

Sparse polynomial exponential sums

by

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1. Introduction. In this paper we estimate the complete exponential sum

$$(1.1) \quad S(f, q) = \sum_{x=1}^q e_q(f(x)),$$

where $e_q(\cdot)$ is the additive character $e_q(\cdot) = e^{2\pi i \cdot / q}$, and f is a sparse integer polynomial,

$$(1.2) \quad f(x) = a_1 x^{k_1} + \dots + a_r x^{k_r}$$

with $0 < k_1 < \dots < k_r$. We always assume that the content of f , (a_1, \dots, a_r) , is relatively prime to the modulus q . Let $d = d(f) = k_r$ denote the degree of f and for any prime p let $d_p(f)$ denote the degree of f read modulo p . A fundamental problem is to determine whether there exists an absolute constant C such that for an arbitrary positive integer q ,

$$(1.3) \quad |S(f, q)| \leq C q^{1-1/d},$$

if f is not a constant function modulo p for each prime $p | q$. It is well known that the exponent $1 - 1/d$ is best possible. For the case of Gauss sums ($r = 1$) Shparlinski [26], [27] showed that one may take $C = 1 + O(d^{-1/4+\varepsilon})$ and this was sharpened to $C = 1 + O(d^{-1+\varepsilon})$ in his subsequent work with Konyagin [14, Theorem 6.7].

The best upper bounds available for general f are

$$|S(f, q)| \leq e^{d+O(d/\log d)} q^{1-1/d},$$

due to Stechkin [29], and

$$|S(f, q)| \leq e^{1.74d} q^{1-1/d},$$

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due to Qi and Ding [25]; see also Chen [2], [3], Hua [11]–[13], Lu [17]–[19], Nechaev [20], [21], Nechaev and Topunov [22], Qi and Ding [23], [24] and Zhang and Hong [31]. These authors noted that in order to make any further improvement one must first obtain a nontrivial upper bound on the prime modulus exponential sum $|S(f, p)|$ for $p < (d-1)^2$, the interval where Weil's [30] bound $|S(f, p)| \leq (d-1)\sqrt{p}$ is worse than the trivial bound. In [5] we obtained a bound of this type in terms of the number of terms r of $f(x)$. Using this bound we establish here

THEOREM 1.1. *For any positive integer r there exists a constant $C(r)$ such that for any polynomial f of type (1.2) and positive integer q relatively prime to the content of f ,*

$$|S(f, q)| \leq C(r)q^{1-1/d}.$$

Although our proof yields $C(r) \leq e^{O(r^4)}$, no attempt was made to obtain the best possible value for $C(r)$.

For prime power moduli one can replace $C(r)$ with an absolute constant as shown by Stechkin [29] and Cochrane and Zheng [8], the latter result being

LEMMA 1.1 [8, Theorem 1.1]. *Let f be a polynomial over \mathbb{Z} of degree d and p a prime with $d_p(f) \geq 1$. Then for any $m \geq 1$,*

$$(1.4) \quad |S(f, p^m)| \leq 4.41p^{m(1-1/d)}.$$

It is also well known (see [20], [3] or [8]) that for $p \geq (d-1)^{2d/(d-2)}$ and $m \geq 1$,

$$(1.5) \quad |S(f, p^m)| \leq p^{m(1-1/d)}.$$

The significance of the constant one in (1.5) lies in the fact that bounds for exponential sums modulo prime powers lead to bounds for a general modulus $q = \prod_{i=1}^k p_i^{e_i}$ via the multiplicative formula

$$(1.6) \quad S(f, q) = \prod_{i=1}^k S(\lambda_i f, p_i^{e_i}),$$

where the λ_i are such that $\sum_{i=1}^k \lambda_i q/p_i^{e_i} = 1$. Thus if (1.5) holds for all prime power divisors of q then it follows that $|S(f, q)| \leq q^{1-1/d}$. It is desirable to extend the inequality in (1.5) to an interval of the type $p > Cd$ for some constant C .

In closing we note that for sums over reduced residue systems,

$$(1.7) \quad S^*(f, q) = \sum_{x=1, (x,q)=1}^q e_q(f(x)),$$

the exponent in the upper bound can be dramatically reduced. Shparlinski [28] showed that

$$|S^*(f, q)| \leq C(d, \varepsilon)q^{1-1/r+\varepsilon},$$

for any sparse polynomial in r terms with content relatively prime to q . Loh [16] obtained a related upper bound but an error in his Lemma 3 leaves his results in doubt.

2. The method of recursion. A standard method for bounding exponential sums modulo prime powers is the method of recursion, also known as the method of critical points. For any polynomial f let $t = t_p(f) = \text{ord}_p(f')$ be the largest power of p dividing all of the coefficients of f' , $d_1 = d_p(p^{-t}f')$, and let $\mathcal{A} = \mathcal{A}(f, p)$ be the set of zeros of the congruence $p^{-t}f'(x) \equiv 0 \pmod{p}$. \mathcal{A} is called the *set of critical points* associated with the sum $S(f, p^m)$, for any $m \geq 2$. Write

$$S(f, p^m) = \sum_{\alpha=0}^{p-1} S_\alpha(f, p^m)$$

with

$$S_\alpha(f, p^m) = \sum_{x \equiv \alpha \pmod{p}} e_{p^m}(f(x)).$$

A fact of central importance is that if m is sufficiently large then $S_\alpha(f, p^m) = 0$ unless α is a critical point.

