

Some results on Oppenheim's "Factorisatio Numerorum" function

by

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1. Introduction. Let $f(n)$ denote the number of distinct unordered factorisations of the natural number n into factors larger than 1. For example, $f(28) = 4$ as 28 has the following factorisations:

$$28, 2 \cdot 14, 4 \cdot 7, 2 \cdot 2 \cdot 7.$$

In this paper, we address three aspects of the function $f(n)$. For the first aspect, in [1], Canfield, Erdős and Pomerance mention without proof that the number of values of $f(n)$ that do not exceed x is $x^{o(1)}$ as $x \rightarrow \infty$. Our first theorem in this note makes this result explicit.

For a set \mathcal{A} of positive integers we put $\mathcal{A}(x) = \{n \in \mathcal{A} : n \leq x\}$.

THEOREM 1. *Let $\mathcal{A} = \{f(m) : m \in \mathbb{N}\}$. Then*

$$\#\mathcal{A}(x) = x^{O(\log \log \log x / \log \log x)}.$$

Recall that Oppenheim [8] and independently Szekeres and Turán [11] considered the average value of $f(n)$ in the interval $(0, x]$ showing that

$$(1) \quad \frac{1}{x} \sum_{0 < n \leq x} f(n) = \frac{e^{2\sqrt{\log x}}}{2\sqrt{\pi}(\log x)^{3/4}} \left(1 + O\left(\frac{1}{\sqrt{\log x}}\right) \right).$$

There is a large body of literature addressing average values of various arithmetic functions in short intervals. Our next result gives a lower bound for the average of $f(n)$ over a short interval $(x, x+y]$ which is of the same order as the average of $f(n)$ over the interval $(0, y]$.

THEOREM 2. *Uniformly for $x > 0$ and $y \geq 2$, we have*

$$\frac{1}{y} \sum_{x < n \leq x+y} f(n) \gg \frac{e^{2\sqrt{\log y}}}{(\log y)^{3/4}}.$$

Finally, there are also several results addressing the behaviour of positive integers n which are multiples of some other arithmetic function of n . See, for example, [3], [5], [9] and [10] for problems related to counting positive integers n which are divisible by either $\omega(n)$, $\Omega(n)$ or $\tau(n)$, where these functions are the number of distinct prime factors of n , the number of total prime factors of n , and the number of divisors of n , respectively. Our next and last result gives upper and lower bounds on the counting function of the set of positive integers n which are multiples of $f(n)$.

THEOREM 3. *Let $\mathcal{B} = \{n : f(n) \mid n\}$. Then*

$$\#\mathcal{B}(x) = \frac{x}{(\log x)^{1+o(1)}} \quad \text{as } x \rightarrow \infty.$$

2. Preliminaries and lemmas. The function $f(n)$ is related to various partition functions. For example, $f(2^n) = p(n)$, where $p(n)$ is the number of partitions of n . Furthermore, $f(p_1 \cdots p_k) = B_k$, where B_k is the k th Bell number which counts the number of partitions of a set with k elements into nonempty disjoint subsets. In general, $f(p_1^{\alpha_1} \cdots p_k^{\alpha_k})$ is the number of partitions of a multiset consisting of α_i copies of $\{i\}$ for each $i = 1, \dots, k$. Throughout the paper, we write $\log x$ for the natural logarithm of x . We use p and q for prime numbers, O and o for the Landau symbols, and \ll and \gg for the Vinogradov symbols. The following asymptotic formula for the k th Bell number is due to de Bruijn [4].

LEMMA 1.

$$\frac{\log B_k}{k} = \log k - \log \log k - 1 + \frac{\log \log k}{\log k} + \frac{1}{\log k} + O\left(\frac{(\log \log k)^2}{(\log k)^2}\right).$$

We also need the Stirling numbers of the second kind $S(k, l)$ which count the number of partitions of a k -element set into l nonempty disjoint subsets. Clearly,

$$(2) \quad B_k = \sum_{l=1}^k S(k, l).$$

We now formulate and prove a few lemmas about the function $f(n)$ which will come in handy later on.

The next lemma is an easy observation, so we state it without proof.

