## Arithmetic progressions in sums of subsets of sparse sets

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1. Introduction. Let $A=\left\{a_{1}, \ldots, a_{k}\right\}$ be a subset of integers. Denote by

$$
\mathcal{S}(A)=\left\{\sum_{i=1}^{k} \varepsilon_{i} a_{i}: \varepsilon_{i} \in\{0,1\}\right\}
$$

the subsets sumset of $A$ and let $\mathrm{L}(S)$ stand for the length of the longest arithmetic progression in $S$. The problem of finding large arithmetic structures in $\mathcal{S}(A)$ (or generally in sumsets) is one of the most fundamental in combinatorial number theory. It has been intensively studied, especially in the case of sufficiently dense sets $A \subseteq\{1, \ldots, n\}\left(|A| \geq n^{\alpha}\right.$ for some $\alpha>1 / 2$; see for example [1], 3], [4, [5]). A complete solution of this problem for sets of polynomial size was given recently by Szemerédi and Vu in [6], [7] and [8]. They proved, among other things, that if $A \subseteq[n]$ and $|A| \gg d_{d} n^{1 / d}$, where $d \geq 2$ is a fixed integer, then

$$
\begin{equation*}
\mathrm{L}(\mathcal{S}(A)) \ggg{ }_{d}|A|^{1+1 /(d-1)} . \tag{1.1}
\end{equation*}
$$

However, much less is known in the case of sparse sets, with $|A|=n^{o(1)}$. The only paper dealing with this question is [2], where the authors showed that for every set $A \subseteq\{1, \ldots, n\}$ we have

$$
\mathrm{L}(\mathcal{S}(A)) \gg|A| / \log ^{2} n
$$

Actually, they proved that one can find an arithmetic progression of the form $a, 2 a, \ldots, \mathrm{~L} a$ in $\mathcal{S}(A)$.

The aim of this note is to provide new estimates on $\mathrm{L}(\mathcal{S}(A))$ for sets $A \subseteq\{1, \ldots, n\}$ of size $n^{o(1)}$. First, we improve the bound of Erdős and Sárközy by a $\log n$ factor. Then, we establish a bound for sets with at least $\exp \left(\log ^{1 / 2+o(1)} n\right)$ elements, which is very similar to 1.1). Finally, we also provide some examples to show that our estimates are close to best possible.

[^0]We will use the following notation. For subsets of integers $A, B$ we put $A+B=\{a+b: a \in A, b \in B\}$. By $\log x$ we always mean $\log _{e} x$, where $e=2.71 \ldots$ is the Euler number. Furthermore, throughout the paper, we assume that $n$ is large enough if necessary.
2. Sparse sets. By a $d$-cube we mean any set of the form

$$
\mathcal{C}=\mathcal{C}\left(x ; x_{1}, \ldots, x_{d}\right)=\left\{x+\sum_{i=1}^{d} \varepsilon_{i} x_{i}: \varepsilon_{i} \in\{0,1\}\right\}
$$

where $x, x_{1}, \ldots, x_{d} \in \mathbb{Z}$. Observe that $\mathcal{C}=x+\left\{0, x_{1}\right\}+\cdots+\left\{0, x_{d}\right\}$. Our approach is based on the following elementary lemma.

Lemma 2.1. Let $A, B$ be finite sets of integers such that $|A+B| \leq K|A|$. If $B$ contains a d-cube, then

$$
\mathrm{L}(A+B) \geq\left\lceil\left(K^{1 / d}-1\right)^{-1}\right\rceil+1 \geq \frac{d}{(e-1) \log K}+1
$$

Proof. Suppose that $\mathcal{C}\left(x ; x_{1}, \ldots, x_{d}\right) \subseteq B$ and put $T_{0}=A+x, T_{i}=$ $T_{i-1}+\left\{0, x_{i}\right\}$ for $i=1, \ldots, d$. We have

$$
T_{0} \subseteq T_{1} \subseteq \cdots \subseteq T_{d} \subseteq A+B
$$

so that

$$
|A|=\left|T_{0}\right| \leq\left|T_{1}\right| \leq \cdots \leq\left|T_{d}\right| \leq K|A|
$$