LEMMA 2.1 [6, Proposition 4.1]. *Suppose that p is an odd prime and $m \geq t + 2$, or $p = 2$ and $m \geq t + 3$, or $p = 2$, $t = 0$ and $m = 2$. Then if α is not a critical point, $S_\alpha(f, p^m) = 0$. Consequently,*

$$S(f, p^m) = \sum_{\alpha \in \mathcal{A}} S_\alpha(f, p^m).$$

For any $\alpha \in \mathcal{A}$ define

$$(2.1) \quad \begin{aligned} \sigma &= \sigma_\alpha := \text{ord}_p(f(px + \alpha) - f(\alpha)), \\ g_\alpha(x) &:= p^{-\sigma}(f(px + \alpha) - f(\alpha)). \end{aligned}$$

LEMMA 2.2 [6, Proposition 4.1] (The recursion relationship). *Suppose that p is an odd prime and $m \geq t + 2$, or $p = 2$ and $m \geq t + 3$, or $p = 2$, $t = 0$ and $m = 2$. Then if $\alpha \in \mathcal{A}$,*

$$(2.2) \quad S_\alpha(f, p^m) = e_{p^m}(f(\alpha))p^{\sigma-1}S(g_\alpha, p^{m-\sigma}),$$

where

$$(2.3) \quad S(g_\alpha, p^{m-\sigma}) = \begin{cases} \sum_{x=1}^{p^{m-\sigma}} e_{p^{m-\sigma}}(g_\alpha(x)) & \text{if } m > \sigma, \\ p^{m-\sigma} & \text{if } m \leq \sigma. \end{cases}$$

Under the hypotheses of the lemma we have

$$(2.4) \quad |S(f, p^m)| \leq \sum_{\alpha \in \mathcal{A}} |S_\alpha(f, p^m)| = \sum_{\alpha \in \mathcal{A}} p^{\sigma_\alpha - 1} |S(g_\alpha, p^{m - \sigma_\alpha})|.$$

In particular, since there are at most d_1 critical points we immediately have the upper bound

$$(2.5) \quad |S(f, p^m)| \leq d_1 p^{m-1}.$$

In [8] we established the following bounds for $S_\alpha(f, p^m)$ and $S(f, p^m)$:

LEMMA 2.3 [8, Theorem 2.1]. *Let f be a polynomial over \mathbb{Z} and p a prime with $d_p(f) \geq 1$. Suppose that p is odd and $m \geq t + 2$, or $p = 2$ and $m \geq t + 3$. Set $\lambda = (5/4)^5 \approx 3.05$ and $d_1 = d_p(p^{-t}f')$. Then*

(i) *For any critical point α of multiplicity ν we have*

$$(2.6) \quad |S_\alpha(f, p^m)| \leq \min\{\nu, \lambda\} p^{t/(\nu+1)} p^{m(1-1/(\nu+1))},$$

with equality if $\nu = 1$.

(ii) $|S(f, p^m)| \leq \lambda p^{t/(d_1+1)} p^{m(1-1/(d_1+1))}$.

Related results using the method of critical points were obtained by Chalk [1], Cochrane [4], Cochrane and Zheng [6], [7], Ding [9], [10], and Loh [15].

For any critical point α set

$$(2.7) \quad \tau := \text{ord}_p(g'_\alpha(x)), \quad g_1(x) := p^{-\tau} g'_\alpha(x).$$

The following relations are well known (see e.g. [6, Lemma 3.1]) and play a central role in the proof of the preceding lemma.

LEMMA 2.4.

$$(2.8) \quad \sigma \geq \begin{cases} t + 2 & \text{if } p \text{ is odd or } \nu > 1, \\ t + 1 & \text{if } p = 2 \text{ and } \nu = 1. \end{cases}$$

$$(2.9) \quad \sigma \leq \nu + 1 + t - \tau.$$

$$(2.10) \quad d_p(g_\alpha) \leq \begin{cases} \sigma - t + \text{ord}_p(d_p(g_\alpha)) \leq \nu + 1 + \text{ord}_p(d_p(g_\alpha)), \\ \sigma \leq \nu + 1 + t - \tau. \end{cases}$$

$$(2.11) \quad d_p(g_1) \leq \sigma + \tau - t - 1 \leq \nu.$$

$$(2.12) \quad p^\tau \mid d_p(g_\alpha).$$

An immediate consequence that we frequently make reference to is

LEMMA 2.5. *Suppose that α is a critical point of multiplicity ν with $\nu \geq 2$ and $p > \nu + 2$. Then $d_p(g_\alpha) \leq \nu + 1$.*

Proof. Let $d_p = d_p(g_\alpha)$. Suppose that $\text{ord}_p(d_p) \geq 1$. If $d_p = p$ then by (2.10) we have $p = d_p \leq \nu + 2$ contradicting our assumption. Otherwise $d_p \geq 2p$ and we have $p \leq d_p/2 \leq d_p - \text{ord}_p(d_p) \leq \nu + 1$, again a contradiction. Thus $p \nmid d_p$ and we obtain (by (2.10)) $d_p \leq \nu + 1$. ■

3. Preliminary upper bounds. We begin with a couple of auxiliary lemmas.

LEMMA 3.1. *Define $\lambda_i = i$ for $i = 1, 2, 3$ and $\lambda_i = \lambda$ for $i \geq 4$, where $\lambda = (5/4)^5 \approx 3.05$. Then for $1 \leq i \leq d$ we have*

$$d\lambda_i\lambda^{(i-d)/(i+1)} \leq i\lambda.$$

Proof. For any fixed $i \geq 1$ the function $f_i(x) := (\lambda_i/i)x\lambda^{(i-x)/(i+1)}$ attains its maximum value at $x = (i+1)/\log(\lambda) < i+1$, and is decreasing for larger values of x . Thus for $d \geq i$, the maximum value of $f_i(d)$ occurs at $d = i$ or $d = i+1$. Now, $f_i(i) = \lambda_i \leq \lambda$ and $f_i(i+1) = \lambda_i(1+1/i)\lambda^{-1/(i+1)} \leq \lambda$, as can be seen by considering the different cases $i = 1, 2, 3$ and $i \geq 4$. ■