LEMMA 2. *If $a \mid b$, then $f(a) \leq f(b)$.*

We let p_n denote the n th prime number and $\alpha_1(n)$ denote the maximal exponent of a prime appearing in the prime factorisation of n . Let n be a positive integer with prime factorisation

$$n = q_1^{\alpha_1} \cdots q_k^{\alpha_k},$$

where q_1, \dots, q_k are distinct primes and $\alpha_1(n) := \alpha_1 \geq \dots \geq \alpha_k$. We put $n_0 = p_1^{\alpha_1} \dots p_k^{\alpha_k}$, and observe that $f(n) = f(n_0)$. This observation will play a crucial role in the proof of Theorem 1.

The following lemma gives upper bounds for $\alpha_1(n)$ and $\omega(n)$ when $f(n) \leq x$.

LEMMA 3. *Let $n = q_1^{\alpha_1} \dots q_k^{\alpha_k}$, where $\alpha_1 \geq \dots \geq \alpha_k$ and $f(n) \leq x$. Then*

- (i) $\alpha_1 = O((\log x)^2)$;
- (ii) $k = \omega(n) = O(\log x / \log \log x)$.

Proof. It follows from Lemma 2 that

$$f(n) \geq f(q_1^{\alpha_1}) = p(\alpha_1).$$

Using the asymptotic formula

$$(3) \quad p(n) = (1 + o(1)) \frac{\exp(\pi \sqrt{2n/3})}{4n\sqrt{3}} \quad \text{as } n \rightarrow \infty,$$

due to Hardy and Ramanujan [6], we conclude that $\exp(c\sqrt{\alpha_1}) \leq x$ with some constant $c > 0$. Hence, (i) follows. In order to prove (ii), let $n'_0 = p_1 \dots p_k$. By Lemma 2, we have $f(n'_0) \leq f(n) \leq x$. Furthermore, $f(n'_0) = B_k$. It now follows from Lemma 1 that

$$\exp((1 + o(1))k \log k) = B_k \leq x$$

as $k \rightarrow \infty$, yielding

$$k = O\left(\frac{\log x}{\log \log x}\right),$$

which completes the proof of the lemma. ■

3. Proofs of the theorems

3.1. Proof of Theorem 1. For a positive integer n , we let again n_0 and $\alpha_1(n)$ be the functions defined earlier. We let $\mathcal{A}(x) = \{m_1, \dots, m_t\}$ be such that $m_1 < \dots < m_t$ and let $\mathcal{N} = \{n_1, \dots, n_t\}$ be positive integers such that n_i is minimal among all positive integers n with $f(n) = m_i$ for all $i = 1, \dots, t$. It is clear that if $n \in \mathcal{N}$, then n is of the form n_0 . Since $\#\mathcal{A}(x) = t = \#\mathcal{N}$, it suffices to bound the cardinality of \mathcal{N} .

We partition this set as $\mathcal{N} = \mathcal{N}_1 \cup \mathcal{N}_2 \cup \mathcal{N}_3$, where

$$\begin{aligned} \mathcal{N}_1 &= \{n \in \mathcal{N} : \alpha_1(n) \leq \log \log x\}, \\ \mathcal{N}_2 &= \left\{n \in \mathcal{N} : \omega(n) \leq \frac{\log x}{(\log \log x)^2}\right\}, \quad \mathcal{N}_3 = \mathcal{N} \setminus (\mathcal{N}_1 \cup \mathcal{N}_2). \end{aligned}$$

If $n \in \mathcal{N}_1$, then n has at most $O(\log x / \log \log x)$ prime factors (by Lemma 3), each appearing with an exponent of at most $\log \log x$.

Therefore,

$$(4) \quad \#\mathcal{N}_1 = (\log \log x)^{O(\log x / \log \log x)} = x^{O\left(\frac{\log \log \log x}{\log \log x}\right)}.$$

Next, we observe that an integer in \mathcal{N}_2 has at most $\log x / (\log \log x)^2$ prime factors, each appearing with an exponent $O((\log x)^2)$ (by Lemma 3). Thus,

$$(5) \quad \begin{aligned} \#\mathcal{N}_2 &\leq (O(\log x)^2)^{\frac{\log x}{(\log \log x)^2}} = \exp\left(\frac{(2 + o(1)) \log x}{\log \log x}\right) \\ &= x^{o\left(\frac{\log \log \log x}{\log \log x}\right)} \quad \text{as } x \rightarrow \infty. \end{aligned}$$

Finally, let $n \in \mathcal{N}_3$, and write it as

$$n = p_1^{\alpha_1} \cdots p_i^{\alpha_i} p_{i+1}^{\alpha_{i+1}} \cdots p_k^{\alpha_k},$$

where we put

$$i := \max\{j \leq k : \alpha_j \geq y\} \quad \text{with } y := \lfloor \log \log x / \log \log \log x \rfloor.$$

