

## Solving parametrized systems of Pell equations

by

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*To my long time friend Pierre Samuel,  
octogenarian in love with Diophantus*

**1. Introduction.** Let  $F > 1$  be a square-free integer. In his papers [3]–[6] Ljunggren studied the quartic equation

$$x^4 - Fy^2 = 1$$

(and similar equations) and proved that the equation has at most two solutions  $(x, y)$  in positive integers. He also gave an algorithm to find the solutions.

When  $F = p$  is a prime number, Ljunggren showed that

$$x^4 - py^2 = 1$$

has no solution in positive integers when  $p \neq 5, 29$ . Moreover if  $p = 5$  the only solution is  $(3, 4)$  and if  $p = 29$  the only solution is  $(99, 1820)$ . In [12] Samuel gave another proof for  $p = 5$ .

The proof of this result leads to the study of the systems

$$\begin{cases} x^2 - 2y^2 = \pm 1, \\ x^2 - 2py^2 = \mp 1. \end{cases}$$

In our paper we shall need binary recurring sequences. Let  $P > 0$ ,  $Q \neq 0$  be integers such that  $D = P^2 - 4Q > 0$ . We shall consider the Lucas sequences  $(U_n)_n$ ,  $(V_n)_n$  with parameters  $(P, Q)$ :

$$\begin{aligned} U_0 &= 0, & U_1 &= 1, & V_0 &= 2, & V_1 &= P, \\ U_n &= PU_{n-1} - QU_{n-2} \text{ for } n \geq 2, & V_n &= PV_{n-1} - QV_{n-2} \text{ for } n \geq 2, \\ U_n &= (-1/Q^n)U_{-n} \text{ for } n < 0; & V_n &= (1/Q^n)V_{-n} \text{ for } n < 0. \end{aligned}$$

As is easily seen, the above recurrences still hold for any integer  $n$ . When needed, we shall use the notation  $U_n = U_n(P, Q)$ ,  $V_n = V_n(P, Q)$ .

Given the square-free integer  $F > 1$ , let  $\varepsilon = c + d\sqrt{F}$  be the fundamental unit of the ring  $\mathbb{Z}[\sqrt{F}]$ , so  $c > 0$ ,  $d > 0$ . As is well known,  $c$  and  $d$  are effectively bounded in terms of  $F$ . Let  $P = 2c$ ,  $Q = c^2 - Dd^2 = \pm 1$ , let  $U_n = U_n(2c, Q)$ ,  $V_n = V_n(2c, Q)$ . Then

$$\varepsilon^n = \frac{V_n}{2} + dU_n\sqrt{F}.$$

As is easily seen,  $V_n$  is even for every  $n \in \mathbb{Z}_-$  and  $V_{2n} \equiv 2 \pmod{4}$  for every  $n \in \mathbb{Z}$ . If  $s \geq 1$  we define

$$k_s = \frac{1}{4}(2 + Q^s V_{2s}), \quad h_s = \frac{1}{4}(2 - Q^s V_{2s}),$$

so  $h_s, k_s \neq 0, 1$  and  $k_s + h_s = 1$ .

(1.1) THEOREM. *Let  $F > 1$ ,  $G \geq 1$  be square-free integers, let  $s \geq 1$ ,  $f \neq 0$  and  $g = fk_s$  or  $g = fh_s$ . Then there exists an integer  $N > 0$ , effectively computable in terms of  $F, G, f$  and  $s$ , such that if  $p = 1$  or  $p$  is a prime number, if  $x \geq 0, y \geq 0, z > 0$  are integers such that*

$$\begin{cases} x^2 - Fy^2 = f, \\ x^2 - pGz^2 = g, \end{cases}$$

then  $x, y, z, p < N$ .

## 2. Preliminaries

*A. Binary recurring sequences.* Let  $P > 0, Q \neq 0$  with  $D = P^2 - 4Q > 0$ . We gather some properties of  $U_n = U_n(P, Q)$  and  $V_n = V_n(P, Q)$  which will be needed in this paper.

