \chi_1))(z_i, x \cap z_i, y \cap z_i). \end{aligned}$$

For  $i \notin R$ , we have  $\chi_s \rightarrow x \cap z_i = y \cap z_i$  for  $s = 0, 1$  and hence  $\text{enl}(s_R(\chi_0)) \wedge \text{enl}(s_R(\chi_1)) = (\chi_0 \vee \phi_-) \wedge (\chi_1 \vee \phi_-) \leftrightarrow (\chi_0 \wedge \chi_1) \vee \phi_- = \text{enl}(s_R(\chi_0 \wedge \chi_1))$ .

LEMMA 3.6. Let  $\phi$  be special and satisfiable,  $R = R(\phi)$ , and let  $\{\bigwedge_{i < r} \chi_{w,i} \mid w \in W\}$  be a set of representatives for  $\phi$ .

(i) For any  $\bigwedge_{i < r} \chi'_i \rightarrow \bigvee_{w \in W} \bigwedge_{i < r} s_R(\chi_{w,i})$ , there is a  $w \in W$  such that  $\bigwedge_{i < r} \chi'_i \rightarrow \bigwedge_{i < r} s_R(\chi_{w,i})$ .

(ii)  $\text{enl}(s(\phi)) \leftrightarrow \bigvee_{w \in W} \bigwedge_{i < r} \text{enl}(s_R(\chi_{w,i}))$ .

(iii) For any  $\bigwedge_{i < r} \chi'_i \rightarrow \bigvee_{w \in W} \bigwedge_{i < r} \text{enl}(s_R(\chi_{w,i}))$ , there is a  $w \in W$  such that  $\bigwedge_{i < r} \chi'_i \rightarrow \bigwedge_{i < r} \text{enl}(s_R(\chi_{w,i}))$ .

*Proof.* We will first prove (iii). Then the proof of (i) which is similar but easier will be clear. Let  $\bigwedge_{i < r} \chi'_i(z_i, x \cap z_i, y \cap z_i)$  be consistent, otherwise one can take any  $w \in W$ .

For  $i < r$  there is an  $n_i$ ,  $0 < n_i < 15$ , and there are  $\widehat{\chi}_{i,0}, \dots, \widehat{\chi}_{i,n_i-1} \in \{\phi^0, \dots, \phi^{14}\}$  such that

$$\bigwedge_{i < r} \chi'_i(z_i, x \cap z_i, y \cap z_i) \leftrightarrow \bigwedge_{i < r} (\widehat{\chi}_{i,0} \vee \dots \vee \widehat{\chi}_{i,n_i-1})(z_i, x \cap z_i, y \cap z_i).$$

We will show the claim by induction on  $\prod_{i < r} n_i$ .

*Case*  $\prod_{i < r} n_i = 1$ . Take an atomless Boolean algebra  $\mathfrak{A}$  and  $\bar{c} \in A$  such that  $\bar{c}$  is an enumeration of all the atoms in the generated subalgebra. Take  $a, b \in A$  such that  $\mathfrak{A} \models \bigwedge_{i < r} \chi'_i(c_i, a \cap c_i, b \cap c_i)$ . Then there is some  $w \in W$  with  $\mathfrak{A} \models \bigwedge_{i < r} \text{enl}(s_R(\chi_{w,i}(c_i, a \cap c_i, b \cap c_i)))$ . Since  $\bigwedge_{i < r} \chi'_i(z_i, x \cap z_i, y \cap z_i)$  defines an  $\mathcal{L}_{\omega\omega}$ -2-type of  $(x, y)$  over  $\bar{z}$ , we have  $\bigwedge_{i < r} \chi'_i(z_i, x \cap z_i, y \cap z_i) \rightarrow \bigwedge_{i < r} \text{enl}(s_R(\chi_{w,i}(z_i, x \cap z_i, y \cap z_i)))$ .

*Induction step.* We consider the step from  $\prod_{i < r} n_i$  to  $(n_0 + 1) \times \prod_{0 < i < r} n_i$ , the other cases are similar.

$$(\widehat{\chi}_{0,0} \vee \dots \vee \widehat{\chi}_{0,n_0}) \wedge \bigwedge_{0 < i < r} \chi'_i \leftrightarrow \left( \widehat{\chi}_{0,0} \wedge \bigwedge_{0 < i < r} \chi'_i \right) \vee \left( (\widehat{\chi}_{0,1} \vee \dots \vee \widehat{\chi}_{0,n_0}) \wedge \bigwedge_{0 < i < r} \chi'_i \right).$$