Observe that the divisors $p_{i+1}^{\alpha_{i+1}} \cdots p_t^{\alpha_t}$ of the numbers $n \in \mathcal{N}_3$ can be chosen in at most

$$(6) \quad y^k = y^{O\left(\frac{\log x}{\log \log x}\right)} = \exp\left(O\left(\frac{\log x \log \log \log x}{\log \log x}\right)\right)$$

ways. Furthermore, by Lemma 3, the numbers $n' = p_1^{\alpha_1} \cdots p_i^{\alpha_i}$ can trivially be chosen in at most

$$(O((\log x)^2))^i = \exp(O(i \log \log x))$$

ways. Thus, writing \mathcal{N}_4 for the subset of \mathcal{N}_3 such that $i \leq \log x / (\log \log x)^2$, we get

$$(7) \quad \#\mathcal{N}_4 \leq \exp\left(O\left(\frac{\log x}{\log \log x}\right)\right).$$

From now on, we look at $n \in \mathcal{N}_5 = \mathcal{N}_3 \setminus \mathcal{N}_4$.

For each t , we let k_t be such that $S(t, k_t)$ is maximal among the numbers $S(t, k)$ for $k = 1, \dots, t$. By formula (2), the definition of k_t , and Lemma 1,

$$S(t, k_t) \geq \frac{Bt}{t} = \frac{\exp((1 + o(1))t \log t)}{t} = \exp((1 + o(1))t \log t)$$

as $t \rightarrow \infty$. We now claim that

$$f(n) \geq f(n') \geq f((p_1 \cdots p_i)^y) \geq \frac{S(i, k_i)^y}{(y k_i)!}.$$

The first two inequalities follow immediately from Lemma 2, so let us prove the last one.

Note that $S(i, k_i)$ counts the number of factorisations of $p_1 \cdots p_i$ into precisely k_i factors. Therefore, $(S(i, k_i))^y$ counts the number of factorisa-

tions of $(p_1 \cdots p_i)^y$ into $k_i y$ square-free factors, where we count each such factorisation at most $(k_i y)!$ times. This establishes the claim.

Since i tends to infinity as $x \rightarrow \infty$ for all $n \in \mathcal{N}_5$, we get

$$S(i, k_i)^y \geq \exp((1 + o(1))yi \log i)$$

as $x \rightarrow \infty$. Furthermore, we trivially have

$$(k_i y)! \leq (k_i y)^{k_i y} = \exp(k_i y \log(k_i y)).$$

Thus,

$$(8) \quad f(n) \geq \frac{S(i, k_i)^y}{(k_i y)!} \geq \exp((1 + o(1))yi \log i - k_i y \log(k_i y))$$

as $x \rightarrow \infty$. We next show that for our choices of y and i we have

$$k_i y \log(k_i y) = o(yi \log i) \quad \text{as } x \rightarrow \infty.$$

Indeed, using the fact

$$k_i = (1 + o(1)) \frac{i}{\log i} \quad \text{as } i \rightarrow \infty$$

(see, for example, [2]), we see that the above condition is equivalent to

$$\log y = o((\log i)^2),$$

which holds as $x \rightarrow \infty$ because

$$y = \log \log x / \log \log \log x \quad \text{and} \quad i > \log x / (\log \log x)^2.$$

Now the inequality $f(n) \leq x$ together with (8) and the fact that $\log i \geq (1 + o(1)) \log \log x$ implies that

$$(9) \quad i \leq (1 + o(1)) \frac{\log x}{y \log \log x} \quad \text{as } x \rightarrow \infty.$$

Thus, the numbers n' can be chosen in at most

$$(10) \quad (O((\log x)^2))^i \leq (O((\log x)^2))^{(1+o(1))\frac{\log x}{y \log \log x}} = x^{O(\frac{\log \log \log x}{\log \log x})}$$

ways. As we have already seen at (6), the complementary divisor $n/n' = p_{i+1}^{\alpha_{i+1}} \cdots p_t^{\alpha_t}$ of n can be chosen in at most

$$(11) \quad x^{O(\log \log \log x / \log \log x)}$$

ways also. Thus, the total number of choices for n in \mathcal{N}_5 is

$$(12) \quad \#\mathcal{N}_5 \leq x^{O(\log \log \log x / \log \log x)}.$$

Hence, from estimates (7) and (12), we get

$$(13) \quad \#\mathcal{N}_3 \leq \#\mathcal{N}_4 + \#\mathcal{N}_5 \leq x^{O(\log \log \log x / \log \log x)}.$$

From estimates (4), (5) and (13), we finally get

$$\#\mathcal{N} \leq \#\mathcal{N}_1 + \#\mathcal{N}_2 + \#\mathcal{N}_3 \leq x^{O(\log \log \log x / \log \log x)},$$

which completes the proof of the theorem.