Let  $\alpha, \beta$  be the roots of  $X^2 - PX + Q$ , so

$$\alpha = \frac{P + \sqrt{D}}{2}, \quad \beta = \frac{P - \sqrt{D}}{2},$$

$$\alpha + \beta = P, \quad \alpha\beta = Q, \quad \alpha - \beta = \sqrt{D}.$$

For each  $n \in \mathbb{Z}$ :

$$U_n = \frac{\alpha^n - \beta^n}{\alpha - \beta}, \quad V_n = \alpha^n + \beta^n.$$

If  $m, n \in \mathbb{Z}$ :

$$2U_{m+n} = U_m V_n + U_n V_m,$$

$$U_{m+n} = U_m V_n - Q^n U_{m-n}, \quad V_{m+n} = V_m V_n - Q^n V_{m-n},$$

$$U_{2m} = U_m V_m, \quad V_{2m} = V_m^2 - 2Q^m,$$

$$V_m^2 - DU_m^2 = 4Q^m.$$

The following lemma will be required:

(2.1) LEMMA. *For every  $s \geq 1$  and  $n \in \mathbb{Z}$  we have*

- (a)  $V_n^2 - V_{n-s}V_{n+s} = Q^{n-s}(2Q^s - V_{2s})$ ,
- (b)  $V_n^2 - DU_{n-s}U_{n+s} = Q^{n-s}(2Q^s + V_{2s})$ ,
- (c)  $D(U_n^2 - U_{n-s}U_{n+s}) = Q^{n-s}(-2Q^s + V_{2s})$ ,
- (d)  $DU_n^2 - V_{n-s}V_{n+s} = -Q^{n-s}(2Q^s + V_{2s})$ .

*Proof.* We just prove (a):

$$\begin{aligned} V_n^2 - V_{n-s}V_{n+s} &= (\alpha^n + \beta^n)^2 - (\alpha^{n-s} + \beta^{n-s})(\alpha^{n+s} + \beta^{n+s}) \\ &= \alpha^{2n} + \beta^{2n} + 2(\alpha\beta)^n \\ &\quad - [\alpha^{2n} + \beta^{2n} - (\alpha\beta)^{n-s}(\alpha^{2s} + \beta^{2s})] \\ &= Q^{n-s}(2Q^s - V_{2s}). \blacksquare \end{aligned}$$

Now we assume  $\gcd(P, Q) = 1$ . If  $m, n \geq 1$  and  $d = \gcd(m, n)$  then

$$\begin{aligned} \gcd(U_m, U_n) &= U_d, \\ \gcd(V_m, V_n) &= \begin{cases} V_d & \text{if } m/d \text{ and } n/d \text{ are odd,} \\ 1 \text{ or } 2 & \text{otherwise,} \end{cases} \\ \gcd(U_m, V_n) &= \begin{cases} V_d & \text{if } m/d \text{ is even,} \\ 1 \text{ or } 2 & \text{otherwise,} \end{cases} \\ \gcd(U_n, Q) &= 1, \quad \gcd(V_n, Q) = 1. \end{aligned}$$

If  $a, b \in \mathbb{Z}$ , not both equal to 0, let

$$W_n = aU_n + bV_n \quad \text{for all } n \in \mathbb{Z}.$$

Then

$$W_n = PW_{n-1} - QW_{n-2} \quad \text{for all } n \in \mathbb{Z}.$$

We also have

$$W_{m+n} = W_mV_n - Q^nW_{m-n} \quad \text{for all } m, n \in \mathbb{Z}.$$

If  $\gcd(P, Q) = 1$  then

$$\gcd(W_n, Q) = \gcd(W_1, Q) \quad \text{for all } n \geq 1.$$

The following lemma will also be required:

(2.2) LEMMA. *Assume that  $W_n \neq 0$  for all  $n \geq 1$  and  $\gcd(W_1, Q) = 1$ . If  $t \geq 1$  then  $\gcd(W_n, W_{n+t})$  divides  $W_1W_2 \cdots W_t$  for all  $n \geq 1$ .*