By induction hypothesis there are  $w', w'' \in W$  such that

$$\begin{aligned} \widehat{\chi}_{0,0} \wedge \bigwedge_{0 < i < r} \chi'_i &\rightarrow \bigwedge_{i < r} \text{enl}(s_R(\chi_{w',i})), \\ (\widehat{\chi}_{0,1} \vee \dots \vee \widehat{\chi}_{0,n_0}) \wedge \bigwedge_{0 < i < r} \chi'_i &\rightarrow \bigwedge_{i < r} \text{enl}(s_R(\chi_{w'',i})). \end{aligned}$$

Thus we have

$$\begin{aligned} &\left( \left( \widehat{\chi}_{0,0} \wedge \bigwedge_{0 < i < r} \chi'_i \right) \vee \left( (\widehat{\chi}_{0,1} \vee \dots \vee \widehat{\chi}_{0,n_0}) \wedge \bigwedge_{0 < i < r} \chi'_i \right) \right) \rightarrow \\ &\quad (\text{enl}(s_R(\chi_{w',0})) \vee \text{enl}(s_R(\chi_{w'',0}))) \wedge \bigwedge_{0 < i < r} (\text{enl}(s_R(\chi_{w',i})) \wedge \text{enl}(s_R(\chi_{w'',i}))). \end{aligned}$$

Note that in the last conjunction we get “and” and not only “or”, because

$$\bigwedge_{0 < i < r} \chi'_i \rightarrow \bigwedge_{0 < i < r} \text{enl}(s_R(\chi_{w',i})) \wedge \bigwedge_{0 < i < r} \text{enl}(s_R(\chi_{w'',i})),$$

as the situation below any  $z_i$  is independent of the situation below the other  $z_j$ .

From 3.5(i), (ii) and (v) we get

$$\begin{aligned} & \left( \widehat{\chi}_{0,0} \wedge \bigwedge_{0 < i < r} \chi'_i \right) \vee \left( (\widehat{\chi}_{0,1} \vee \dots \vee \widehat{\chi}_{0,n_0}) \wedge \bigwedge_{0 < i < r} \chi'_i \right) \\ & \quad \rightarrow \text{enl}(s_R(\chi_{w',0} \vee \chi_{w'',0})) \wedge \bigwedge_{0 < i < r} \text{enl}(s_R(\chi_{w',i} \wedge \chi_{w'',i})). \end{aligned}$$

Since  $\{\bigwedge_{i < r} \chi_{w,i}(z_i, x \cap z_i, y \cap z_i) \mid w \in W\}$  is a set of representatives for  $\phi(z, x, y)$  and since  $w', w'' \in W$ , we have  $(\chi_{w',0} \vee \chi_{w'',0}) \wedge \bigwedge_{0 < i < r} (\chi_{w',i} \wedge \chi_{w'',i}) \rightarrow \phi$  and there is a  $w \in W$  such that

$$(\chi_{w',0} \vee \chi_{w'',0}) \wedge \bigwedge_{0 < i < r} (\chi_{w',i} \wedge \chi_{w'',i}) \rightarrow \bigwedge_{i < r} \chi_{w,i}.$$

For such a  $w$  we have

$$\text{enl}(s_R(\chi_{w',0} \vee \chi_{w'',0})) \wedge \bigwedge_{0 < i < r} \text{enl}(s_R(\chi_{w',i} \wedge \chi_{w'',i})) \rightarrow \bigwedge_{i < r} \text{enl}(s_R(\chi_{w,i})),$$

and thus the induction step is complete and (iii) is shown.

(ii) Assume  $s(\phi)$  is satisfiable, otherwise both sides are not satisfiable. Let  $S = \{\bigwedge_{i < r} \chi_{w,i} \mid w \in W\}$  be a set of representatives for  $\phi$ , and  $S' = \{\bigwedge_{i < r} \chi'_{w',i} \mid w' \in W'\}$  be a set of representatives for  $s(\phi) = \bigvee_{w \in W} \bigwedge_{i < r} s_R(\chi_{w,i})$  such that  $W' \supseteq \widehat{W} := \{w \in W \mid \bigwedge_{i < r} s_R(\chi_{w,i}) \text{ is satisfiable}\}$  and  $\chi'_{w',i} = s_R(\chi_{w,i})$  for  $w \in \widehat{W}$ .