3.2. Proof of Theorem 2. For ease of notation we put

$$S(x, y) := \sum_{x < n \leq x+y} f(n).$$

Let z be some function of y tending to infinity with it such that $z \log z < o(\sqrt{\log y})$ as $y \rightarrow \infty$. Assume that $0 < x \leq zy$. Write

$$S(x, y) = S(0, x + y) - S(0, x).$$

Observe that

$$\log(x + y) = \log y + O(\log z),$$

therefore

$$\begin{aligned} \exp(2\sqrt{\log(x + y)}) &= \exp(2\sqrt{\log y + O(\log z)}) \\ &= \exp\left(2\sqrt{\log y} + O\left(\frac{\log z}{\sqrt{\log y}}\right)\right) \\ &= e^{2\sqrt{\log y}} \left(1 + O\left(\frac{\log z}{\sqrt{\log y}}\right)\right), \end{aligned}$$

and a similar estimate holds for $\exp(2\sqrt{\log x})$. Furthermore,

$$\frac{1}{(\log(x + y))^{3/4}} = \frac{1}{(\log y + O(\log z))^{3/4}} = \frac{1}{(\log y)^{3/4}} \left(1 + O\left(\frac{\log z}{\log y}\right)\right),$$

and again a similar estimate holds for $1/(\log x)^{3/4}$. Thus, using estimate (1), we see that in the range $0 < x \leq zy$ the desired sum is

$$S(x, y) = S(0, x + y) - S(0, x) = \frac{ye^{2\sqrt{\log y}}}{2\sqrt{\pi}(\log y)^{3/4}} \left(1 + O\left(\frac{z \log z}{(\log y)^{1/2}}\right)\right).$$

This is even an asymptotic as $y \rightarrow \infty$ if we take $z := (\log y)^{1/2}(\log \log y)^{-2}$. We next assume that $x > yz$. For each integer $n \in (0, y]$, let $m(n)$ be the largest multiple of n in $(x, x + y]$ and write it as $m(n) = m_0(n) \cdot n$. Observe that $m_0(n) \geq x/n > x/y$. Thus, if $x \geq y^2$, then $x/n > y$. Let $\mathcal{M} = \{m(n) : n \in (0, y]\}$ and observe that in this range

$$\sum_{x < n \leq x+y} f(n) \geq \sum_{m \in \mathcal{M}} f(m) \geq \sum_{0 < n \leq y} f(n),$$

where the last inequality follows by considering only factorisations of $m \in \mathcal{M}$ which are of the form

$$n_1 \cdots n_k \cdot m_0(n)$$

for some $n \in (0, y]$, by remarking also that since $m_0(n) > y$, distinct factorisations of n will yield distinct factorisations of $m \in \mathcal{M}$. Thus, if $x > y^2$, the above argument yields

$$S(x, y) \geq S(0, y) = \frac{ye^{2\sqrt{\log y}}}{2\sqrt{\pi}(\log y)^{3/4}} \left(1 + O\left(\frac{1}{\sqrt{\log y}}\right)\right).$$

We now suppose that $yz \leq x \leq y^2$. We let

$$S(0, y) - S(0, y/2) = \sum_{y/2 < n \leq y} f(n) = S(0, y) \left(\frac{1}{2} + O\left(\frac{1}{\sqrt{\log y}}\right) \right).$$

