COLLOQUIUM MATHEMATICUM

VOL. LXVIII

THE SOLVABILITY OF THE DIOPHANTINE EQUATION $D_1 x^2 - D_2 y^4 = 1$

BY

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1. Introduction. Let \mathbb{Z} , \mathbb{N} denote the sets of integers and positive integers respectively. Let $D_1, D_2 \in \mathbb{N}$ such that $gcd(D_1, D_2) = 1$ and D_1D_2 is not a square. Many papers concerning the equation

(1)
$$D_1 x^2 - D_2 y^4 = 1, \quad x, y \in \mathbb{N},$$

were written by Cohn, Ljunggren, Maĭrullin, Mordell and Obláth. In this paper we deal with the solvability of (1). Clearly, if (x, y) is a solution of (1), then (x, y^2) is a solution of the equation

(2)
$$D_1 u^2 - D_2 v^2 = 1, \quad u, v \in \mathbb{Z}$$

with x > 0 and $y^2 > 0$. Since D_1D_2 is not a square, (2) has a unique solution (u_1, v_1) such that $u_1 > 0$, $v_1 > 0$ and $u_1\sqrt{D_1} + v_1\sqrt{D_2} \le u\sqrt{D_1} + v\sqrt{D_2}$, where (u, v) runs over all solutions of (2) with u > 0 and v > 0. The solution (u_1, v_1) is called the *least solution* of (2). In this paper, using the Ko–Terjanian–Rotkiewicz method (cf. [3]), we prove the following result:

THEOREM. If $\min(D_1, D_2) > 1$, then (1) has solutions (x, y) if and only if the least solution (u_1, v_1) of (2) satisfies

(3) $v_1 = dk^2, \quad d, k \in \mathbb{N}, \ d \ is \ square \ free,$

and $(\varepsilon_1^d - \overline{\varepsilon}_1^d)/(2\sqrt{D_2})$ is a square, where

(4)
$$\varepsilon_1 = u_1 \sqrt{D_1} + v_1 \sqrt{D_2}, \quad \overline{\varepsilon}_1 = u_1 \sqrt{D_1} - v_1 \sqrt{D_2}.$$

2. Lemmas

LEMMA 1 ([2]). For $\min(D_1, D_2) > 1$, if (2) has solutions (u, v), then all solutions (u, v) of (2) with u > 0 and v > 0 are given by

$$u\sqrt{D_1} + v\sqrt{D_2} = (u_1\sqrt{D_1} + v_1\sqrt{D_2})^t$$

where $t \in \mathbb{N}$ with $2 \nmid t$, and (u_1, v_1) is the least solution of (2).

1991 Mathematics Subject Classification: 11D25, 11A15.

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Supported by the National Natural Science Foundation of China.

LEMMA 2 ([1, p. 117]). For any $n \in \mathbb{N}$ and any complex numbers α , β ,

$$\alpha^{n} + \beta^{n} = \sum_{i=0}^{[n/2]} (-1)^{i} \frac{n}{i} \binom{n-i-1}{i-1} (\alpha + \beta)^{n-2i} (\alpha \beta)^{i}.$$

LEMMA 3. For $\min(D_1, D_2) > 1$, let (u, v) be a solution of (2) with u > 0, v > 0, and let

$$\varepsilon = u\sqrt{D_1} + v\sqrt{D_2}, \quad \overline{\varepsilon} = u\sqrt{D_1} - v\sqrt{D_2}.$$

Further, for any $m \in \mathbb{Z}$ with $2 \nmid m$, let

(5)
$$F(m) = \frac{\varepsilon^m - \overline{\varepsilon}^m}{\varepsilon - \overline{\varepsilon}}.$$

Then the $F(m) \in \mathbb{Z}$ satisfy:

- (i) F(m) = -F(-m).
- (ii) If m > 0, then $F(m) \in \mathbb{N}$ satisfies $F(m) \equiv m \pmod{4D_2v^2}$.
- (iii) For any $m' \in \mathbb{Z}$ with $2 \nmid m'$, $F(m) \equiv F(m 2m') \pmod{F(m')}$.

Proof. Since $\varepsilon \overline{\varepsilon} = 1$, we have F(m) = -F(-m). For m > 0, by Lemma 2, we get

(6)
$$F(m) = \frac{\varepsilon^m + (-\overline{\varepsilon})^m}{\varepsilon + (-\overline{\varepsilon})} = \sum_{i=0}^{(m-1)/2} \frac{m}{i} {m-i-1 \choose i-1} (\varepsilon - \overline{\varepsilon})^{m-2i-1} (\varepsilon \overline{\varepsilon})^i$$
$$= \sum_{i=0}^{(m-1)/2} \frac{m}{i} {m-i-1 \choose i-1} (4D_2 v^2)^{(m-1)/2-i}$$
$$\equiv m \pmod{4D_2 v^2}.$$

This implies (ii).