By definition,  $\text{enl}(s(\phi)) = \bigvee_{w' \in W'} \bigwedge_{i < r} \text{enl}(\chi'_{w',i})$ . By (i), for any  $w' \in W'$  there is some  $w \in W$  such that  $\bigwedge_{i < r} \chi'_{w',i} \rightarrow \bigwedge_{i < r} s_R(\chi_{w,i})$  and hence  $\bigwedge_{i < r} \text{enl}(\chi'_{w',i}) \rightarrow \bigwedge_{i < r} \text{enl}(s_R(\chi_{w,i}))$ . Thus  $\text{enl}(s(\phi)) \rightarrow \bigvee_{w \in W} \bigwedge_{i < r} \text{enl}(s_R(\chi_{w,i}))$ . The other direction follows immediately from the choice of  $S'$  and the definition of  $\text{enl}$ .

LEMMA 3.7. *Let  $\phi$  be a special formula. Then  $\text{enl}(\text{enl}(s(\phi))) \leftrightarrow \text{enl}(s(\phi))$ .*

Proof. Assume  $s(\phi)$  is satisfiable, otherwise both sides are not satisfiable. Let  $S, W$  be as above and  $S'' = \{\bigwedge_{i < r} \chi''_{w'',i} \mid w'' \in W''\}$  be a set of representatives for  $\text{enl}(s(\phi))$ . By definition,  $\text{enl}(\text{enl}(s(\phi))) = \bigvee_{w'' \in W''} \bigwedge_{i < r} \text{enl}(\chi''_{w'',i})$ . For  $w'' \in W''$  we have  $\bigwedge_{i < r} \chi''_{w'',i} \rightarrow \text{enl}(s(\phi))$ , hence by 3.6(ii),  $\bigwedge_{i < r} \chi''_{w'',i} \rightarrow \bigvee_{w \in W} \bigwedge_{i < r} \text{enl}(s_R(\chi_{w,i}))$ . By 3.6(iii) there is some  $w \in W$  such that  $\bigwedge_{i < r} \chi''_{w'',i} \rightarrow \bigwedge_{i < r} \text{enl}(s_R(\chi_{w,i}))$ , whence  $\bigwedge_{i < r} \text{enl}(\chi''_{w'',i}) \rightarrow \bigwedge_{i < r} \text{enl}(\text{enl}(s_R(\chi_{w,i})))$ . It is easy to check that for qf  $\chi(z_i, x \cap z_i, y \cap z_i)$  by definition

$$\text{enl}(\text{enl}(\chi(z_i, x \cap z_i, y \cap z_i))) \rightarrow \text{enl}(\chi(z_i, x \cap z_i, y \cap z_i)).$$

Therefore  $\bigwedge_{i < r} \text{enl}(\chi''_{w'',i}) \rightarrow \bigwedge_{i < r} \text{enl}(s_R(\chi_{w,i}))$ , and putting things together yields  $\bigvee_{w'' \in W''} \bigwedge_{i < r} \text{enl}(\chi''_{w'',i}) \rightarrow \bigvee_{w \in W} \bigwedge_{i < r} \text{enl}(s_R(\chi_{w,i}))$ , and, by 3.6(ii),  $\bigvee_{w'' \in W''} \bigwedge_{i < r} \text{enl}(\chi''_{w'',i}) \rightarrow \text{enl}(s(\phi))$ .

The other direction is obvious.

LEMMA 3.8.  $\neg \text{big}(s(\phi)) \rightarrow \neg \text{big}(\text{enl}(s(\phi)))$  is valid for special  $\phi$ .

Proof. Let  $\mathfrak{A}$  be any atomless Boolean algebra. Assume  $\mathfrak{A} \models \text{big}(\text{enl}(s(\phi(\overset{r}{z}, x, y))))(\overset{r}{c})$ . We show that  $\mathfrak{A} \models \text{big}(s(\phi(\overset{r}{z}, x, y))))(\overset{r}{c})$ . Since the 1-types of  $x$  and of  $y$  over  $\overset{r}{c}$  are determined by  $\mathfrak{A} \models \exists y \text{enl}(s(\phi(\overset{r}{c}, x, y)))$  and  $\mathfrak{A} \models \exists x \text{enl}(s(\phi(\overset{r}{c}, x, y)))$ , there is just one pair  $(I_2, I_3)$  such that

$$\begin{aligned} \mathfrak{A} \models & \bigvee_{\{(I_0, I_1) \mid I_0 \dot{\cup} I_1 \dot{\cup} I_2 \dot{\cup} I_3 = \{0, \dots, r-1\}, I_0 \neq \emptyset\}} \forall xy \\ & \left( \left( \bigwedge_{i \in I_0} \phi^{012}(c_i, x \cap c_i, y \cap c_i) \wedge \bigwedge_{i \in I_1} x \cap c_i = y \cap c_i \neq 0, c_i \right. \right. \\ & \left. \left. \wedge \bigwedge_{i \in I_2} x \cap c_i = y \cap c_i = 0 \wedge \bigwedge_{i \in I_3} x \cap c_i = y \cap c_i = c_i \right) \rightarrow \text{enl}(s(\phi(\overset{r}{c}, x, y))) \right). \end{aligned}$$