*Proof.* The lemma is trivial for  $n = 1, \dots, t$ . Let  $t \leq n$  and assume the lemma true for  $1, 2, \dots, n$ . Let  $d = \gcd(W_{n+1}, W_{n+1+t})$ . We have  $W_{n+1+t} = W_{n+1}V_t - Q^tW_{n+1-t}$ . Since  $1 \leq n+1-t \leq n$  and  $\gcd(W_{n+1}, Q) = \gcd(W_1, Q) = 1$ , it follows that  $d$  divides  $\gcd(W_{n+1-t}, W_{n+1})$ , which, by induction, divides  $W_1W_2 \cdots W_t$ , thus concluding the proof.  $\blacksquare$

We shall need the following theorem (see Shorey & Tijdeman [14], Shorey & Stewart [13], Pethő [9]), which we quote in the special case needed in this paper.

With the preceding notations:

(2.3) THEOREM. *Assume that  $\gcd(P, Q) = 1$  and  $D \neq 0$ . Let  $a, b \in \mathbb{Z}$ , not both equal to 0, let  $W_n = aU_n + bV_n$  for all  $n \in \mathbb{Z}$ . Let  $A > 0$  be a square-free integer. Then there exists  $N > 0$ , effectively computable in terms of  $P, Q, a, b, A$  such that if  $n \geq 0$  and  $W_n = A\Box$  (where  $\Box$  denotes a non-zero integer which is a square) then  $n < N$ .*

The proof of this theorem involves inequalities of Baker for linear forms in logarithms and the constant  $N$  provided by the proof is usually very large.

For special sequences, the explicit determination of squares and double-squares has been achieved. We quote a few results for sequences with parameters  $P$  even and  $Q = \pm 1$ .

For  $P = 2, Q = 1$ ,  $U_n$  and  $V_n$  are the Pell numbers and we have:

$$\begin{aligned} \{n \mid U_n = \Box\} &= \{1, 7, -1, -7\}, & \{n \mid U_n = 2\Box\} &= \{2\}, \\ \{n \mid V_n = \Box\} &= \emptyset, & \{n \mid V_n = 2\Box\} &= \{0, 1\}. \end{aligned}$$

The above results are due to Ljunggren [5]; the determination of the square Pell numbers required deep arguments.

Ljunggren [5] and Cohn [1] studied the sequences of numbers  $U_n(4, -1)$  and  $V_n(4, -1)$ :

$$\begin{aligned} \{n \mid U_n(4, -1) = \Box\} &= \{1, 2, -1\}, & \{n \mid U_n(4, -1) = 2\Box\} &= \{4\}, \\ \{n \mid V_n(4, -1) = \Box\} &= \{1\}, & \{n \mid V_n(4, -1) = 2\Box\} &= \{0, 2, -2\}. \end{aligned}$$

In [1] Cohn obtained more results about squares and double squares in the sequences  $U_n(2c, \pm 1)$  and  $V_n(2c, \pm 1)$ , for special values of  $2c$ .

The reader may obtain more information about recurring sequences in Ribenboim [11] (see Chapter 1 entitled “The Fibonacci Numbers and the Arctic Ocean”). More specifically about Pell numbers, see Ribenboim [10].

*B. Pell equations.* We keep the same notations:  $F > 1$ ,  $f \neq 0$ ,  $\varepsilon = c + d\sqrt{F}$  is the fundamental unit of  $\mathbb{Z}[\sqrt{F}]$ , so  $c \geq 1$ ,  $d \geq 1$ ;  $P = 2c$ ,  $Q = c^2 - Fd^2 = \pm 1$ ,  $U_n = U_n(2c, Q)$ ,  $V_n = V_n(2c, Q)$ . We consider solutions of  $x^2 - Fy^2 = f$ .

Two solutions  $(x, y)$  and  $(x', y')$  of the Pell equation are said to be *equivalent* if there exists  $n \in \mathbb{Z}$  such that

$$Q^n = 1 \quad \text{and} \quad \frac{x + y\sqrt{F}}{x' + y'\sqrt{F}} = \varepsilon^n.$$

If  $Q = 1$  let  $c' + d'\sqrt{F} = c + d\sqrt{F} = \varepsilon$ . If  $Q = -1$  let

$$c' + d'\sqrt{F} = (c^2 + d^2F) + 2cd\sqrt{F} = \varepsilon^2.$$

We note that if  $c = 1$  then  $Q = -1$  so  $c' > 1$ .