Hence, for some $0 \leq j \leq d-1,\left|T_{j+1}\right|=\left|T_{j}+\left\{0, x_{j+1}\right\}\right| \leq K^{1 / d}\left|T_{j}\right|$, so that

$$
\left|\left(T_{j}+x_{j+1}\right) \cap T_{j}\right| \geq\left(2-K^{1 / d}\right)\left|T_{j}\right|
$$

It is easy to see that if $|(S+x) \cap S| \geq(1-\delta)|S|$, then there is an arithmetic progression in $S$ of length $\lceil 1 / \delta\rceil$ and common difference $x$. Thus, $\mathrm{L}\left(T_{j}\right) \geq$ $\left\lceil\left(K^{1 / d}-1\right)^{-1}\right\rceil$ and

$$
\mathrm{L}(A+B) \geq \mathrm{L}\left(T_{j+1}\right) \geq\left\lceil\left(K^{1 / d}-1\right)^{-1}\right\rceil+1
$$

If $d \geq \log K$, then

$$
K^{1 / d} \leq 1+\frac{\log K}{d} \sum_{i \geq 1} \frac{1}{i!}=1+(e-1) \frac{\log K}{d}
$$

and the assertion follows.
Now we can improve the bound of Erdős and Sárközy.
Corollary 2.2. If $A \subseteq\{1, \ldots, n\}$, then

$$
\mathrm{L}(\mathcal{S}(A)) \geq \frac{|A|}{4(e-1) \log n}
$$

Proof. Observe that $\mathcal{S}(A) \subseteq\left[n^{2}\right]$ and $\mathcal{C}\left(a_{1} ; a_{2}, \ldots, a_{k}\right) \subseteq \mathcal{S}(A)$, where $A=\left\{a_{1}, \ldots, a_{k}\right\}$. Put

$$
\mathcal{C}_{1}=\{0\} \quad \text { and } \quad \mathcal{C}_{2}=\mathcal{C}\left(a_{1} ; a_{2}, \ldots, a_{k}\right)
$$

Clearly, $\mathcal{C}_{1}+\mathcal{C}_{2} \subseteq \mathcal{S}(A)$, so that

$$
\left|\mathcal{C}_{1}+\mathcal{C}_{2}\right|=\left|\mathcal{C}_{2}\right|\left|\mathcal{C}_{1}\right| \leq|\mathcal{S}(A)|\left|\mathcal{C}_{1}\right| \leq n^{2}\left|\mathcal{C}_{1}\right|
$$

By Lemma 2.1 applied with $A=\mathcal{C}_{1}, B=\mathcal{C}_{2}$ and $K=n^{2}$ we have

$$
\mathrm{L}(\mathcal{S}(A)) \geq \mathrm{L}(\mathcal{C}) \geq \frac{d\left(\mathcal{C}_{2}\right)}{2(e-1) \log n^{2}}+1 \geq \frac{|A|}{4(e-1) \log n}
$$

The set of powers of 3 contained in $\{1, \ldots, n\}$ shows that in general one cannot improve the bound given by Corollary 2.2. However, in the next section, we show that this is possible for sufficiently dense sets.
3. Sets of size at least $e^{(\log n)^{1 / 2+o(1)}}$. We start with a version of Lemma 2.1 more suitable for our purpose.

Lemma 3.1. Suppose that $A \subseteq\{1, \ldots, n\}$ and $|A| \geq 100 \log n$. Then there exists a subset $A_{1}$ of $A$ such that $\left|A_{1}\right| \geq|A| / 2$ and for any $x_{1}, \ldots, x_{h} \in A_{1}$ and any integers $i_{1}, \ldots, i_{h}$ there is an arithmetic progression in $\mathcal{S}\left(A \backslash A_{1}\right)$ of length $|A|(10 \log n)^{-1}\left(\sum_{j}\left|i_{j}\right|\right)^{-1}$ and common difference $\sum_{j} i_{j} x_{j}$.