Take  $I_0 \subseteq$ -maximal such that

$$\begin{aligned} \mathfrak{A} \models & \forall xy \left( \left( \bigwedge_{i \in I_0} \phi^{012}(c_i, x \cap c_i, y \cap c_i) \wedge \bigwedge_{i \in I_1} x \cap c_i = y \cap c_i \neq 0, c_i \right. \right. \\ & \left. \left. \wedge \bigwedge_{i \in I_2} x \cap c_i = y \cap c_i = 0 \wedge \bigwedge_{i \in I_3} x \cap c_i = y \cap c_i = c_i \right) \rightarrow \text{enl}(s(\phi(\overset{r}{c}, x, y))) \right). \end{aligned}$$

Let  $R = R(\phi)$  and  $\{\chi_{w,i} \mid w \in W\}$  be a set of representatives for  $\phi$ . By 3.6(ii) and (iii) there is a  $w \in W$  such that

$$\begin{aligned} \mathfrak{A} \models & \forall xy \left( \left( \bigwedge_{i \in I_0} \phi^{012}(c_i, x \cap c_i, y \cap c_i) \wedge \bigwedge_{i \in I_1} x \cap c_i = y \cap c_i \neq 0, c_i \right. \right. \\ & \left. \left. \wedge \bigwedge_{i \in I_2} x \cap c_i = y \cap c_i = 0 \wedge \bigwedge_{i \in I_3} x \cap c_i = y \cap c_i = c_i \right) \right. \\ & \left. \rightarrow \bigwedge_{i < r} \text{enl}(s_R(\chi_{w,i}(c_i, x \cap c_i, y \cap c_i))) \right). \end{aligned}$$

We claim that also

$$\begin{aligned} \mathfrak{A} \models & \forall xy \left( \left( \bigwedge_{i \in I_0} \phi^{012}(c_i, x \cap c_i, y \cap c_i) \wedge \bigwedge_{i \in I_1} x \cap c_i = y \cap c_i \neq 0, c_i \right. \right. \\ & \left. \left. \wedge \bigwedge_{i \in I_2} x \cap c_i = y \cap c_i = 0 \wedge \bigwedge_{i \in I_3} x \cap c_i = y \cap c_i = c_i \right) \right. \\ & \left. \rightarrow \bigwedge_{i < r} s_R(\chi_{w,i}(c_i, x \cap c_i, y \cap c_i)) \right). \end{aligned}$$

Indeed, by the definition of  $\text{enl}$  we have for any  $s_R(\chi_{w,i}(z_i, x \cap z_i, y \cap z_i))$ : For  $i \in I_0$ , if  $\phi^{012} \rightarrow \text{enl}(s_R(\chi_{w,i}))$ , then  $\phi^{012} \rightarrow s_R(\chi_{w,i})$ . For  $i \in I_2$ , if  $x \cap z_i = y \cap z_i = 0 \rightarrow \text{enl}(s_R(\chi_{w,i}))$ , then  $x \cap z_i = y \cap z_i = 0 \rightarrow s_R(\chi_{w,i})$ .

For  $i \in I_3$ , if  $x \cap z_i = y \cap z_i = z_i \rightarrow \text{enl}(s_R(\chi_{w,i}))$ , then  $x \cap z_i = y \cap z_i = z_i \rightarrow s_R(\chi_{w,i})$ .

For  $i \in I_1$  the formula  $x \cap z_i = y \cap z_i \neq 0, z_i \wedge \text{enl}(s_R(\chi_{w,i})) \wedge \neg s_R(\chi_{w,i})$  is consistent only if  $\phi^{012} \rightarrow s_R(\chi_{w,i})$ . But then we could take  $I'_0 := I_0 \cup \{i\}$  and  $I'_1 = I_1 \setminus \{i\}$  and replace  $(I_0, I_1)$  by  $(I'_0, I'_1)$ , which contradicts the maximality of  $I_0$ .

Now we are ready to prove (\*\*) for special formulas of the form  $s(\phi)$ .