A solution  $(a, b)$  with  $a \geq 0$  and  $b \geq 0$  is called a *fundamental solution* if the following inequalities are satisfied:

$$0 \leq a \leq \sqrt{\frac{(c' + \delta)|f|}{2}}, \quad 0 \leq b \leq d' \sqrt{\frac{|f|}{2(c' + \delta)}}$$

where  $\delta = |f|/f$ .

Nagell proved (see [7] and [8]):

(2.4) THEOREM. *Every solution  $(x, y)$  with  $x \geq 0, y \geq 0$  of  $x^2 - Fy^2 = f$  is equivalent to a fundamental solution.*

**3. Proof of Theorem (1.1).** We divide the proof into three parts.

1°) Let  $S$  be the set of all  $(x, y, z, p)$  such that  $x \geq 0, y \geq 0, z > 0, p = 1$  or  $p$  is a prime number and

$$\begin{cases} x^2 - Fy^2 = f, \\ x^2 - pGz^2 = fk_s. \end{cases}$$

[The proof when  $g = fh_s$  is similar and will not be given.]

Let  $T$  be the set of all  $(x, y)$  such that  $x \geq 0, y \geq 0, x^2 - Fy^2 = f$  and there exists  $(z, p)$  such that  $(x, y, z, p) \in S$ . Clearly, it suffices to show that the set  $T$  is effectively computable.

By the theorem of Nagell (2.4) if the equation  $x^2 - Fy^2 = f$  has solutions, then it has a non-empty effectively computable set of fundamental solutions and every solution  $(x, y)$ , with  $x \geq 0, y \geq 0$  is given by a relation  $x + y\sqrt{F} = (a + b\sqrt{F})\varepsilon^n$ , where  $a \geq 0, b \geq 0, (a, b)$  is a fundamental solution,  $\varepsilon = c + d\sqrt{F}$  is the fundamental unit of  $\mathbb{Z}[\sqrt{F}]$ ,  $Q = c^2 - d^2F, Q^n = 1$ .

We fix an arbitrary fundamental solution  $(a, b)$  and write

$$x_n + y_n\sqrt{F} = (a + b\sqrt{F})\varepsilon^n = (a + b\sqrt{F})\left(\frac{V_n}{2} + dU_n\sqrt{F}\right),$$

where  $U_n = U_n(2c, Q)$  and  $V_n = V_n(2c, Q)$ . So

$$x_n = a\frac{V_n}{2} + bdFU_n, \quad y_n = adU_n + b\frac{V_n}{2}.$$

It suffices to show that the set  $R = \{n > s \mid (x_n, y_n) \in T\}$  is effectively bounded.

2°) We show that if  $Q^n = 1$  and  $s \geq 1$  then

$$x_n^2 - Fy_{n-s}y_{n+s} = fk_s.$$

Indeed:

$$x_n^2 = \left(a\frac{V_n}{2} + bdFU_n\right)^2 = a^2\frac{V_n^2}{4} + \frac{b^2F}{4}DU_n^2 + abdFU_{2n}.$$

Next

$$\begin{aligned}
Fy_{n-s}y_{n+s} &= F\left(b\frac{V_{n-s}}{2} + adU_{n-s}\right)\left(b\frac{V_{n+s}}{2} + adU_{n+s}\right) \\
&= \frac{b^2F}{4}V_{n-s}V_{n+s} + abdFU_{2n} + a^2d^2FU_{n-s}U_{n+s} \\
&= \frac{b^2F}{4}V_{n-s}V_{n+s} + abdFU_{2n} + \frac{a^2}{4}DU_{n-s}U_{n+s}.
\end{aligned}$$

In the above calculation we used identities indicated in Section 2. It follows that

$$x_n^2 - Fy_{n-s}y_{n+s} = \frac{a^2}{4}(V_n^2 - DU_{n-s}U_{n+s}) + \frac{b^2F}{4}(DU_n^2 - V_{n-s}V_{n+s}).$$