LEMMA 3.2. *If $p > cd_1$ for some constant c then for $1 \leq i \leq d_1 - 1$ we have*

$$(4p/(cd_1))^{(i-d_1)/(i+1)} \leq i/d_1.$$

Proof. We first note that

$$(d_1/i)^{(i+1)/(d_1-i)} \leq 4 \quad \text{for } 1 \leq i \leq d_1 - 1.$$

This can be checked directly for $i = 1, 2, 3$. For $i \geq 4$ it follows from Lemma 3.1. Then $p > cd_1 \geq (c/4)d_1(d_1/i)^{(i+1)/(d_1-i)}$ and the result follows. ■

LEMMA 3.3. *Let p be a prime and f be any integer polynomial with $t = 0$ and either $d_1 = 0, 1$ or $p > d_1^{2+4/(d_1-1)}$ where $d_1 = d_p(p^{-t}f')$. Then for $m \geq 2$,*

$$(3.1) \quad |S(f, p^m)| \leq p^{m(1-1/(d_1+1))}.$$

Proof. If $d_1 = 0$ then there are no critical points and the sum is zero. If $d_1 = 1$ then there is a single critical point of multiplicity one and the result follows from Lemma 2.3(i). Suppose that $d_1 \geq 2$. Let $\mathcal{A} = \mathcal{A}(f, p) \subset \mathbb{F}_p$ be the set of critical points. We prove by induction on m that, under the hypotheses of the theorem,

$$(3.2) \quad |S_\alpha| \leq p^{m(1-1/(\nu+1))}$$

for any critical point $\alpha \in \mathcal{A}$. We first note that (3.1) is an immediate consequence of (3.2). Indeed, if $p^m \leq (p/d_1)^{d_1+1}$ then using the trivial upper bound $|S_\alpha(f, p^m)| \leq p^{m-1}$ we have $|S(f, p^m)| \leq \sum_{\alpha \in \mathcal{A}} |S_\alpha(f, p^m)| \leq d_1 p^{m-1} \leq p^{m(1-1/(d_1+1))}$. Next, if there is a critical point α of multiplicity d_1 then it is the only critical point and we have $|S(f, p^m)| = |S_\alpha(f, p^m)| \leq p^{m(1-1/(d_1+1))}$.

Finally, suppose that $p^m > (p/d_1)^{d_1+1}$ and that every critical point is of multiplicity less than d_1 . Letting n_i denote the number of critical points of

multiplicity i we deduce from (3.2) that

$$(3.3) \quad |S(f, p^m)| \leq \sum_{\alpha \in \mathcal{A}} |S_\alpha(f, p^m)| \leq \sum_{i=1}^{d_1-1} n_i p^{m(1-1/(i+1))} \\ = p^{m(1-1/(d_1+1))} \left(\sum_{i=1}^{d_1-1} n_i p^{m(i-d_1)/((i+1)(d_1+1))} \right).$$

Then from $p^m > (p/d_1)^{d_1+1}$, $p > 4d_1$ and Lemma 3.2 with $c = 4$ we obtain

$$(3.4) \quad |S(f, p^m)| \leq p^{m(1-1/(d_1+1))} \left(\sum_{i=1}^{d_1-1} n_i (p/d_1)^{(i-d_1)/(i+1)} \right) \\ \leq \left(\sum_{i=1}^{d_1-1} n_i i / d_1 \right) p^{m(1-1/(d_1+1))} \leq p^{m(1-1/(d_1+1))}.$$

We now proceed to establish (3.2). If $\nu = 1$ then by Lemma 2.3 we have equality in (3.2). So we may assume that $\nu \geq 2$. When $m = 2$ the bound is trivial, $|S_\alpha| \leq p \leq p^{2(1-1/(\nu+1))}$. Suppose $m \geq 3$. If $\sigma \geq m$ then the result follows trivially,

$$|S_\alpha| \leq p^{m-1} \leq p^{m(1-1/(\nu+1))} p^{(\sigma-\nu-1)/(\nu+1)} \leq p^{m(1-1/(\nu+1))},$$

the latter inequality following from (2.9). Suppose next that $\sigma = m - 1$. Put $d_p = d_p(g_\alpha)$. Since $p > d_1^2 \geq \nu^2 \geq \nu + 2$ it follows from Lemma 2.5 that $d_p \leq \nu + 1 \leq d_1 + 1$. If $d_p \geq 3$ then $p \geq (d_p - 1)^{2+4/(d_p-2)}$, so by the Weil bound, $|S(g_\alpha, p)| \leq (d_p - 1)\sqrt{p} \leq p^{1-1/d_p} \leq p^{1-1/(\nu+1)}$. If $d_p = 1$ or 2 the same bound is elementary. It follows from the recursion formula of Lemma 2.2 that

$$|S_\alpha| = p^{\sigma-1} |S(g_\alpha, p)| \\ \leq p^{m-1-1/(\nu+1)} = p^{(\sigma-\nu-1)/(\nu+1)} p^{m(1-1/(\nu+1))} \leq p^{m(1-1/(\nu+1))}.$$

Suppose finally that $m \geq \sigma + 2$. We note that $\tau = 0$ since by (2.12), $p^\tau \leq d_p(g_\alpha) \leq \nu + 1 \leq d_1 + 1 < p$, and so we can apply the induction assumption to $S(g_\alpha, p^{m-\sigma})$. Putting $d_2 = d_p(g'_\alpha) \leq \nu \leq d_1$ and noting that either $d_2 = 0, 1$ or $p \geq d_2^{2+2/(d_2-1)}$ we obtain

$$|S_\alpha| = p^{\sigma-1} |S(g_\alpha, p^{m-\sigma})| \leq p^{\sigma-1} p^{(m-\sigma)(1-1/(d_2+1))} \\ \leq p^{\sigma-1} p^{(m-\sigma)(1-1/(\nu+1))} \leq p^{m(1-1/(\nu+1))}. \blacksquare$$