LEMMA 3.9. *Let  $\phi$  be special and  $\vec{c} \in B$  be an  $r$ -tuple that consists of atoms in the generated subalgebra.*

(i) *If  $\neg \text{big}(\phi)$  and  $\text{enl}(\phi) \rightarrow \phi$  are valid, then for any  $\alpha$  with  $\vec{c} \in B_\alpha$  the relation  $\phi(\vec{c}, x, y)$  is small in  $\mathfrak{B}_\alpha$ .*

(ii) *If  $\neg \text{big}(s(\phi))$  is valid, then for any  $\alpha$  with  $\vec{c} \in B_\alpha$  the relation  $\text{enl}(s(\phi(\vec{c}, x, y)))$  is small in  $\mathfrak{B}_\alpha$ .*

PROOF. (i) Let  $\mathfrak{B} \models \neg \text{big}(\phi(\vec{c}, x, y))(\vec{c})$  and  $\vec{c} \in B_\alpha$  be atoms in the generated subalgebra. Set  $\mathfrak{B}_\alpha =: \mathfrak{A}$ , and let  $M \neq \emptyset$  be a maximally homogeneous set for  $\phi(\vec{c}, x, y)$  in  $\mathfrak{A}$ , and  $(a, b)_A \in P(A)$ , i.e.  $(a, b)_A$  is an interval in  $\mathfrak{A}$ . Take  $(a', b')_A \leq (a, b)_A$  such that there is just one  $i \in r$ , say  $i_0$ , with  $(b' \setminus a') \subseteq c_i$  and  $c_i \cap a' \neq 0$  and  $b' \cap c_i \neq c_i$ . We assume  $\mathfrak{B}$  (and also  $\mathfrak{A}$  and  $\mathcal{P}(\omega)$ ) satisfy

$$\forall x \in (a', b') (\exists y \phi(\vec{c}, x, y) \wedge \exists y \phi(\vec{c}, y, x))(\vec{c}),$$

for otherwise  $(a', b')_A \in D_A(M, \phi(\vec{c}, x, y), 1, 0)$ .

Since  $\mathfrak{B} \models \neg \text{big}(\phi)(\vec{c})$ , we have  $(a', b')_A \cap M \neq (a', b')_A$ . We fix a  $d \in (a', b')_A \setminus M$  and an  $m \in M$  such that  $\mathfrak{A} \models \neg \phi(\vec{c}, d, m) \vee \neg \phi(\vec{c}, m, d)$ , say  $\mathfrak{A} \models \neg \phi(\vec{c}, d, m)$ , and show that there is an  $(a'', b'')_A \leq (a', b')_A$  such that for any  $x \in (a'', b'')_{\mathcal{P}(\omega)}$  we have  $x \in M$  or  $\mathcal{P}(\omega) \models \neg \phi(\vec{c}, x, m)$ .

Then (i) will be proved, because such an  $(a'', b'')_A$  is in  $D_A(M, \phi(\vec{c}, x, y), 1, 0)$ . Fix a set  $\{\bigwedge_{i < r} \chi_{w,i} \mid w \in W\}$  of representatives for  $\phi$ .

CLAIM.  $d \cap c_{i_0} \neq c_{i_0} \setminus m$ .

PROOF.  $\phi(\vec{c}, x, y) = \bigvee_{w \in W} \bigwedge_{i < r} \chi_{w,i}(z_i, x \cap z_i, y \cap z_i)$ , w.l.o.g.  $W = \{0, 1, \dots, s-1\}$ . Hence  $\mathfrak{A} \models \bigwedge_{w \in W} \bigvee_{i < r} \neg \chi_{w,i}(c_i, d \cap c_i, m \cap c_i)$ , say for  $w = 0, 1, \dots, s' - 1$

$$\mathfrak{A} \models \bigvee_{i < r, i \neq i_0} \neg \chi_{w,i}(c_i, d \cap c_i, m \cap c_i),$$

and for  $w = s', s' + 1, \dots, s - 1$

$$\mathfrak{A} \models \bigwedge_{i < r, i \neq i_0} \chi_{w,i}(c_i, d \cap c_i, m \cap c_i) \wedge \neg \chi_{w,i_0}(c_{i_0}, d \cap c_{i_0}, m \cap c_{i_0}).$$

We may assume  $s > 0$  and  $s' \leq s - 1$ , because otherwise  $(a', b')_A \in D_A(M, \phi(\bar{c}, x, y), 1, 0)$ . Since

$$\mathfrak{A} \models \forall xy \left( \left( \bigwedge_{s' \leq w < s} \bigwedge_{i < r, i \neq i_0} \chi_{w,i}(c_i, x \cap c_i, y \cap c_i) \right. \right. \\ \left. \left. \wedge \bigvee_{s' \leq w < s} \chi_{w,i_0}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0}) \right) \rightarrow \phi(\bar{c}, y, x) \right),$$

we have

$$\mathfrak{A} \models \forall xy \left( \left( \bigwedge_{s' \leq w < s} \bigwedge_{i < r, i \neq i_0} \chi_{w,i}(c_i, x \cap c_i, y \cap c_i) \right. \right. \\ \left. \wedge \left( \bigvee_{s' \leq w < s} \chi_{w,i_0}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0}) \vee (x \cap c_{i_0} = (-y) \cap c_{i_0}) \right) \right) \\ \left. \wedge \exists x \chi_{w,i_0}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0}) \wedge \exists y \chi_{w,i_0}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0}) \right) \\ \rightarrow \text{enl}(\phi(\bar{c}, x, y)).$$

