

Families of periodic Jacobi–Perron algorithms in any dimension with period lengths going to infinity

by

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1. Introduction. Let $\alpha = (\alpha_1, \dots, \alpha_n)$ be a vector in \mathbb{R}^n ($n \geq 1$). The *Jacobi–Perron Algorithm* (JPA) expansion [6] of α is given by two sequences:

- $(a^{(\nu)})_{\nu \geq 0}$ in \mathbb{Z}^n where $a^{(\nu)} = (a_1^{(\nu)}, \dots, a_n^{(\nu)})$,
- $(\alpha^{(\nu)})_{\nu \geq 0}$ in \mathbb{R}^n where $\alpha^{(\nu)} = (\alpha_1^{(\nu)}, \dots, \alpha_n^{(\nu)})$,

defined by:

$$\begin{cases} \alpha^{(0)} = \alpha, \\ a_i^{(\nu)} = [\alpha_1^{(\nu)}] \quad \text{for } \nu \geq 0 \text{ and } 1 \leq i \leq n, \\ \alpha_n^{(\nu+1)} = \frac{1}{\alpha_1^{(\nu)} - a_1^{(\nu)}} \quad \text{if } \alpha_1^{(\nu)} \neq a_1^{(\nu)}, \\ \alpha_i^{(\nu+1)} = \frac{\alpha_{i+1}^{(\nu)} - a_{i+1}^{(\nu)}}{\alpha_1^{(\nu)} - a_1^{(\nu)}} \quad \text{for } 1 \leq i < n, \end{cases}$$

where $[x]$ is the integer part of x .

The JPA expansion is *periodic* if there exist integers $k \geq 0$ and $l > 0$ such that $a_i^{(k+\nu)} = a_i^{(k+\nu+l)}$ for all $\nu \geq 0$ and $0 < i \leq n$. If k and l are the least integers satisfying this equality then k is the *preperiod length* and l is the *period length*. If $k = 0$, the expansion is *purely periodic*.

If the JPA expansion of $\alpha = (\alpha_1, \dots, \alpha_n)$ is purely periodic with period length l , then

$$\epsilon = \prod_{\nu=0}^{l-1} \alpha_n^{(\nu)}$$

is a unit $K = \mathbb{Q}[\alpha_1, \dots, \alpha_n]$. This is a unit found by Hasse and Bernstein [3].

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2. A family of expansions. For all fixed integers $c \geq 2, m, n \geq 1$ and $1 \leq t \leq n$, let

$$f_t(X) = X^{n+1} - c^m X^n - (c - 1)X^t - \sum_{i=1}^{t-1} (c^m - 1)X^i - c^m.$$

LEMMA 1. $f_t(X)$ has only one real root α such that

$$\begin{cases} c^m + c - 1 < \alpha < c^m + c & \text{for } t = n, \\ c^m < \alpha < c^m + 1 & \text{for } 1 \leq t \leq n - 1. \end{cases}$$

Proof. We use the rule of signs of Descartes [5]: the possible number of positive roots of a polynomial is equal to the number of sign changes in the coefficients of the polynomial or less than that number by a multiple of 2.

It follows that $f_t(X)$ has only one positive real root. Moreover,

$$\begin{cases} f_t(c^m + c - 1) < 0 \text{ and } f_t(c^m + c) > 0 & \text{if } t = n, \\ f_t(c^m) < 0 \text{ and } f_t(c^m + 1) > 0 & \text{if } 1 \leq t \leq n - 1. \end{cases}$$

So, the result is proved. ■

Write $f_t(X) = X^{n+1} - \sum_{i=0}^n b_t^{(i)} X^i$. We define

$$\begin{cases} \alpha_n = \alpha, \\ \alpha_i = \alpha(\alpha_{i+1} - b_t^{(i+1)}) & \text{for } 1 \leq i \leq n - 1. \end{cases}$$

2.1. Main theorem. The main result of this paper is

THEOREM 2. *The JPA expansion of $(\alpha_1, \dots, \alpha_n)$ is purely periodic with period length $l = (n + t)m + 1$. Moreover $\epsilon = \alpha \left(\frac{\alpha^t}{\alpha - c^m}\right)^m$ is a unit of $\mathbb{Q}(\alpha)$.*

REMARK. If $n = 2$ for $t = 1$ and $t = 2$ we obtain the two families given by C. Levesque and G. Rhin [4]. These results are also explained in my thesis [1].

2.2. Proof of the theorem

(1) *The case $n = t = 1$ (continued fraction).* Writing α for the real root of $f(X) = X^2 - (c^m + c - 1)X - c^m$ such that $c^m + c - 1 < \alpha < c^m + c$, we have to prove that the continued fraction expansion of α is purely periodic with period length $l = 2m + 1$ and that $\epsilon = \alpha \left(\frac{\alpha}{\alpha - c^m}\right)^m$ is a unit of $\mathbb{Q}(\alpha)$.