For any $m, m' \in \mathbb{Z}$ with $2 \nmid mm'$, by Lemma 2, we have

$$\varepsilon^{m-m'} + \overline{\varepsilon}^{m-m'} = \varepsilon^{|m-m'|} + \overline{\varepsilon}^{|m-m'|}$$
$$= \sum_{j=0}^{|m-m'|/2} (-1)^j \frac{|m-m'|}{j} \binom{|m-m'|-j-1}{j-1}$$
$$\times (4D_1 u^2)^{|m-m'|/2-j} \in \mathbb{Z}.$$

Hence, from

$$\frac{\varepsilon^m - \overline{\varepsilon}^m}{\varepsilon - \overline{\varepsilon}} = (\varepsilon\overline{\varepsilon})^{m'} \left(\frac{\varepsilon^{m-2m'} - \overline{\varepsilon}^{m-2m'}}{\varepsilon - \overline{\varepsilon}} \right) + (\varepsilon^{m-m'} + \overline{\varepsilon}^{m-m'}) \left(\frac{\varepsilon^{m'} - \overline{\varepsilon}^{m'}}{\varepsilon - \overline{\varepsilon}} \right)$$

we see that (iii) is true. The lemma is proved.

LEMMA 4. Let $m, m_1 \in \mathbb{N}$ with $m > m_1 > 1$ and $gcd(m, m_1) = 1$. Then there exist $m_2, \ldots, m_s, a_1, \ldots, a_{s-1} \in \mathbb{N}$ such that

(7)
$$m_1 > m_2 > \ldots > m_s = 1, \quad 2 \nmid m_2 \ldots m_s,$$

(8)
$$m = 2a_1m_1 + \delta_1m_2, \quad m_{j-1} = 2a_jm_j + \delta_jm_{j+1}, \quad j = 2, \dots, s-1,$$

where $\delta_i \in [-1, 1]$ for $i = 1, \dots, s-1$

where $\delta_i \in \{-1, 1\}$ for i = 1, ..., s - 1.

Proof. Use the Euclidean algorithm.

LEMMA 5. Let $m, m_1 \in \mathbb{N}$ satisfy $m > m_1 > 1$ and $gcd(m, m_1) = 1$, and let $m_2, \ldots, m_s, \delta_1, \ldots, \delta_{s-1}$ be defined as in Lemma 4. Then

$$\left(\frac{m}{m_1}\right) = (-1)^{\sum_{i=1}^{s-1} \frac{\delta_i - 1}{2} \cdot \frac{m_i - 1}{2} + \sum_{j=1}^{s-2} \frac{m_j - 1}{2} \cdot \frac{m_{j+1} - 1}{2}},$$

where (m/m_1) is the Jacobi symbol.

Proof. This is clear from the basic properties of the Jacobi symbol.

LEMMA 6. Let $m \in \mathbb{N}$ satisfy m > 1, $m \equiv 1 \pmod{4}$ and suppose m is not a square. Then there exists $m_1 \in \mathbb{N}$ such that $m > m_1 > 1$, $2 \nmid m_1$ and $(m/m_1) = -1$.

Proof. By assumption, $m = p_1 \dots p_r m'^2$, where p_1, \dots, p_r are distinct odd primes and $m' \in \mathbb{N}$ with $2 \nmid m'$. Then there exists a non-residue *a* modulo p_1 . Further, by the Chinese remainder theorem, there exists a $b \in \mathbb{N}$ such that

(9)
$$b \equiv a \pmod{p_1}, \quad b \equiv 1 \pmod{p_j}, \quad j = 2, \dots, r.$$

Let

(10)
$$c = \begin{cases} b & \text{if } \gcd(b, m') = 1, \\ b + p_1 \dots p_r & \text{if } \gcd(b, m') > 1. \end{cases}$$

Since $gcd(b, p_1 \dots p_r) = 1$ by (9), we see from (10) that $c \in \mathbb{Z}$ with gcd(c, m) = 1. Hence, by (9) and (10), we get