4. Multiplicity estimates. Next, we obtain an upper bound on the multiplicity of a nonzero zero of a sparse polynomial

$$f(x) = a_1 x^{k_1} + \dots + a_r x^{k_r} \pmod{p}.$$

Let $a \not\equiv 0 \pmod{p}$ be a zero of multiplicity $\nu \pmod{p}$, that is,

$$(x - a)^\nu \parallel f(x) \pmod{p}.$$

For $1 \leq i \leq r$ let

$$S(i, \alpha) = \{k_j : k_j \equiv k_i \pmod{p^\alpha}\},$$

and set

$$(4.1) \quad \alpha_i = \max\{\alpha : |S(i, \alpha)| \geq 2\},$$

$$(4.2) \quad r_i = |S(i, \alpha_i)|.$$

LEMMA 4.1. *The multiplicity ν of any nonzero zero of $f(x) \pmod{p}$ satisfies $\nu < \min_i r_i p^{\alpha_i}$. In particular, if p does not divide any $k_i - k_j$ with $i \neq j$ then $\nu < r$.*

Lemma 4.1 follows from the more precise

LEMMA 4.2. *Suppose that k_1, \dots, k_t are the smallest distinct exponents modulo p so that*

$$f(x) = x^{k_1} f_1(x)^p + \dots + x^{k_t} f_t(x)^p \pmod{p},$$

where

$$f_i(x) = \sum_{k_j = k_i + l_j p} a_j x^{l_j}.$$

Then if $f(x)$ has a nonzero zero a of multiplicity $\nu \pmod{p}$, we have

$$\nu = kp + u$$

where $u < t$ and $(x - a)^k$ is the highest power dividing all the f_1, \dots, f_t .

Proof. Suppose that $(x - a)^k \mid f_1, \dots, f_t$ with $(x - a)^{k+1} \nmid f_1$, and write $f_i(x) = (x - a)^k g_i(x) \pmod{p}$, $\nu = kp + u$, so that

$$(x - a)^u \parallel g(x) = x^{k_1} g_1(x)^p + \dots + x^{k_t} g_t(x)^p.$$

Writing $\nabla = x \frac{d}{dx}$ we must have $\nabla^j g(a) \equiv 0 \pmod{p}$ for $j = 0, \dots, u - 1$. That is,

$$k_1^j a^{k_1} g_1(a)^p + \dots + k_t^j a^{k_t} g_t(a)^p \equiv 0 \pmod{p}$$

for $j = 0, \dots, u - 1$. Since $|\det(k_i^j)_{i=1, \dots, t, j=0, \dots, t-1}| = \prod |k_i - k_j| \not\equiv 0 \pmod{p}$ and $a^{k_1} g_1(a)^p \not\equiv 0 \pmod{p}$ we must therefore have $u < t$. ■

Proof of Lemma 4.1. Pick an arbitrary k_i , $i = 1, \dots, t$, and use the preceding lemma and induction on α_i : If $\alpha_i = 0$ then plainly $k = 0$ and $\nu = u < t \leq r = r_i$. If $\alpha_i \geq 1$ then since $(x - a)^k \mid f_i(x)$ we have (by induction) $k < r_i p^{\alpha_i - 1}$ and $u < p$ giving

$$\nu = pk + u \leq (r_i p^{\alpha_i - 1} - 1)p + (p - 1) < r_i p^{\alpha_i}. \quad \blacksquare$$

In practice we apply the multiplicity estimate to the polynomial $p^{-t} f'(x)$ and so we let $r_1 = r_1(f, p)$ be the number of nonzero terms modulo p of the polynomial $p^{-t} f'(x)$. For critical points having multiplicity less than r_1 we have the following upper bound.

LEMMA 4.3. *Let f be a sparse polynomial as in (1.2) and suppose that either $r_1 = 1, 2$ or $p > (r_1 - 1)^{2r_1/(r_1-2)}$. Then if $m \geq t + 2$ and α is a critical point of multiplicity $\nu < r_1$ we have*

$$(4.3) \quad |S_\alpha| \leq p^{t/(\nu+1)} p^{m(1-1/(\nu+1))}.$$

Proof. If $\nu = 1$ the result follows from Lemma 2.3, and so we may assume $\nu \geq 2$ and $r_1 \geq 3$. Let $d_p = d_p(g_\alpha)$. Since $p > \nu^{2+4/(\nu-1)}$, we get $d_p \leq \nu + 1$ by Lemma 2.5, and thus $p > (d_p - 1)^{2d_p/(d_p-2)}$. Also, since $p^\tau \leq d_p(g_\alpha) \leq \nu + 1 \leq r_1 + 1 < p$ we must have $\tau = 0$.

If $\sigma \geq m$ the result follows trivially,

$$|S_\alpha| \leq p^{m-1} = p^{(m-\nu-1)/(\nu+1)} p^{m(1-1/(\nu+1))} \leq p^{t/(\nu+1)} p^{m(1-1/(\nu+1))}.$$

Suppose next that $\sigma = m - 1$. Then applying the bound in (1.5) to $S(g_\alpha, p)$, we obtain

$$\begin{aligned} |S_\alpha| &= p^{\sigma-1} |S(g_\alpha, p)| \leq p^{\sigma-1} p^{1-1/(\nu+1)} \\ &= p^{(\sigma-\nu-1)/(\nu+1)} p^{m(1-1/(\nu+1))} \leq p^{t/(\nu+1)} p^{m(1-1/(\nu+1))}. \end{aligned}$$

Finally, if $\sigma \leq m - 2$ then we can apply Lemma 3.3 to $S(g_\alpha, p^{m-\sigma})$, since $d_2 := d_p(g'_\alpha) \leq \nu < r_1$ and so $p \geq d_2^{2+4/(d_2-1)}$. We obtain

$$\begin{aligned} |S_\alpha| &\leq p^{\sigma-1} |S(g_\alpha, p^{m-\sigma})| \\ &\leq p^{\sigma-1} p^{(m-\sigma)(1-1/(d_2+1))} = p^{(\sigma-\nu-1)/(\nu+1)} p^{m(1-1/(\nu+1))}. \quad \blacksquare \end{aligned}$$

5. Bounds for exponential sums with p small relative to d . First we consider sums modulo p . From the bound of Weil, one deduces (see [8, Lemma 3.1]) the upper bound

$$(5.1) \quad |S(f, p)| \leq 1.75p^{1-1/d}$$

for any polynomial f with $d_p(f) \geq 1$. Moreover the constant 1.75 may be replaced by 1 provided $p \gg d^2$. For our purposes here we need the constant 1 for $p \gg d$. We obtain this from the following result established in the authors' work [5, Corollary 1.1].