By the assumptions on  $\phi(\bar{z}, x, y)$  and on  $\bar{c}$  there is just one 1-type of  $x \cap c_{i_0}$  over  $c_{i_0}$  consistent with  $\phi(\bar{c}, x, y)$  such that for every  $w \in W$  the formula  $\exists y \chi_{w,i_0}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0})$  is implied by this type. The same holds for the 1-type of  $y \cap c_{i_0}$  over  $c_{i_0}$ , which coincides with the 1-type of  $x \cap c_{i_0}$  over  $c_{i_0}$ , and the formula  $\exists x \chi_{w,i_0}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0})$ . Since  $m \cap c_{i_0}$  and  $d \cap c_{i_0}$  have this 1-type, we get

$$\mathfrak{A} \models \exists x \bigvee_{s' \leq w < s} \chi_{w,i_0}(c_{i_0}, x \cap c_{i_0}, m \cap c_{i_0}) \\ \wedge \exists y \bigvee_{s' \leq w < s} \chi_{w,i_0}(c_{i_0}, d \cap c_{i_0}, y \cap c_{i_0}).$$

Note that  $\mathfrak{A} \models \neg \phi(\bar{c}, d, m)$  and  $\phi$  is equivalent to  $\text{enl}(\phi)$ . Therefore  $d \cap c_{i_0} \neq c_{i_0} \setminus m$  and the claim is proved.

We now give  $(a'', b'')_A$  case by case.

Case 1:  $d \cap c_{i_0} \neq m \cap c_{i_0}$ . Then

$$\mathfrak{A} \models \bigvee_{i=0,1,2,4,8} \phi^i(c_{i_0}, d \cap c_{i_0}, m \cap c_{i_0}).$$

Assume that  $\mathfrak{A} \models \phi^i(c_{i_0}, d \cap c_{i_0}, m \cap c_{i_0})$ .

If  $i = 0$  or  $i = 2$ , take an  $e'$  such that  $0 \subset e' \subset c_{i_0} \cap m \cap (-d)$ , and  $(a'', b'')_A = (d, b' \setminus e')_A$ . If  $i = 1$  or  $i = 8$ , take  $(a'', b'')_A = (a', d)_A$ . Finally, if  $i = 4$ , take  $(a'', b'')_A = (d, b')_A$ .

Then, in each subcase, for any  $x \in (a'', b'')_{\mathcal{P}(\omega)}$  we have

$$\mathcal{P}(\omega) \models \text{tp}(x, m/\bar{c}) = \text{tp}(d, m/\bar{c}) \text{ and hence } \mathcal{P}(\omega) \models \neg\phi(\bar{c}, x, m).$$

Case 2:  $d \cap c_{i_0} = m \cap c_{i_0}$ .

Subcase 2.1:

$$\mathfrak{A} \models \exists xy \left( \phi^{012}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0}) \wedge \neg \bigvee_{s' \leq w < s} \chi_{w, i_0}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0}) \right).$$

Since  $\phi^{012}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0})$  determines the  $\mathcal{L}_{\omega\omega}$ -1-type of  $y \cap c_{i_0}$  over  $c_{i_0}$ , and  $m$  has the same one, we have

$$\mathfrak{A} \models \exists x \left( \phi^{012}(c_{i_0}, x \cap c_{i_0}, m \cap c_{i_0}) \wedge \neg \bigvee_{s' \leq w < s} \chi_{w, i_0}(c_{i_0}, x \cap c_{i_0}, m \cap c_{i_0}) \right).$$

There is an example  $d'$  for  $x$  with  $d' \cap c_{i_0} \in (a' \cap c_{i_0}, b' \cap c_{i_0})_A$ , because  $m \cap c_{i_0} = d \cap c_{i_0} \in (a' \cap c_{i_0}, b' \cap c_{i_0})_A$  and hence within the given 1-type of  $x \cap c_{i_0}$  over  $c_{i_0}$  the formula  $\phi^i(c_i, x \cap c_i, m \cap c_i)$  can be realized with some  $x \cap c_{i_0} \in (a' \cap c_{i_0}, b' \cap c_{i_0})_A$  for  $i = 0, 1, 2$ . We can argue with  $(d' \cap c_{i_0}) \cup (d \setminus c_{i_0})$  as with  $d$  in case 1 for  $i = 0, 1, 2$ .