We will use the following formulas:

$$\alpha - (c^m + c - 1) = \frac{c^m}{\alpha}, \quad \alpha - c^m = \frac{c\alpha}{\alpha + 1}.$$

The algorithm starts with

$$\alpha_1^{(0)} = \alpha, \quad a_1^{(0)} = c^m + c - 1.$$

Next,

$$\left\{ \begin{array}{l} \alpha_1^{(1)} = \frac{1}{\alpha - (c^m + c - 1)} = \frac{\alpha}{c^m}, \quad a_1^{(1)} = 1, \\ \alpha_1^{(2)} = \frac{1}{\alpha/c^m - 1} = \frac{c^m}{\alpha - c^m}, \\ \quad = \frac{c^m(\alpha + 1)}{c\alpha} = c^{m-1} + \frac{c^{m-1}}{\alpha}, \quad a_1^{(2)} = c^{m-1}. \end{array} \right.$$

Now, by induction, it is easy to prove the following formulas (for $0 \leq s \leq m - 1$):

$$\left\{ \begin{array}{l} \alpha_1^{(2s+1)} = \alpha/c^{m-s}, \quad a_1^{(2s+1)} = c^s, \\ \alpha_1^{(2s+2)} = c^{m-s-1} + c^{m-s-1}/\alpha, \quad a_1^{(2s+2)} = c^{m-s-1}. \end{array} \right.$$

So, when $s = m - 1$ we have

$$\alpha_1^{(2m)} = 1 + 1/\alpha, \quad a_1^{(2m)} = 1,$$

and finally we obtain

$$\alpha_1^{(2m+1)} = \alpha = \alpha_1^{(0)},$$

which leads to $l = 2m + 1$.

The Hasse–Bernstein formula produces in $\mathbb{Q}(\alpha)$ the unit

$$\epsilon = \prod_{\nu=0}^{l-1} \alpha_1^{(\nu)} = \alpha \left(\frac{\alpha}{\alpha - c^m} \right)^m.$$

(2) *The case $n \neq 1$.* In the course of the proof, the equality $1 + \alpha + \alpha^2 + \dots + \alpha^{t-1} + \alpha^n = \frac{c\alpha^t}{\alpha - c^m}$ will be useful.

Assume $t \neq 1$. For the first vectors, we easily obtain

$$\left\{ \begin{array}{l} \alpha_n^{(0)} = \alpha, \quad a_n^{(0)} = \begin{cases} c^m + c - 1 & \text{if } t = n, \\ c^m & \text{if } \leq t \leq n - 1; \end{cases} \\ \text{for } t < i < n, \\ \alpha_i^{(0)} = \alpha^{n-i}(\alpha - c^m) \\ \quad = (c - 1) \frac{\alpha^t}{\alpha^i} + \sum_{r=1}^{t-1} \frac{(c^m - 1)\alpha^{t-r}}{\alpha^i} + \frac{c^m}{\alpha^i}, \quad a_i^{(0)} = 0; \\ \alpha_t^{(0)} = \alpha^{n-t}(\alpha - c^m) \\ \quad = (c - 1) + \sum_{r=1}^{t-1} \frac{c^m - 1}{\alpha^r} + \frac{c^m}{\alpha^t}, \quad a_t^{(0)} = c - 1; \\ \text{for } 1 \leq i < t, \\ \alpha_i^{(0)} = \sum_{r=0}^{i-1} \frac{c^m - 1}{\alpha^r} + \frac{c^m}{\alpha^i}, \quad a_i^{(0)} = c^m - 1. \end{array} \right.$$

Now, we show by induction on s that, for $0 \leq s \leq m - 1$,

$$(2.1) \quad \begin{cases} \text{for } 2 \leq i \leq n, \\ \alpha_i^{((n+t)s)} = \alpha_i^{(0)}, & a_i^{((n+t)s)} = a_i^{(0)}; \\ \alpha_1^{((n+t)s)} = (c^{m-s} - 1) + c^{m-s}/\alpha, & a_1^{((n+t)s)} = c^{m-s} - 1. \end{cases}$$