By Lemma (2.1) we have

$$\begin{aligned}
x_n^2 - Fy_{n-s}y_{n+s} &= \frac{a^2}{4}Q^{ns}(2Q^s + V_{2s}) - \frac{b^2F}{4}Q^{ns}(2Q^s + V_{2s}) \\
&= \frac{f}{4}(2 + Q^sV_{2s}) = fh_s,
\end{aligned}$$

because  $Q^n = 1$ . [For the proof of the theorem when  $g = fh_s$  we need the relation

$$x_n^2 - x_{n-s}x_{n+s} = fh_s,$$

which is established in a similar way.]

3°) By Lemma (2.2) for every  $n > s$ ,  $\gcd(y_{n-s}, y_{n+s})$  divides  $y_1y_2 \cdots y_{2s}$ ; we note that the integer  $y_1y_2 \cdots y_{2s}$  is effectively computable in terms of  $F$ ,  $s$  and the chosen fundamental solution  $(a, b)$ . For every positive integer  $e$  dividing  $y_1y_2 \cdots y_{2s}$ , let

$$R_e = \{n \in R \mid \gcd(y_{n-s}, y_{n+s}) = e\}.$$

Let  $n \in R_e$ , so from  $Fy_{n-s}y_{n+s} = pGz^2$  it follows that

$$e^2F^2\frac{y_{n-s}}{e} \cdot \frac{y_{n+s}}{e} = pFG\Box = p^\delta H\Box,$$

where  $\delta = 0$  or  $1$ ,  $p \nmid H$ ,  $H$  is square-free and  $H \mid FG$ . Hence

$$\frac{y_{n-s}}{e} \cdot \frac{y_{n+s}}{e} = p^\delta H\Box.$$

Let  $\mathcal{H} = \{(H', H'') \mid H'H'' \text{ is square-free, } H'H'' \mid FG \text{ and } \gcd(H', H'') = 1\}$ , so  $\mathcal{H}$  is effectively computable. For each  $(H', H'')$  let  $R'_{e, (H', H'')}$  be the set of all  $n \in R_e$  such that

$$\frac{y_{n-s}}{e} = p^\delta H'\Box, \quad \frac{y_{n+s}}{e} = H''\Box.$$

By Theorem (2.3) the set  $\{n+s \mid n \in R'_{e, (H', H'')}\}$ , hence also  $R'_{e, (H', H'')}$  is effectively bounded. Similarly, let  $R''_{e, (H', H'')}$  be the set of  $n \in R_e$  such

that

$$\frac{y_{n-s}}{e} = H'\square, \quad \frac{y_{n+s}}{e} = p^\delta H''\square.$$

Then again  $R''_{e,(H',H'')}$  is effectively bounded. Since this holds for each  $(H', H'') \in \mathcal{H}$  the set  $R_e$  is effectively bounded for each  $e \mid y_1 y_2 \cdots y_{2s}$ . Hence  $R$  is effectively bounded and this concludes the proof of the theorem. ■

**4. A numerical example.** We give an example where our method is applied with success to determine explicitly all solutions. To begin we prove a lemma.

(4.1) LEMMA. *Let  $U_n, V_n$  be the Pell numbers for all  $n \in \mathbb{Z}$ . Then*

- (a)  $\{n \neq 0 \mid U_n = 3\square\} = \{4\}$ ,  $\{n \mid V_n = 3\square\} = \emptyset$ ,
- (b)  $\{n \neq 0 \mid U_n = 6\square\} = \emptyset$ .

*Proof.* (a) Let  $U_n = 3\square$ . By considering the sequence  $U_n$  modulo 3 we see that 4 divides  $n$ . Let  $n = 4h$ , so  $U_n = U_{2h}V_{2h}$ , with  $\gcd(U_{2h}, V_{2h}) = 2$ . Then either  $V_{2h} = 2\square$  or  $U_{2h} = 2\square$ . So  $h = 1$  and  $n = 4$ . If  $V_n = 3\square$ , since  $2 \mid V_n$  but  $4 \nmid V_n$  this is impossible.