Subcase 2.2:

$$\mathfrak{A} \models \forall xy \left( \phi^{012}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0}) \rightarrow \bigvee_{s' \leq w < s} \chi_{w, i_0}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0}) \right).$$

Again we have

$$\begin{aligned} \mathfrak{A} \models \forall xy \left( \left( \bigwedge_{s' \leq w < s} \bigwedge_{i < r, i \neq i_0} \chi_{w, i}(c_i, x \cap c_i, y \cap c_i) \right. \right. \\ \left. \left. \wedge \bigvee_{s' \leq w < s} \chi_{w, i}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0}) \right) \rightarrow \phi(\bar{c}, x, y) \right). \end{aligned}$$

Since

$$\phi^{012}(z_{i_0}, x \cap z_{i_0}, y \cap z_{i_0}) \rightarrow \bigvee_{s' \leq w < s} \chi_{w, i}(z_{i_0}, x \cap z_{i_0}, y \cap z_{i_0}),$$

by the definition of  $\text{enl}$  we have

$$\begin{aligned} \forall xy \left( \left( \bigwedge_{i < r, i \neq i_0} \text{enl} \left( \bigwedge_{s' \leq w < s} \chi_{w, i}(z_i, x \cap z_i, y \cap z_i) \right) \right. \right. \\ \wedge \left( \bigvee_{s' \leq w < s} \chi_{w, i}(z_{i_0}, x \cap z_{i_0}, y \cap z_{i_0}) \vee (x \cap z_{i_0} = y \cap z_{i_0}) \right) \\ \wedge \exists x \bigvee_{s' \leq w < s} \chi_{w, i_0}(z_{i_0}, x \cap z_{i_0}, y \cap z_{i_0}) \\ \left. \left. \wedge \exists y \bigvee_{s' \leq w < s} \chi_{w, i_0}(z_{i_0}, x \cap z_{i_0}, y \cap z_{i_0}) \right) \right) \rightarrow \text{enl}(\phi(\bar{z}, x, y)). \end{aligned}$$

In  $\mathfrak{A}$  we get

$$\begin{aligned} \mathfrak{A} \models \forall xy \left( \left( \bigwedge_{s' \leq w < s} \bigwedge_{i < r, i \neq i_0} \text{enl}(\chi_{w,i}(c_i, x \cap c_i, y \cap c_i)) \right. \right. \\ \wedge \left( \bigvee_{s' \leq w < s} \chi_{w,i_0}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0}) \vee (x \cap c_{i_0} = y \cap c_{i_0}) \right. \\ \wedge \exists x \bigvee_{s' \leq w < s} \chi_{w,i_0}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0}) \\ \left. \left. \left. \wedge \exists y \bigvee_{s' \leq w < s} \chi_{w,i_0}(c_{i_0}, x \cap c_{i_0}, y \cap c_{i_0}) \right) \right) \right) \rightarrow \text{enl}(\phi(\vec{c}, x, y)). \end{aligned}$$

As in the first subcase, we get

$$\begin{aligned} \mathfrak{A} \models \exists x \bigvee_{s' \leq w < s} \chi_{w,i_0}(c_{i_0}, x \cap c_{i_0}, m \cap c_{i_0}) \\ \wedge \exists y \bigvee_{s' \leq w < s} \chi_{w,i_0}(c_{i_0}, d \cap c_{i_0}, y \cap c_{i_0}) \wedge d \cap c_{i_0} = m \cap c_{i_0}. \end{aligned}$$

Putting things together yields  $\mathfrak{A} \models \text{enl}(\phi(\vec{c}, d, m))$  and hence  $\mathfrak{A} \models \phi(\vec{c}, d, m)$ , a contradiction to the choice of  $d$  and  $m$ .

(ii) By 3.8,  $\neg \text{big}(s(\phi)) \rightarrow \neg \text{big}(\text{enl}(s(\phi)))$ , and, by 3.7,  $\text{enl}(\text{enl}(s(\phi))) \rightarrow \text{enl}(s(\phi))$  is valid. Therefore (ii) follows from (i) applied to  $\text{enl}(s(\phi))$ .