Proof. We define a set $B=\left\{a_{1}, \ldots, a_{l}\right\} \subseteq A$ recursively in the following way. Let $a_{1} \in A$ be arbitrary. Suppose that $B_{i}=\left\{a_{1}, \ldots, a_{i}\right\}$ has already been chosen. If there is $a_{i+1} \in A \backslash\left\{a_{1}, \ldots, a_{i}\right\}$ such that

$$
\begin{equation*}
\left|\mathcal{S}\left(B_{i} \cup\left\{a_{i+1}\right\}\right)\right| \geq\left(1+5|A|^{-1} \log n\right)\left|\mathcal{S}\left(B_{i}\right)\right| \tag{3.1}
\end{equation*}
$$

then we define $B_{i+1}=\left\{a_{1}, \ldots, a_{i+1}\right\}$. If such an element does not exist, we stop the algorithm putting $B=B_{i}$. We have to prove that this procedure terminates after at most $|A| / 2$ steps. Clearly, for each $i$,

$$
\left|\mathcal{S}\left(B_{i}\right)\right| \geq\left(1+5|A|^{-1} \log n\right)^{i-1}
$$

Thus, using the inequality $(1+1 / t)^{t+1 / 2}>e$ for $t>0$, we see that

$$
\left|\mathcal{S}\left(B_{i}\right)\right| \geq e^{\frac{5(i-1) \log n}{|A|+2 \log n}}
$$

On the other hand $\left|\mathcal{S}\left(B_{i}\right)\right| \leq|\mathcal{S}(A)| \leq n^{2}$, so that $l \leq|A| / 2$. We show that our assertion holds for $A_{1}=A \backslash B$. Suppose that $x_{1}, \ldots, x_{h} \in A_{1}$ and $i_{1}, \ldots, i_{h} \in \mathbb{Z}$. By (3.1), for every $1 \leq j \leq h$ we have

$$
\left|\mathcal{S}(B) \cap\left(\mathcal{S}(B)+x_{j}\right)\right| \geq\left(1-10|A|^{-1} \log n\right)|\mathcal{S}(B)|
$$

hence

$$
\left|\mathcal{S}(B) \cap\left(\mathcal{S}(B)+\sum_{j} i_{j} x_{j}\right)\right| \geq\left(1-10|A|^{-1} \sum_{j}\left|i_{j}\right| \log n\right)|\mathcal{S}(B)|
$$

so $\mathcal{S}(B)$ contains an arithmetic progression of length $|A|\left(8 \sum_{j}\left|i_{j}\right| \log n\right)^{-1}$ and common difference $\sum_{j} i_{j} x_{j}$.

Now we are ready to prove the main result of this section.
Theorem 3.2. Suppose that $A \subseteq\{1, \ldots, n\}$ and $|A| \geq 8(n / \log n)^{1 / d}$, where $0<d \leq \sqrt{\frac{\log n}{2 \log \log n}}$ is an integer. Then

$$
\mathrm{L}(\mathcal{S}(A)) \geq 2^{-10}(|A| / \log n)^{1+1 / d}
$$

Proof. Let $A_{1}$ and $A_{2}$ be the sets given by Lemma 3.1 applied for $A$ and $A_{1}$, respectively. We have $\left|A_{1}\right| \geq|A| / 2$ and $\left|A_{2}\right| \geq|A| / 4$.

Put $t=\left\lfloor(|A| / \log n)^{1 / d}\right\rfloor$. By hypothesis it follows that $t \geq 2$. We show that the equation

$$
\begin{equation*}
x_{0}+t x_{1}+\cdots+t^{d} x_{d}=y_{0}+t y_{1}+\cdots+t^{d} y_{d} \tag{3.2}
\end{equation*}
$$

can be nontrivially (this means $x_{i} \neq y_{i}$ for some $i$ ) solved in $A_{2}$. Indeed, all sums $x_{1}+t x_{2}+\cdots+t^{d} x_{d+1}, x_{i} \in A_{2}$, are less than $2 t^{d} n$. Since $\left|A_{2}\right|^{d+1} \geq 2 t^{d} n$ two among them are equal, giving a nontrivial solution $x_{0}, \ldots, x_{d}, y_{0}, \ldots, y_{d} \in A_{2}$. Denote by $i$ the largest index $0 \leq j \leq d$ such that $x_{j} \neq y_{j}$. Then

$$
t^{i}\left(x_{i}-y_{i}\right)=\left(y_{1}-x_{1}\right)+\cdots+t^{i-1}\left(y_{i-1}-x_{i-1}\right) \neq 0
$$