LEMMA 5.1. *Let f be an integer polynomial of degree d as in (1.2). Then for any $\delta > 0$, if $p > (9/\delta^{1.06})d$ and $p > C_1(\delta)$, then*

$$(5.2) \quad \left| \sum_{x=1}^p e_p(f(x)) \right| \leq p \left(1 - \frac{1}{rp^\delta} \right).$$

LEMMA 5.2. *Let f be a polynomial as in (1.2) of degree $d = d_p(f) \geq 1 \pmod{p}$ and suppose that $p > C_2$ (an absolute constant), $p > 50d$ and $p > r^4$. Then*

$$|S(f, p)| \leq p^{1-1/d}.$$

Proof. The result is elementary for $d = 1, 2$ and so we assume $d > 2$. If $p > 16d^2$ then the result follows from the Weil bound $|S(f, p)| \leq (d-1)\sqrt{p}$.

Suppose that $p \leq 16d^2$. Applying Lemma 5.1 with $\delta = 1/5$ we deduce that if $p > 50d$ and $p > C_1(1/5)$ then $|S(f, p)| \leq p(1 - 1/(rp^{1/5}))$. Since $p > r^4$ it follows that $|S(f, p)| \leq p(1 - 1/p^{9/20})$, and since $p \leq 16d^2$ the latter is $\leq p^{1-1/d}$ for $p > 10^{60}$. ■

LEMMA 5.3. *Let f be a sparse polynomial as in (1.2) with $p \geq 50(d_1+1)$, $p > C_2$ (the constant in Lemma 5.2), $p > r^4$ and $m \geq t + 2$. Then for any critical point α of multiplicity ν we have*

$$(5.3) \quad |S_\alpha| \leq p^{t/(\nu+1)} p^{m(1-1/(\nu+1))},$$

and

$$(5.4) \quad |S(f, p^m)| \leq p^{t/(d_1+1)} p^{m(1-1/(d_1+1))}.$$

Proof. We first observe that (5.4) is always an immediate consequence of (5.3). Indeed, if $p^{m-t} \leq (p/d_1)^{d_1+1}$ then using the trivial upper bound $|S_\alpha(f, p^m)| \leq p^{m-1}$ we have $|S(f, p^m)| \leq \sum_{\alpha \in \mathcal{A}} |S_\alpha(f, p^m)| \leq d_1 p^{m-1} \leq p^{t/(d_1+1)} p^{m(1-1/(d_1+1))}$. Next, if there is a critical point α of multiplicity d_1 then it is the only critical point and we have $|S(f, p^m)| = |S_\alpha(f, p^m)| \leq p^{t/(d_1+1)} p^{m(1-1/(d_1+1))}$.

Finally, suppose that $p^{m-t} > (p/d_1)^{d_1+1}$ and that every critical point is of multiplicity less than d_1 . Letting n_i denote the number of critical points of multiplicity i we deduce from (5.3) that

$$\begin{aligned} |S(f, p^m)| &\leq \sum_{i=1}^{d_1} n_i p^{t/(i+1)} p^{m(1-1/(i+1))} \\ &\leq p^{t/(d_1+1)} p^{m(1-1/(d_1+1))} \sum_i n_i p^{(m-t)(i-d_1)/((i+1)(d_1+i))} \\ &\leq p^{t/(d_1+1)} p^{m(1-1/(d_1+1))} \sum_i n_i (p/d_1)^{(i-d_1)/(i+1)} \\ &\leq p^{t/(d_1+1)} p^{m(1-1/(d_1+1))}, \end{aligned}$$

the last inequality following from Lemma 3.2 (with $c = 4$) and $\sum_i n_i i \leq d_1$.

We now establish (5.3) by induction on m . The result is trivial if $m = 2$. Suppose that $m > 2$. If $\sigma \geq m$ then from (2.9),

$$|S_\alpha| \leq p^{m-1} \leq p^{t/(\nu+1)} p^{m(1-1/(\nu+1))}.$$

If $\sigma = m - 1$ and $\alpha \neq 0$ then since $p > d_1$ it follows from Lemma 4.1 that $\nu < r$. Also, since $p \geq 50d_1 \geq 50\nu$ we see by Lemma 2.5 that

$$d_p(g) \leq \nu + 1 \leq r < p^{1/4},$$

and so by (1.5), $|S(g_\alpha, p)| \leq p^{1-1/d_p(g)} \leq p^{1-1/(\nu+1)}$. It then follows from the recursion relation that

$$(5.5) \quad |S_\alpha| \leq p^{\sigma-1} |S(g_\alpha, p)| \leq p^{\sigma-1/(\nu+1)} \leq p^{t/(\nu+1)} p^{m(1-1/(\nu+1))},$$

by (2.9). If $\alpha = 0$ then we have to argue differently since the multiplicity may be larger than r . In this case $g_\alpha(x) = f(px)$ is a sparse polynomial with the same number of terms as f . Since $p > 50(d_1 + 1) \geq 50(\nu + 1) \geq 50d_p(g_\alpha)$ we can apply Lemma 5.2 to obtain $|S(g_\alpha, p)| \leq p^{1-1/d_p(g_\alpha)}$, and the result follows as before.