Lemma 3.9, the construction and the monotonicity of  $Q_1^2$  yield:

**THEOREM 3.10.** *For any special  $\phi$ ,*

$$\begin{aligned} \mathfrak{B} \models \forall \vec{z} \left( \left( \text{“}\vec{z} \text{ are the atoms in the generated subalgebra”} \wedge \neg \text{big}(s(\phi))(\vec{z}) \right) \right. \\ \left. \rightarrow \neg Q_1^2 xy s(\phi(\vec{z}, x, y)) \right). \end{aligned}$$

Finally, we show how to get Theorem 3.10 for  $\phi$  instead of  $s(\phi)$ .

**THEOREM 3.11.** *For any special  $\phi$*

$$\begin{aligned} \mathfrak{B} \models \forall \vec{z} \left( \left( \text{“}\vec{z} \text{ are the atoms in the generated subalgebra”} \wedge \neg \text{big}(\phi)(\vec{z}) \right) \right. \\ \left. \rightarrow \neg Q_1^2 xy \phi(\vec{z}, x, y) \right). \end{aligned}$$

**Proof** (by induction on  $\text{card}(R(\phi))$ ). If  $R(\phi) = \emptyset$ , then  $\phi(\vec{z}, x, y) \rightarrow x = y$ , and hence  $\mathfrak{B} \models \neg Q_1^2 xy \phi(\vec{z}, x, y)$ .

Now assume  $\mathfrak{B} \models \forall \vec{z} \left( \left( \text{“}\vec{z} \text{ are the atoms in the generated subalgebra”} \wedge \neg \text{big}(\psi)(\vec{z}) \right) \rightarrow \neg Q_1^2 xy \psi(\vec{z}, x, y) \right)$  for all  $\psi$  with  $R(\psi) \subset R(\phi)$ . We show  $\mathfrak{B} \models Q_1^2 xy \phi(\vec{c}, x, y) \rightarrow \text{big}(\phi)(\vec{c})$  for any  $r$ -tuple  $\vec{c}$  that consists of atoms in the generated subalgebra. Assume  $\mathfrak{B} \models Q_1^2 xy \phi(\vec{c}, x, y)$  and let  $H$  be an uncountable homogeneous set for  $\phi(\vec{c}, x, y)$  in  $\mathfrak{B}$ . By recursion on  $i \leq r$  we define uncountable subsets  $H^{(i)}$ ,  $0 \leq i \leq r$ .

Set  $H^{(0)} := H$ . Assume  $H^{(i)}$  is defined. We distinguish two cases:

Case 1:  $\{x \cap c_i \mid x \in H^{(i)}\}$  is uncountable. Then take  $H^{(i+1)} \subseteq H^{(i)}$  such that  $H^{(i+1)}$  is uncountable and for any  $x, y \in H^{(i+1)}$ , if  $x \neq y$  then  $x \cap c_i \neq y \cap c_i$ .

Case 2:  $\{x \cap c_i \mid x \in H^{(i)}\}$  is countable. Then there is some  $x \in H^{(i)}$  such that  $\{y \in H^{(i)} \mid x \cap c_i = y \cap c_i\}$  is uncountable. Let  $H^{(i+1)}$  be such a set.

For  $i \notin R$ ,  $\{x \cap c_i \mid x \in H^{(i)}\}$  is a singleton, and we are in case 2. Now consider  $H^{(0)}, H^{(1)}, \dots, H^{(r)}$ . If for all  $i \in R$  case 1 is true, then  $H^{(r)}$  shows  $\mathfrak{B} \models Q_1^2 xy s(\phi(\overset{r}{c}, x, y))$ . By 3.10,  $\mathfrak{B} \models \text{big}(s(\phi(\overset{r}{c})))$ . Since  $s(\phi) \rightarrow \phi$ ,  $\mathfrak{B} \models \text{big}(\phi(\overset{r}{c}))$ .

If there is some  $i \in R$  with case 2 being true, fix such an  $i$ . Then  $H^{(i+1)}$  shows  $\mathfrak{B} \models Q_1^2 xy (\phi \wedge x \cap z_i = y \cap z_i)(\overset{r}{c}, x, y)$ . Take  $\psi = \phi \wedge x \cap z_i = y \cap z_i$ . Then  $\psi$  is also special. Since  $\psi \rightarrow \phi$  and  $i \in R(\phi) \setminus R(\psi)$ , we have  $R(\psi) \subset R(\phi)$ . By induction hypothesis, we conclude from  $\mathfrak{B} \models Q_1^2 xy (\phi \wedge x \cap z_i = y \cap z_i)(\overset{r}{c}, x, y)$  that  $\mathfrak{B} \models \text{big}(\psi(\overset{r}{c}))$  and hence  $\mathfrak{B} \models \text{big}(\phi(\overset{r}{c}))$ .

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