(b) If  $U_n = 6\square$  then  $n = 4h$  and we have the following cases:

$$\begin{array}{l} U_{2h} = 3\square \mid 6\square \mid \square \mid 2\square, \\ V_{2h} = 2\square \mid \square \mid 6\square \mid 3\square. \end{array}$$

From (a) and the knowledge of  $m$  such that  $U_m = \square, 2\square$ ,  $V_m = \square, 2\square$  we conclude that  $n = 0$ . ■

(4.2) EXAMPLE. *If  $p$  is a prime, or  $p = 1$ , if  $x, y, z$  are positive integers and*

$$\begin{cases} x^2 - 2y^2 = 1, \\ x^2 - pz^2 = 9, \end{cases}$$

*then  $(x, y, z, p) = (99, 70, 24, 17)$ .*

*Proof.*  $\varepsilon = 1 + \sqrt{2}$  is the fundamental unit of  $\mathbb{Z}[\sqrt{2}]$ , let  $P = 2$ ,  $Q = -1$ ,  $U_n, V_n$  are the Pell numbers,  $k_2 = 9$ , so the method is applicable.  $\varepsilon^2 = 3 + 2\sqrt{2}$  is the fundamental solution of the first equation,  $x_n + y_n\sqrt{2} = \varepsilon^{n+2} = V_{n+2}/2 + U_{n+2}\sqrt{2}$  and we work with  $n$  even since  $Q = -1$ . We have:

$$2y_{n-2}y_{n+2} = pz^2 \neq 0,$$

that is,  $U_n U_{n+4} = 2p\square \neq 0$ .

(a)  $p = 2$ , so  $U_n U_{n+4} = \square$ . If  $n \equiv 2 \pmod{4}$  then  $\gcd(U_n, U_{n+4}) = U_2 = 2$ , so  $U_n = 2\square$  and  $U_{n+4} = 2\square$ , which is impossible. If  $n \equiv 0 \pmod{4}$  then  $\gcd(U_n, U_{n+4}) = U_4 = 12$ , hence  $U_n = 3\square$ ,  $U_{n+4} = 3\square$ , which is impossible by (4.1).

(b)  $p \neq 2$ , so  $U_n U_{n+4} = 2p\Box$ . If  $n \equiv 2 \pmod{4}$  then  $U_n \equiv U_{n+4} \equiv 2 \pmod{4}$  so the 2-adic value of  $U_n U_{n+4}$  is even, hence  $U_n U_{n+4} = 2p\Box$  is impossible.

Now let  $n \equiv 0 \pmod{4}$ , so the following cases are possible:

$$\begin{array}{cccc} U_n = & 3\Box & \Big| & 6\Box & \Big| & 3p\Box & \Big| & 6p\Box, \\ U_{n+4} = & 6p\Box & \Big| & 3p\Box & \Big| & 6\Box & \Big| & 3\Box. \\ & (1) & & (2) & & (3) & & (4) \end{array}$$

(1) It was seen that  $n = 4$ , hence  $U_8 = 408 = 6 \times 17\Box$  so  $p = 17$ ,  $x_4 = V_6/2 = 99$ ,  $y_4 = U_6 = 70$  and this gives the solution  $(x, y, z, p) = (99, 70, 24, 17)$ .

(4)  $n + 4 = 4$ ,  $n = 0$  which is impossible, since then  $z = 0$ .

(2) and (3) are impossible as it was shown in Lemma (4.1). ■

As an exercise the reader may wish to show that if  $x, y, z$  are positive, if  $p$  is a prime number and if

$$\begin{cases} x^2 - 3y^2 = -3, \\ x^2 - 3pz^2 = -12, \end{cases}$$

then  $(x, y, z, p) = (3, 2, 1, 7)$ .

The reader may employ the same method to prove the original result of Ljunggren mentioned in the Introduction: if  $x, y$  are positive integers, if  $p$  is a prime number and  $x^4 - py^2 = 1$  then  $(x, y, p) = (3, 4, 5)$  or  $(99, 1820, 29)$ .

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