Suppose now that $\sigma \leq m - 2$. We first note that by (2.12), $\tau = 0$ since $p > d_p(g_\alpha)$. Set $d_2 = d_p(g'_\alpha)$. If $\alpha \neq 0$ then by (2.11) and Lemma 4.1, $d_2 \leq \nu < r < p^{1/4}$. Thus by Lemma 3.3,

$$|S(g_\alpha, p^{m-\sigma})| \leq p^{(m-\sigma)(1-1/(d_2+1))}.$$

If $\alpha = 0$ then we can apply the induction assumption to the polynomial $g_\alpha = p^{-\sigma}f(px)$ and obtain the same bound. From the recursion relationship we then obtain

$$\begin{aligned} |S(f, p^m)| &\leq p^{\sigma-1}p^{(m-\sigma)(1-1/(d_2+1))} \\ &\leq p^{-1+\sigma/(\nu+1)}p^{m(1-1/(\nu+1))} \leq p^{t/(\nu+1)}p^{m(1-1/(\nu+1))}. \blacksquare \end{aligned}$$

Next we obtain a bound valid for even smaller values of p . Again, let d_1 and $r_1 = r_1(f, p)$ be the degree and number of nonzero terms of the polynomial $p^{-t}f'(x)$ read modulo p .

LEMMA 5.4. *Let f be a sparse polynomial in r terms and p a prime with $p > r^4$, $p > C_3$ and such that $p \nmid (k_j - k_i)$ for all $k_i < k_j \leq d_1$. Then for $m \geq t + 2$ and any critical point α of multiplicity ν we have*

- (i) *If $\alpha \neq 0$ then $|S_\alpha(f, p^m)| \leq p^{t/(\nu+1)}p^{m(1-1/(\nu+1))}$.*
- (ii) *For $\alpha = 0$, $|S_0(f, p^m)| \leq p^{(2r+t)/(\nu+1)}p^{m(1-1/(\nu+1))}$.*
- (iii) $|S(f, p^m)| \leq p^{(2r+t)/(d_1+1)}p^{m(1-1/(d_1+1))}$.

Proof. We take $C_3 = \max\{C_2, 200\}$ where C_2 is the constant in Lemma 5.3. The condition $p \nmid (k_i - k_j)$ implies (by Lemma 4.1) that $\nu < r_1$ for any nonzero critical point. So (i) is implied by Lemma 4.3. If $p \geq 50(d_1 + 1)$ then the lemma is implied by Lemma 5.3 and so we may assume $p < 50(d_1 + 1)$. In particular, it follows that $r \leq d_1$ (if $r \geq 4$ then $r^4 < p < 50(d_1 + 1)$ implies $r < r \cdot r^3/50 < d_1 + 1$; if $r \leq 3$ then $p > 200$ implies $d_1 > 3 \geq r$).

The proof of (ii) is by induction on m , but first we show that (i) and (ii) together imply (iii). If zero is the only critical point then (ii) immediately implies (iii) and so we assume henceforth that $r \geq 2$ and that $\nu(0) < d_1$.

If $m - t \leq 2r$ then the upper bound in (iii) follows from the trivial bound $|S(f, p^m)| \leq p^m$. Next write $m - t = 2r + 1 + j$ with $j \geq 0$ and set

$$\Delta = p^{(t+2r)/(d_1+1)}p^{m(1-1/(d_1+1))},$$

the desired bound. We have

$$|S(f, p^m)| \leq |S_0(f, p^m)| + \sum_{\alpha \neq 0} |S_\alpha(f, p^m)|.$$

For the first term we have the trivial bound

$$(5.6) \quad |S_0(f, p^m)| \leq p^{m-1} = p^{(j-d_1)/(d_1+1)} \Delta.$$

Now there are at most $p - 1$ nonzero critical points, each of multiplicity $\leq r_1 - 1 \leq r - 1$, and so by (i),

$$(5.7) \quad \sum_{\alpha \neq 0} |S_\alpha(f, p^m)| \leq p \cdot p^{t/r} p^{m(1-1/r)} \\ = p^{(j(r-d_1-1)-rd_1-d_1-1)/((d_1+1)r)} \Delta = p^{(j-d_1)/(d_1+1)-(1+j)/r} \Delta.$$

Combining (5.6) and (5.7) we have, for $j \leq d_1/2$,

$$|S(f, p^m)| \leq p^{-d_1/(2(d_1+1))} (1 + p^{-1/r}) \Delta < 2p^{-1/4} \Delta < \Delta,$$

and for $d_1 > j > d_1/2$,

$$|S(f, p^m)| \leq (p^{-d_1/(2r)} + 1) p^{-1/(d_1+1)} \Delta \leq (r^{-2d_1/r} + 1) r^{-4/(d_1+1)} \Delta < \Delta.$$

If $j \geq d_1$ then by the bound in (ii) (replacing ν with $d_1 - 1$) we obtain

$$|S_0(f, p^m)| \leq p^{(-j-1)/(d_1(d_1+1))} \Delta \leq p^{-1/d_1} \Delta.$$

For the remaining critical points we use the upper bound of (5.7) replacing j with d_1 . Thus

$$|S(f, p^m)| \leq (p^{-1/d_1} + p^{(-d_1-1)/r}) \Delta \leq (r^{-4/d_1} + r^{-4(d_1+1)/r}) \Delta < \Delta.$$

We return to the task of proving (ii) by induction on m . The bound follows trivially from $|S_0(f, p^m)| \leq p^{m-1}$ if $m \leq \nu + 1 + t + 2r$, and so we assume $m \geq \nu + 2 + t + 2r$. By (2.9) we have

$$m - \sigma \geq \nu + 2 + t + 2r - (\nu + 1 + t - \tau) = 1 + 2r + \tau \geq \tau + 2,$$

and by the recursion formula of Lemma 2.2, $|S_0(f, p^m)| = p^{\sigma-1} |S(g_0, p^{m-\sigma})|$, where $g_0(x) = p^{-\sigma} f(px)$. Since g_0 has the same degree monomials as f we can apply the induction assumption to g_0 and obtain,

$$|S_0(f, p^m)| \leq p^{\sigma-1} p^{(\tau+2r)/(d_2+1)} p^{(m-\sigma)(1-1/(d_2+1))},$$

where $d_2 := d_p(p^{-\tau} g_0) \leq d_1$. Now by (2.11), $d_2 \leq \nu$ and so replacing d_2 by ν in the previous inequality and using the upper bound in (2.9) we deduce the inequality in (ii). ■