These formulas hold for $s = 0$. Assume they hold for all $0 \leq s \leq m - 1$. Then

$$\left\{ \begin{array}{l} \alpha_n^{((n+t)s+1)} = \alpha/c^{m-s}, & a_n^{((n+t)s+1)} = c^s; \\ \text{for } t \leq i < n, \\ \alpha_i^{((n+t)s+1)} = \alpha^{n-i}(\alpha - c^m)/c^{m-s}, & a_i^{((n+t)s+1)} = 0; \\ \text{for } 1 \leq i < t, \\ \alpha_i^{((n+t)s+1)} = \frac{\alpha}{c^{m-s}} \sum_{r=1}^i \frac{c^m - 1}{\alpha^r} + \frac{c^s}{\alpha^i} \\ = \frac{1}{c^{m-s}\alpha^i} \left(c^m + \sum_{r=1}^i \alpha^r(c^m - 1) \right), & a_i^{((n+t)s+1)} = c^s - 1; \end{array} \right.$$

and

$$\left\{ \begin{array}{l} \alpha_n^{((n+t)s+2)} = \frac{\alpha c^{m-s}}{\alpha(c^{m-s} - 1) + c^m}, & a_n^{((n+t)s+2)} = 1; \\ \text{for } t - 1 \leq i < n, \\ \alpha_i^{((n+t)s+2)} = \frac{\alpha^{n-i}(\alpha - c^m)}{\alpha(c^{m-s} - 1) + c^m}, & a_i^{((n+t)s+2)} = 0; \\ \text{for } 1 \leq i < t - 1, \\ \alpha_i^{((n+t)s+2)} = \frac{\alpha^{i+1}(c^{m-s} - 1) + c^m + \sum_{r=1}^i \alpha^r(c^m - 1)}{\alpha^i(\alpha(c^{m-s} - 1) + c^m)}, & a_i^{((n+t)s+2)} = 0. \end{array} \right.$$

Then it is easy to prove, by induction on ν with $3 \leq \nu \leq t$, that

$$\left\{ \begin{array}{l} \alpha_n^{((n+t)s+\nu)} = \frac{\alpha^{\nu-1}(c^{m-s} - 1) + c^m\alpha + \sum_{r=2}^{\nu-2} \alpha^r(c^m - 1)}{\alpha^{\nu-1}(c^{m-s} - 1) + c^m + \sum_{r=1}^{\nu-2} \alpha^r(c^m - 1)}, & a_n^{((n+t)s+\nu)} = 1; \\ \text{for } t + 1 - \nu \leq i < n, \\ \alpha_i^{((n+t)s+\nu)} = \frac{\alpha^{n-i}(\alpha - c^m)}{\alpha^{\nu-1}(c^{m-s} - 1) + c^m + \sum_{r=1}^{\nu-2} \alpha^r(c^m - 1)}, & a_i^{((n+t)s+\nu)} = 0; \\ \text{for } 1 \leq i < t + 1 - \nu, \\ \alpha_i^{((n+t)s+\nu)} = \frac{\alpha^{i+\nu-1}(c^{m-s} - 1) + c^m + \sum_{r=1}^{i+\nu-2} \alpha^r(c^m - 1)}{\alpha^i(\alpha^{\nu-1}(c^{m-s} - 1) + c^m + \sum_{r=1}^{\nu-2} \alpha^r(c^m - 1))}, & a_i^{((n+t)s+\nu)} = 0. \end{array} \right.$$

As can be verified, all the formulas hold for $\nu = 3$. If we assume that they hold for ν , it is again easy to prove them for $\nu + 1$.

Next,

$$\left\{ \begin{array}{l} \alpha_n^{((n+t)s+t+1)} = \frac{\alpha^{t-1}(c^{m-s} - 1) + c^m + \sum_{r=1}^{t-2} \alpha^r (c^m - 1)}{\alpha^{n-1}(\alpha - c^m)}, \\ \text{for } 1 \leq i < n, \\ \alpha_i^{((n+t)s+t+1)} = 1/\alpha^i, \end{array} \right. \quad \begin{array}{l} a_n^{((n+t)s+t+1)} = c^{m-s-1}, \\ \\ a_i^{((n+t)s+t+1)} = 0. \end{array}$$

Indeed, let us give the details for $a_n^{((n+t)s+t+1)} = c^{m-s-1}$. We have

$$\begin{aligned} \alpha_n^{((n+t)s+t+1)} &= \frac{\alpha^{t-1}c^{m-s} - (\alpha - c^m) \sum_{r=1}^{t-1} \alpha^{r-1}}{\alpha^{n-1}(\alpha - c^m)} \\ &= \frac{c^{m-s-1}}{\alpha^n} \left(\frac{c\alpha^t}{\alpha - c^m} \right) - \sum_{r=1}^{t-1} \frac{1}{\alpha^{n-r}}, \end{aligned}$$

that is,

$$\alpha_n^{((n+t)s+t+1)} = \frac{c^{m-s-1}}{\alpha^n} (1 + \alpha + \alpha^2 + \dots + \alpha^{t-1} + \alpha^n) - \sum_{r=1}^{t-1} \frac{1}{\alpha^{n-r}}.$$

or

$$\alpha_n^{((n+t)s+t+1)} = c^{m-s-1} + \frac{1}{\alpha^n} \left((c^{m-s-1} - 1) \sum_{r=1}^{t-1} \alpha^r + c^{m-s-1} \right),$$

and finally, using the fact that $0 < (c^{m-s-1} - 1) \sum_{r=1}^{t-1} \alpha^r + c^{m-s-1} < \alpha^n$, we obtain $a_n^{((n+t)s+t+1)} = c^{m-s-1}$.