(11)
$$\begin{pmatrix} \frac{c}{m} \end{pmatrix} = \begin{pmatrix} \frac{c}{p_1} \end{pmatrix} \dots \begin{pmatrix} \frac{c}{p_r} \end{pmatrix} \begin{pmatrix} \frac{c}{m'^2} \end{pmatrix} = \begin{pmatrix} \frac{c}{p_1} \end{pmatrix} \dots \begin{pmatrix} \frac{c}{p_r} \end{pmatrix} \\ = \begin{pmatrix} \frac{b}{p_1} \end{pmatrix} \dots \begin{pmatrix} \frac{b}{p_r} \end{pmatrix} = \begin{pmatrix} \frac{a}{p_1} \end{pmatrix} = -1.$$

Let $c_0, m_1 \in \mathbb{Z}$ satisfy $c_0 \equiv c \pmod{m}, 0 \leq c_0 < m$ and

$$m_1 = \begin{cases} c_0 & \text{if } 2 \nmid c_0, \\ m - c_0 & \text{if } 2 \mid c_0. \end{cases}$$

Notice that $m \equiv 1 \pmod{4}$ and $2 \nmid m_1$. We see from (11) that $(m/m_1) = -1$. The lemma is proved.

LEMMA 7. The equation

(12) $F(m) = z^2, \quad m, z \in \mathbb{N}, \ m > 1, \ 2 \nmid m \text{ and } m \text{ is not a square,}$

has no solution (m, z).

This is a special case $(M = 1, L = 4u^2D_1)$ of Theorem 3 of [3].

3. Proof of Theorem. The sufficiency being clear, it suffices to prove the necessity. Assume that (1) has solutions (x, y). Then (1) has a unique solution (x_1, y_1) such that

(13)
$$x_1\sqrt{D_1} + y_1^2\sqrt{D_2} \le x\sqrt{D_1} + y^2\sqrt{D_2}$$

where (x, y) runs over all solutions of (1). Notice that (x_1, y_1^2) is a solution of (2) with $x_1, y_1^2 \in \mathbb{N}$. Let (u_1, v_1) be the least solution of (2). By Lemma 1, we have

(14)
$$x_1\sqrt{D_1} + y_1^2\sqrt{D_2} = (u_1\sqrt{D_1} + v_1\sqrt{D_2})^t,$$

where $t \in \mathbb{N}$ with $2 \nmid t$.

If t = 1, then (14) shows that $v_1 = y_1^2$ and the theorem holds. Let $\varepsilon_1, \overline{\varepsilon}_1$ be defined as in (4), and let

$$F_1(m) = \frac{\varepsilon_1^m - \overline{\varepsilon}_1^m}{\varepsilon_1 - \overline{\varepsilon}_1}$$

for any $m \in \mathbb{Z}$ with $2 \nmid m$. If t > 1, then

(15)
$$y_1^2 = \frac{\varepsilon_1^t - \overline{\varepsilon}_1^t}{2\sqrt{D_2}} = v_1 F_1(t),$$

by (14). We deduce from (15) that

(16)
$$v_1 = c_1 y_{11}^2$$

and

(17)
$$F_1(t) = c_1 y_{12}^2$$

where $c_1, y_{11}, y_{12} \in \mathbb{N}$ satisfy $c_1y_{11}y_{12} = y_1$. By Lemma 3(ii), we have $F(t) \equiv t \pmod{v_1}$, hence, by (16) and (17),

(18)
$$t \equiv 0 \pmod{c_1}.$$

We now suppose that t has a divisor p^2 , where p is an odd prime. Let

(19)
$$u_2\sqrt{D_1} + v_2\sqrt{D_2} = (u_1\sqrt{D_1} + v_1\sqrt{D_2})^{t/p}, u_3\sqrt{D_1} + v_3\sqrt{D_2} = (u_1\sqrt{D_1} + v_1\sqrt{D_2})^{t/p^2}$$

By Lemma 1, (u_2, v_2) and (u_3, v_3) are solutions of (2) with $u_2, v_2, u_3, v_3 \in \mathbb{N}$. Further, let

(20)
$$\varepsilon_{2} = u_{2}\sqrt{D_{1}} + v_{2}\sqrt{D_{2}}, \quad \overline{\varepsilon}_{2} = u_{2}\sqrt{D_{1}} - v_{2}\sqrt{D_{2}}, \\ \varepsilon_{3} = u_{3}\sqrt{D_{1}} + v_{3}\sqrt{D_{2}}, \quad \overline{\varepsilon}_{3} = u_{3}\sqrt{D_{1}} - v_{3}\sqrt{D_{2}},$$

and let

(21)
$$F_2(m) = \frac{\varepsilon_2^m - \overline{\varepsilon}_2^m}{\varepsilon_2 - \overline{\varepsilon}_2}, \quad F_3(m) = \frac{\varepsilon_3^m - \overline{\varepsilon}_3^m}{\varepsilon_3 - \overline{\varepsilon}_3},$$

for any $m \in \mathbb{Z}$ with $2 \nmid m$. Then, by (14), we have

$$y_1^2 = \frac{\varepsilon_2^p - \overline{\varepsilon}_2^p}{2\sqrt{D_2}} = v_2 F_2(p).$$