6. Dealing with the primes that divide $k_i - k_j$ for some $i \neq j$. If $p | (k_j - k_i)$ for some $k_i < k_j \leq d_1$ then there may be nonzero critical points of multiplicity exceeding r and so we have to argue more carefully. Let $f(x)$ be a sparse polynomial as in (1.2) of degree d and set $d_1 = d_p(p^{-t} f'(x))$. For any pair (i, j) with $1 \leq i < j \leq r$ let p_{ij} be the maximal prime divisor of $k_j - k_i$ (taking $p_{ij} = 1$ in case $k_j - k_i = 1$) and put

$$(6.1) \quad \mathcal{P} = \{p_{ij} : 1 \leq i < j \leq r\}.$$

Assume now that $p > 4r$, $p \mid (k_j - k_i)$ for some $k_i < k_j \leq d_1$ but that $p \notin \mathcal{P}$. Let

$$p_{ls} = \min\{p_{ij} : p \mid (k_j - k_i), k_i < k_j \leq d_1\},$$

and define

$$M := rd_1/p_{ls}.$$

Then if $p^e \parallel (k_j - k_i)$ is the maximum power of p dividing any of the differences $k_j - k_i$ that actually occur in the critical point congruence for $S(f, p^m)$, it follows from Lemma 4.1 that the multiplicity ν of any nonzero critical point satisfies

$$(6.2) \quad \nu < rp^e \leq r(k_j - k_i)/p_{ij} \leq M.$$

Let $S^*(f, p^m)$ denote the sum over a reduced residue system (modulo p^m) as in (1.7). For $j \geq 0$ define μ_j, t_j by

$$(6.3) \quad p^{\mu_j} \parallel (a_1 p^{jk_1}, \dots, a_r p^{jk_r}), \quad p^{\mu_j + t_j} \parallel (a_1 k_1 p^{jk_1}, \dots, a_r k_r p^{jk_r}).$$

Then we can write

$$(6.4) \quad S(f, p^m) = \sum_{j=0}^m S^*(f(p^j x), p^{m-j}) = \sum_{j=0}^m p^{\mu_j - j} S_j^*,$$

where for $0 \leq j \leq m$,

$$(6.5) \quad S_j^* = S^*(p^{-\mu_j} f(p^j x), p^{m-\mu_j}).$$

The critical point congruence associated with the sum S_j^* is just

$$g_j(x) := p^{-\mu_j - t_j} (a_1 k_1 p^{jk_1} x^{k_1 - 1} + \dots + a_r k_r p^{jk_r} x^{k_r - 1}) \equiv 0 \pmod{p}.$$

Viewing $g_j(x)$ as a polynomial over \mathbb{F}_p we observe that for any $j < m$ the largest degree term of $g_{j+1}(x)$ is at most the smallest degree term of $g_j(x)$. Indeed, if $p^{t_j + \mu_j} \parallel a_l k_l p^{jk_l}$ then $p^{t_j + \mu_j + k_l} \parallel a_l k_l p^{(j+1)k_l}$ and $p^{t_j + \mu_j + k_l + 1} \parallel a_l k_l p^{(j+1)k_l}$ for $l > I$. It follows that the degrees of the g_j are nonincreasing (with j) and that at most r of the $g_j(x)$ can have more than one nonzero term. The rest of the $g_j(x)$ are monomials and therefore the associated sums S_j^* are zero, provided $m - \mu_j \geq 2$. Thus there are at most r values of $j \leq m$ for which $m - \mu_j \geq 2$ and S_j^* is nonzero. Moreover, for these nonzero sums the multiplicity of any nonzero critical point is bounded above by M .

Say $d_1 = k_I - 1$ for some I . Then since $p^t \parallel a_I k_I$ it is easily seen that for $0 \leq j \leq m$,

$$(6.6) \quad \mu_j + t_j \leq t + j(d_1 + 1).$$

We split the sum in (6.4) into two parts according as $m - t_j - \mu_j \geq 8M$ or not. If this inequality holds then since S_j^* has at most p critical points,

each of multiplicity $\leq M$, it follows from Lemma 2.3 that

$$p^{\mu_j-j}|S_j^*| \leq p^{\mu_j-j} 4pp^{t_j/M} p^{(m-\mu_j)(1-1/M)} = \frac{4p}{p^{(m-\mu_j-t_j)/(2M)}} \frac{p^{m-j}}{p^{(m-\mu_j-t_j)/(2M)}}.$$

Now $(m - \mu_j - t_j)/(2M) \geq 4$. Also, since $p > 2r$, $2M < d_1$ and so by (6.6),

$$\frac{m - \mu_j - t_j}{2M} \geq \frac{m - t - j(d_1 + 1)}{d_1 + 1} = \frac{m - t}{d_1 + 1} - j.$$

It follows that

$$(6.7) \quad p^{\mu_j-1}|S_j^*| \leq \frac{4}{p^3} p^{t/(d_1+1)} p^{m(1-1/(d_1+1))}.$$

We consider next the set of j for which $m - t_j - \mu_j < 8M$ and let j_0 denote the least such j . Then

$$\sum_{j \geq j_0} p^{\mu_j-j}|S_j^*| \leq p^{m-j_0} = p^{(m-t)/(d_1+1)-j_0} p^{t/(d_1+1)} p^{m(1-1/(d_1+1))}.$$

Now

$$(m - t) - j_0(d_1 + 1) \leq 8M + t_{j_0} + \mu_{j_0} - t - j_0(d_1 + 1) \leq 8M = 8rd_1/p_{l_s},$$

by (6.6). Thus

$$(6.8) \quad \sum_{j \geq j_0} p^{\mu_j-j}|S_j^*| \leq p^{8r/p_{l_s}} p^{t/(d_1+1)} p^{m(1-1/(d_1+1))}.$$