Put $b=x_{i}-y_{i}$ and $c=\left(y_{0}-x_{0}\right)+\cdots+t^{i-1}\left(y_{i-1}-x_{i-1}\right)$. By Lemma 3.1, $\mathcal{S}\left(A \backslash A_{1}\right)$ contains an arithmetic progression $P_{1}$ of length $\left|A_{1}\right| /(20 \log n) \geq$ $|A| /(40 \log n)$ and common difference $b$, and $\mathcal{S}\left(A_{1} \backslash A_{2}\right)$ contains an arithmetic progression $P_{2}$ of length

$$
\left|A_{2}\right| /\left(10\left(1+t+\cdots+t^{i-1}\right) \log n\right) \geq|A| /\left(80 t^{i-1} \log n\right)
$$

and common difference $c$. As $t b=c, P_{1}+P_{2}$ is an arithmetic progression of length at least $2^{-9} t|A| / \log n \geq 2^{-10}(|A| / \log n)^{1+1 / d}$ and common difference $b$. The assertion follows from the inclusion

$$
P_{1}+P_{2} \subseteq \mathcal{S}\left(A \backslash A_{1}\right)+\mathcal{S}\left(A_{1} \backslash A_{2}\right) \subseteq \mathcal{S}(A)
$$

Using the same argument one can prove an analogous result for $\mathbb{Z} / p \mathbb{Z}$. An improvement relies on the fact that the equation (3.2) can be solved in $\mathbb{Z} / p \mathbb{Z}$ in sparser subsets.

Theorem 3.3. Let $A \subseteq \mathbb{Z} / p \mathbb{Z}$ and $|A| \geq 8(p / \log p)^{1 / d}$, where $0<d \leq$ $\sqrt{\frac{\log p}{2 \log \log p}}$ is an integer. Then

$$
\mathrm{L}(\mathcal{S}(A)) \geq 2^{-10}(|A| / \log p)^{1+1 /(d-1)}
$$

4. A construction of sets with small $\mathrm{L}(\mathcal{S}(A))$. We prove here that bounds given by Corollary 2.2 and Theorem 3.2 are near optimal. First, we observe that an example of sets given by Szemerédi and Vu can be applied
for all $d \ll \sqrt{\log n}$. We recall briefly their construction. Let $n, d, m \in \mathbb{N}$ with $2 m^{d} \leq n^{1 / d}$ and set

$$
P=\left\{\sum_{i=0}^{d-1} x_{i}\left(2 m^{d+1}\right)^{i}: 1 \leq x_{i} \leq m\right\}
$$

Thus,

$$
\mathcal{S}(P) \subseteq\left\{\sum_{i=0}^{d-1} x_{i}\left(2 m^{d+1}\right)^{i}: 1 \leq x_{i} \leq m^{d+1}\right\}
$$

Clearly $P \subseteq\{1, \ldots, n\}$ is a $d$-dimensional, proper generalized arithmetic progression, hence $|P|=m^{d}$ and $\mathcal{S}(P)$ is 2-Freiman isomorphic (see [9]) to a subset of $\left[m^{d+1}\right]^{d}$, so that $\mathrm{L}(\mathcal{S}(P)) \leq m^{d+1}=|P|^{1+1 / d}$. In particular, for every $k$ there is a subset of $\left\{1, \ldots, 2^{k^{2}}\right\}$ of size $2^{k-1}$ containing an arithmetic progression of length at most $2^{k}$ in its subsets sumset. Now, we adapt this construction to sparser sets.

Theorem 4.1. For every positive integer $n$ and $100 \log n \leq t \leq 2^{\sqrt{\log _{2} n}-1}$ there is a set $A \subseteq\{1, \ldots, n\}$ such that $|A|=t$ and

$$
\mathrm{L}(\mathcal{S}(A)) \leq 10 \frac{|A|}{\log _{2} n}\left(\log _{2} \frac{|A|}{\log _{2} n}\right)^{2}
$$

Proof. Put

$$
K=\frac{10 t}{\log _{2} n}\left(\log _{2} \frac{t}{\log _{2} n}\right)^{2} \quad \text { and } \quad k=\lfloor\log K\rfloor .
$$