This implies that

(22)
$$v_2 = c_2 y_{11}^{\prime 2}, \quad F_2(p) = c_2 y_{12}^{\prime 2},$$

where $c_2, y'_{11}, y'_{12} \in \mathbb{N}$ satisfy $c_2 y'_{11} y'_{12} = y_1$. By Lemma 3(ii), $F_2(p) \equiv p \pmod{v_2}$, hence, by (22), $p \equiv 0 \pmod{c_2}$. This implies that either $c_2 = 1$ or $c_2 = p$. From (22), if $c_2 = 1$, then $F_2(p)$ is a square, which is impossible by Lemma 7. Therefore, $c_2 = p$ and

(23)
$$v_2 = py_{11}^{\prime 2}, \quad F_2(p) = py_{12}^{\prime 2},$$

by (22). On the other hand, we see from (19)-(21) that

(24)
$$v_2 = v_3 F_3(p).$$

The combination of (23) and (24) yields

(25)
$$v_3 = \begin{cases} c_3 y_{111}^2, \\ c_3 p y_{112}^2, \end{cases} \quad F_3(p) = \begin{cases} c_3 p y_{112}^2, \\ c_3 y_{111}^2, \end{cases}$$

where $c_3, y_{111}, y_{112} \in \mathbb{N}$ satisfy $c_3y_{111}y_{112} = y'_{11}$. Notice that $F_3(p)$ is never a square by Lemma 7. By much the same argument as above, we can find from (25) that $c_3 = 1$ or p, and v_3 is a square. Since (u_3, v_3) is a solution of (2), it follows that $(u_3, \sqrt{v_3})$ is a solution of (1) satisfying $u_3\sqrt{D_1} + v_3\sqrt{D_2} < x_1\sqrt{D_1} + y_1^2\sqrt{D_2}$ by (19), which contradicts our assumption (13). Thus, t is square free and so is c_1 by (18).

If $t \neq c_1$, then t has an odd prime divisor q with $q \nmid c_1$ by (18). Let

(26)
$$u_4\sqrt{D_1} + v_4\sqrt{D_2} = (u_1\sqrt{D_1} + v_1\sqrt{D_2})^{t/q},$$

(27)
$$\varepsilon_4 = u_4 \sqrt{D_1 + v_4} \sqrt{D_2}, \quad \overline{\varepsilon}_4 = u_4 \sqrt{D_1 - v_4} \sqrt{D_2},$$

and let

(28)
$$F_4(m) = \frac{\varepsilon_4^m - \overline{\varepsilon}_4^m}{\varepsilon_4 - \overline{\varepsilon}_4}$$

for any $m \in \mathbb{Z}$ with $2 \nmid m$. Then, by (14) and (26)–(28), we have $y_1^2 = v_4 F_4(q)$, whence

(29)
$$v_4 = c_4 y_{13}^2, \quad F_4(q) = c_4 y_{14}^2,$$

where $c_4, y_{13}, y_{14} \in \mathbb{N}$ satisfy $c_4y_{13}y_{14} = y_1$. Using the same method, by (26) and (29), we can prove that $c_4 = q$ and

(30)
$$v_4 = qy_{13}^2 = \frac{\varepsilon_1^{t/q} - \overline{\varepsilon}_1^{t/q}}{2\sqrt{D_2}} = v_1 F_1(t/q).$$

Substituting (16) into (30) gives

(31)
$$F_1\left(\frac{t}{q}\right) = \frac{qy_{13}^2}{c_1y_{11}^2}.$$

Notice that $F_1(t/q) \in \mathbb{N}$ and $q \nmid c_1$. We see from (31) that $q \mid y_{11}, q \mid y_{13}, q \mid F_1(t/q)$ and $q \mid v_1$ by (16), a contradiction. Thus, we deduce that $t = c_1$ and the necessity is proved by (14) and (16), since t is square free. The proof is complete.

R e m a r k. By much the same argument as in the proof of the Theorem, we can prove a similar result for the case $\min(D_1, D_2) = 1$.

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Reçu par la Rédaction le 26.4.1993