From (6.7) and (6.8) we conclude that

$$\begin{aligned} |S(f, p^m)| &\leq (4r/p^3 + p^{8r/p_{l_s}}) p^{t/(d_1+1)} p^{m(1-1/(d_1+1))} \\ &\leq p^{8r/p_{l_s}} (1 + 1/p^2) p^{t/(d_1+1)} p^{m(1-1/(d_1+1))}. \end{aligned}$$

This establishes

LEMMA 6.1. *Suppose that $p \mid (k_j - k_i)$ for some $k_i < k_j \leq d_1$, $p > 4r$ and $p \notin \mathcal{P}$. Then*

$$|S(f, p^m)| \leq (1 + 1/p^2) p^{8r/p_{l_s}} p^{t/(d_1+1)} p^{m(1-1/(d_1+1))}$$

for some $p_{l_s} \in \mathcal{P}$ with $p \mid (k_l - k_s)$, $p < p_{l_s}$.

7. Proof of Theorem 1.1. For any prime power p^m and polynomial f let

$$(7.1) \quad R(f, p^m) = \frac{|S(f, p^m)|}{p^{m(1-1/d)}}.$$

Let f be a sparse polynomial with r terms and let q be a positive integer such that $d_p(f) \geq 1$ for all prime divisors p of q . Write

$$\prod_{p^m \parallel q} R(f, p^m) = P_1 P_2 P_3 P_4 P_5 P_6$$

where the P_i are products over the prime power divisors of q satisfying the following constraints (counting prime powers only once if they happen to satisfy more than one constraint):

$$(7.2) \quad P_1 = \prod_{m=1} R(f, p^m),$$

$$(7.3) \quad P_2 = \prod_{1 < m \leq t+1} R(f, p^m),$$

$$(7.4) \quad P_3 = \prod_{\substack{p \leq C_3 \text{ or} \\ p \leq r^4 \text{ or } p \in \mathcal{P}}} R(f, p^m),$$

$$(7.5) \quad P_4 = \prod_{\substack{p > r^4 \\ p | (k_j - k_i) \text{ for some } k_i < k_j \leq d_1, p \notin \mathcal{P}}} R(f, p^m),$$

$$(7.6) \quad P_5 = \prod_{\substack{m \geq t+2, 50d > p > r^4, p > C_3 \text{ and} \\ p \nmid (k_j - k_i) \text{ for all } k_i < k_j \leq d_1}} R(f, p^m),$$

$$(7.7) \quad P_6 = \prod_{\substack{m \geq t+2, p > r^4, p > C_2 \text{ and} \\ p \geq 50d}} R(f, p^m),$$

where C_2, C_3 are the constants in Lemmas 5.3 and 5.4 respectively, and \mathcal{P} is the set (6.1) of exceptional primes. By (1.6) the theorem follows if we show that each of the products P_i is bounded by a constant depending only on r .

By Lemma 5.2, the Weil bound (5.1) and the trivial bound $R(f, p) \leq p^{1/d}$ we have

$$\begin{aligned} P_1 &\leq \prod_{p < C_2} R(f, p) \prod_{p \leq r^4} R(f, p) \prod_{p < 50d} R(f, p) \\ &\leq (1.75)^{C_2+r^4} \prod_{p < 50d} p^{1/d} \ll (1.75)^{r^4}. \end{aligned}$$

For the next few products we need the following

LEMMA 7.1. *Let f be a sparse polynomial with r terms of degree d . For any prime p let $t_p = \text{ord}_p(f'(x))$. Then letting p run through the set of all primes for which $d_p(f) \geq 1$ we have*

$$\prod_{p, d_p(f) \geq 1} p^{t_p} \leq d^r.$$

Proof. Let $f(x) = a_1 x^{k_1} + \dots + a_r x^{k_r}$ and p be a prime with $d_p(f) \geq 1$. Then for some i , $p \nmid a_i$, and so for this value of i , $p^{t_p} \mid k_i$. Thus the product over all such p^{t_p} is a divisor of $k_1 \dots k_r$. ■

(We continue to write t for t_p .) For P_2 the condition $1 < m \leq t + 1$ implies that $t \geq 1$ and so $m \leq 2t$. Thus we trivially have

$$P_2 \leq \prod_p p^{m/d} \leq \prod_p p^{2t/d} \leq d^{2r/d} \leq 2.1^r.$$

The number of primes in the product P_3 is less than $r^4/2 + r^2 + C_3 < r^4 + C_3$ and so by Lemma 1.1, $P_3 \leq 5^{r^4+C_3}$. For P_4 we apply Lemma 6.1, to obtain

$$\begin{aligned} P_4 &\leq \prod_p \left(1 + \frac{1}{p^2}\right) \left(\prod_p p^{t/d}\right) \prod_{1 \leq i < j \leq r} \prod_{p \leq p_{i,j}} p^{8r/p_{i,j}} \\ &\ll d^{r/d} \prod_{1 \leq i < j \leq r} C_5^{4r} \ll 1.5^r C_5^{2r^3} \end{aligned}$$

for some absolute constant C_5 . We may take $C_5 = \sup_x e^{\theta(x)/x}$, where $\theta(x) = \sum_{p \leq x} \log p$.

For P_5 , we apply Lemma 5.4(iii) to obtain,

$$P_5 \leq \prod_{p < 50d} p^{(2r+t)/d} \leq \prod_p p^{t/d} \prod_{p < 50d} p^{2r/d} \leq d^{r/d} e^{2r\theta(50d)/d} \leq 1.5^r C_5^{100r}.$$

Finally, we apply Lemma 5.3 to P_6 to obtain

$$P_6 \leq \prod_p p^{t/d} \leq d^{r/d} \leq 1.5^r.$$

Thus the product $P_1 P_2 P_3 P_4 P_5 P_6$ is bounded above by a constant depending only on r .

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