To each factorisation $n_1 \cdots n_k$ of some $n \in \mathcal{I} := [y/2, y]$ we associate, as before, the factorisation $n_1 \cdots n_k \cdot m_0(n)$ of $m(n)$. Observe that $m_0(n) \in (x/n, x/n + y/n] \subset \mathcal{J} := (x/y, 2x/y + 2]$. Let $f_1(n)$ be the number of factorisations of n with two or more parts in \mathcal{J} . Note that $f_1(n) = 0$ unless $(x/y)^2 \leq y$. Writing a factorisation counted by $f_1(n)$ as

$$a \cdot b \cdot m_1 \cdots m_s, \quad \text{where } a, b \in \mathcal{J},$$

we get

$$\sum_{y/2 \leq n \leq y} f_1(n) \leq \sum_{\substack{a \leq b \\ a, b \in \mathcal{J}}} \sum_{m \leq y/ab} f(m) = \sum_{\substack{a \leq b \\ a, b \in \mathcal{J}}} S(0, y/ab).$$

We split the above sum at $ab \leq y/2$. In the low range, we use the fact that the function $u \mapsto e^{2\sqrt{\log u}}/(\log u)^{3/4}$ is increasing, to get

$$\sum_{\substack{a \leq b \\ a, b \in \mathcal{J} \\ ab < y/2}} S(0, y/ab) \leq \frac{ye^{2\sqrt{\log y}}}{2\sqrt{\pi}(\log y)^{3/4}} \left(\sum_{\substack{a \leq b \\ a, b \in \mathcal{J}}} \frac{1}{ab} \right) \left(1 + O\left(\frac{1}{(\log y)^{1/2}}\right) \right).$$

Observe that

$$\begin{aligned} \sum_{\substack{a \leq b \\ a, b \in \mathcal{J}}} \frac{1}{ab} &\leq \sum_{a \geq x/y} \frac{1}{a^2} + \frac{1}{2} \left(\sum_{a \in \mathcal{J}} \frac{1}{a} \right)^2 \\ &\leq \left(\log\left(\frac{2x}{y} + 2\right) - \log\left(\frac{x}{y}\right) + O\left(\frac{1}{z}\right) \right)^2 + O\left(\frac{1}{z}\right) \\ &= \frac{1}{2} \left(\log 2 + O\left(\frac{1}{z}\right) \right)^2 + O\left(\frac{1}{z}\right) = \frac{(\log 2)^2}{2} + O\left(\frac{1}{z}\right). \end{aligned}$$

In the larger range, we have $S(0, y/ab) = 1$. Thus, under the assumption that $(x/y)^2 \leq y$,

$$\sum_{\substack{a \leq b \\ a, b \in \mathcal{J} \\ ab > y/2}} S(0, y/ab) \leq \sum_{a, b \in \mathcal{J}} 1 \ll (x/y)^2 \leq y.$$

Putting everything together, we get

$$\sum_{y/2 \leq n \leq y} f_1(n) \leq S(0, y) \left(\frac{(\log 2)^2}{2} + O\left(\frac{(\log \log y)^2}{\sqrt{\log y}}\right) \right).$$

Therefore,

$$\begin{aligned} \sum_{y/2 \leq n \leq y} (f(n) - f_1(n)) &\geq S(0, y) \left(\frac{1}{2} - \frac{(\log 2)^2}{2} + O\left(\frac{(\log \log y)^2}{\sqrt{\log y}}\right) \right) \\ &\gg S(0, y). \end{aligned}$$

We now look only at the factorisations $m_1 \cdots m_k m_0(n)$ of $m(n)$ for $n \in [y/2, y]$ arising from factorisations $m_1 \cdots m_k$ of n counted by $f(n) - f_1(n)$. These might not be distinct but since the factorisation $m_1 \cdots m_k$ of n has at most one part in \mathcal{J} , the interval containing $m_0(n)$ for all n under scrutiny, it follows that each such factorisation is counted at most twice. This shows

$$S(x, y) \geq \frac{1}{2} \sum_{y/2 \leq n \leq y} (f(n) - f_1(n)) \gg S(0, y),$$

which is what we wanted to prove.

3.3. Proof of Theorem 3. We observe that all primes are in \mathcal{A} since $f(p) = 1$ for all prime p . Thus,

$$\#\mathcal{A}(x) \gg \frac{x}{\log x}.$$

This completes the lower bound part of the theorem. To obtain the upper bound, we cover the set $\mathcal{A}(x)$ by three sets \mathcal{A}_1 , \mathcal{A}_2 and \mathcal{A}_3 as follows:

$$\begin{aligned} \mathcal{A}_1 &= \{n \leq x : \Omega(n) > 10 \log \log x\}, \\ \mathcal{A}_2 &= \left\{ n \leq x : \omega(n) < \frac{\log \log x}{\log \log \log x} \right\}, \\ \mathcal{A}_3 &= \{n \leq x : n \equiv 0 \pmod{f(n)}, n \notin \mathcal{A}_1 \cup \mathcal{A}_2\}. \end{aligned}$$