By the construction above, there is a set $X \subseteq\left\{1, \ldots, 2^{k^{2}}\right\}$ such that $|X|=$ $2^{k-1}$ and $\mathrm{L}(\mathcal{S}(X)) \leq 2^{k}$. Let $A$ consist of all numbers of the form $x 2^{i\left(k^{2}+k\right)}$, where $x \in X$ and $0 \leq i \leq l=\left\lfloor\log _{2} n /\left(k^{2}+k\right)-1\right\rfloor$. Clearly, $A \subseteq\{1, \ldots, n\}$ and

$$
|A|=(l+1) 2^{k-1} \geq \frac{\log _{2} n}{4 \log _{2}^{2} K} \frac{10 t}{\log _{2} n}\left(\log _{2} \frac{t}{\log _{2} n}\right)^{2} \geq t
$$

Observe that

$$
\mathcal{S}(A)=\left\{\sum_{i=0}^{l} x 2^{i\left(k^{2}+k\right)}: x \in \mathcal{S}(X)\right\}
$$

Therefore, from $\sum_{x \in X} x<2^{k^{2}+k-1}$, we deduce that $\mathcal{S}(A)$ is 2-Freiman isomorphic to $\mathcal{S}(X) \times \cdots \times \mathcal{S}(X)(l+1$ times $)$, hence

$$
\mathrm{L}(\mathcal{S}(A))=\mathrm{L}(\mathcal{S}(X)) \leq 2^{k} \leq K=\frac{10 t}{\log _{2} n}\left(\log _{2} \frac{t}{\log _{2} n}\right)^{2}
$$

and the assertion follows.
5. Concluding remarks. The bound (1.1) of Szemerédi and Vu and the examples above support the claim that our estimates are not optimal. One would like to replace $d$ by $d-1$ and remove the $\log n$ factor in Theorem 3.2. It also seems that Corollary 2.2 is not best possible for sets of size $\omega(n) \log n$ if $\omega(n) \rightarrow \infty$ as $n \rightarrow \infty$. Essentially there are two places where our argument could be refined. More specifically, the results can be strengthen provided we can solve any of the following three problems.

Problem 5.1. Is it true that for every $A \subseteq[n]$, there is a subset $A^{\prime} \subseteq A$ of size (roughly) $|A| / 2$ such that

$$
\left|\mathcal{S}\left(A^{\prime}\right)\right| \gg|\mathcal{S}(A)|^{1-\varepsilon}
$$

where $\varepsilon \rightarrow 0$ as $|A| / \log n \rightarrow \infty$ ?
If we can answer the above question in the affirmative, then by Lemma 2.1 applied with $\mathcal{S}\left(A^{\prime}\right), \mathcal{S}\left(A \backslash A^{\prime}\right)$ and $K=|\mathcal{S}(A)|^{\varepsilon}$ we have

$$
\left|\mathcal{S}\left(A^{\prime}\right)+\mathcal{S}\left(A \backslash A^{\prime}\right)\right| \leq|\mathcal{S}(A)| \ll K\left|\mathcal{S}\left(A^{\prime}\right)\right|
$$

so

$$
\mathrm{L}\left(\mathcal{S}\left(A^{\prime}\right)\right) \gg|A| /(\varepsilon \log n)
$$

which improves Corollary 2.2 . It would even be sufficient for us if the following weaker question had a positive answer.

Problem 5.2. Is it true that for every set $A$ there is a subset $B \subseteq A$ with $|B| \geq|A| / 2$ having the property described in Problem 5.1?

Probably one can also replace $d$ by $d-1$. To do it one has to find a more efficient argument than the one we used, based on solutions to a linear equation. This could be done if the following question had a positive answer.

Problem 5.3. Is it true that there exists an absolute constant $C>0$ such that for every set $A \subseteq[n]$ with $|A|>C n^{1 / d}$ (we may allow $C$ to depend at most exponentially on d) there are $a_{1}, \ldots, a_{u}, b_{1}, \ldots, b_{v} \in \pm A$ and $M \in \mathbb{Z}$ such that

$$
M \sum_{i=1}^{u} a_{i}=\sum_{j=1}^{v} b_{j} \neq 0
$$

and $|A| / u>M>|A|^{1 /(d-1)} v$ ?
Acknowledgements. This research was partly supported by MNSW grant N N201 543538.

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Received on 18.3.2010
and in revised form on 29.6.2010


[^0]:    2010 Mathematics Subject Classification: Primary 11B25; Secondary 11P99. Key words and phrases: arithmetic progressions, sumsets.