By induction on ν with $2 \leq \nu \leq n - t + 1$, it is still easy to prove the following formula:

$$\left\{ \begin{array}{l} \text{for } n - \nu + 2 \leq i \leq n, \\ \alpha_i^{((n+t)s+t+\nu)} = \alpha_i^{(0)}, \quad a_i^{((n+t)s+t+\nu)} = \alpha_i^{(0)}; \\ \alpha_{n-\nu+1}^{((n+t)s+t+\nu)} = \frac{c^{m-s-1} - 1}{\alpha^{n-\nu+1}} \left(\sum_{r=1}^{t-1} \alpha^r \right) + \frac{c^{m-s-1}}{\alpha^{n-\nu+1}}, \quad a_{n-\nu+1}^{((n+t)s+t+\nu)} = 0; \\ \text{for } 1 \leq i \leq n - \nu, \\ \alpha_i^{((n+t)s+t+\nu)} = 1/\alpha^i, \quad a_i^{((n+t)s+t+\nu)} = 0. \end{array} \right.$$

All the formulas hold for $\nu = 2$ and if we assume they hold for ν , it is again easy to prove them for $\nu + 1$.

Then, always by induction on ν with $2 \leq \nu \leq t$ we have

$$\left\{ \begin{array}{l} \text{for } t - \nu + 2 \leq i \leq n, \\ \alpha_i^{((n+t)s+n+\nu)} = \alpha_i^{(0)}, \quad a_i^{((n+t)s+n+\nu)} = a_i^{(0)}; \\ \alpha_{t-\nu+1}^{((n+t)s+n+\nu)} = \frac{c^{m-s-1} - 1}{\alpha^{t-\nu+1}} \left(\sum_{r=1}^{t-\nu+1} \alpha^r \right) + \frac{c^{m-s-1}}{\alpha^{t-\nu+1}}, \\ \quad a_{t-\nu+1}^{((n+t)s+n+\nu)} = c^{m-s-1} - 1; \\ \text{for } 1 \leq i \leq t - \nu, \\ \alpha_i^{((n+t)s+n+\nu)} = 1/\alpha^i, \quad a_i^{((n+t)s+n+\nu)} = 0. \end{array} \right.$$

In the special case of $\nu = t$,

$$\left\{ \begin{array}{l} \text{for } 2 \leq i \leq n, \\ \alpha_i^{((n+t)s+n+t)} = \alpha_i^{(0)}, \quad a_i^{((n+t)s+n+t)} = a_i^{(0)}; \\ \alpha_1^{((n+t)s+n+t)} = (c^{m-s-1} - 1) + c^{m-s-1}/\alpha, \quad a_1^{((n+t)s+n+t)} = c^{m-s-1} - 1. \end{array} \right.$$

Therefore, (2.1) holds for $s + 1$, and so all the previous formulas follow by induction.

For $t = 1$, it is again easy to prove that, for $0 \leq s \leq m$,

$$\left\{ \begin{array}{l} \alpha_i^{((n+t)s)} = \alpha_i^{(0)} \quad \text{for } 2 \leq i \leq n, \\ \alpha_1^{((n+t)s)} = \begin{cases} c^{m-s}/\alpha & \text{if } 1 \leq s \leq m, \\ c - 1 + c^m/\alpha & \text{if } s = 0. \end{cases} \end{array} \right.$$

So, for all t such that $1 \leq t \leq n$,

$$\left\{ \begin{array}{l} \text{for } 2 \leq i \leq n, \\ \alpha_i^{((n+t)m)} = \alpha_i^{(0)}, \quad a_i^{((n+t)m)} = a_i^{(0)}; \\ \alpha_1^{((n+t)m)} = 1/\alpha, \quad a_1^{((n+t)m)} = 0. \end{array} \right.$$

Finally, we obtain

$$\alpha_i^{((n+t)m+1)} = \alpha_i^{(0)}, \quad a_i^{((n+t)m+1)} = a_i^{(0)} \quad \text{for } 1 \leq i \leq n,$$

leading to $l = (n + t)m + 1$.

We deduce that the Hasse–Bernstein formula provides the unit

$$\epsilon = \prod_{\nu=0}^{l-1} \alpha_n^{(\nu)} = \alpha \left(\frac{\alpha^t}{\alpha - c^m} \right)^m . \blacksquare$$

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