We recall the bound

$$\#\{n \leq x : \Omega(n) = k\} \ll \frac{kx \log x}{2^k}$$

valid uniformly in k (see, for example, Lemma 13 in [7]). Using the above estimate, we get

$$(14) \quad \#\mathcal{A}_1 \leq x \sum_{k > 10 \log \log x} \frac{k}{2^k} \ll \frac{x \log \log x}{2^{10 \log \log x}} = o\left(\frac{x}{\log x}\right)$$

as $x \rightarrow \infty$. To find an upper bound for \mathcal{A}_2 , we use the bounds (see page 200 of [12])

$$\#\{n \leq x : \omega(n) = k\} \ll \frac{x}{(k-1)!} \frac{(\log \log x + c_1)^{k-1}}{\log x},$$

where $c_1 > 0$ is some constant. Using the elementary estimate $m! \geq (m/e)^m$ with $m = k-1$, we get

$$\#\{n \leq x : \omega(n) = k\} \ll \frac{x}{\log x} \left(\frac{e \log \log x + c_2}{k-1} \right)^{k-1},$$

where $c_2 = ec_1$. The right hand side is an increasing function of k in our range for k versus x when x is large. Since $k < \log \log x / \log \log \log x$, we deduce that

$$(15) \quad \#\mathcal{A}_2 \ll \frac{x}{\log x} (O(\log \log \log x))^{\frac{\log \log x}{\log \log \log x}} = \frac{x}{(\log x)^{1+o(1)}}$$

as $x \rightarrow \infty$.

Finally, we estimate \mathcal{A}_3 . Each $n \in \mathcal{A}_3$ can be written as

$$n = q_1^{\alpha_1} \cdots q_k^{\alpha_k},$$

where q_1, \dots, q_k are distinct primes, $\alpha_1 \geq \cdots \geq \alpha_k$, $\alpha_1 + \cdots + \alpha_k \leq 10 \log \log x$ and $k > K := \lfloor \log \log x / \log \log \log x \rfloor$. Let \mathcal{T} be the set of all such tuples $(k, \alpha_1, \dots, \alpha_k)$. For each such n , we have

$$\begin{aligned} f(n) &\geq B_K \geq \exp((1 + o(1))K \log K) \geq \exp((1 + o(1)) \log \log x) \\ &= (\log x)^{1+o(1)} \end{aligned}$$

as $x \rightarrow \infty$. The number of tuples $(k, \alpha_1, \dots, \alpha_k)$ satisfying the above conditions is at most

$$\#\mathcal{T} \ll \log \log x \sum_{n \leq 10 \log \log x} p(n),$$

where again $p(n)$ is the partition function of n . Using estimate (3), we conclude that

$$\#\mathcal{T} \ll (\log \log x)^2 \exp(O(\sqrt{\log \log x})) = (\log x)^{o(1)} \quad \text{as } x \rightarrow \infty.$$

Thus,

$$(16) \quad \#\mathcal{A}_3 \leq \sum_{(k, \alpha_1, \dots, \alpha_k) \in \mathcal{T}} \frac{x}{f(p_1^{\alpha_1} \cdots p_k^{\alpha_k})} \leq \frac{x \#\mathcal{T}}{B_K} = \frac{x}{(\log x)^{1+o(1)}}$$

as $x \rightarrow \infty$. Now inequalities (14), (15) and (16) yield the desired upper bound and complete the proof.

4. Comments. Quite likely, the results of Theorems 1 and 2 are not best possible. In this respect, we suggest the following questions:

QUESTION 1. *Is it true that $\#\mathcal{A}(x) = \exp(O(\sqrt{\log x}))$?*

QUESTION 2. *In the notations used in the proof of Theorem 2, is it true that*

$$S(x, y) \geq (1 + o(1))S(0, y) \quad \text{as } y \rightarrow \infty?$$

Namely, is it true that the average value of $f(n)$ in the interval $(0, y]$ is an asymptotic lower bound for the average value of $f(n)$ in any interval of length y as $y \rightarrow \infty$?

Concerning Question 2 above, observe that our proof indicates that this is indeed the case except when $x \in [yz, y^2]$, where $z = (\log y)^{1/2}(\log \log y)^{-2}$.

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