

DEFINITION 1. Let $A_k = \{1, \dots, k\}$ with $k \geq 2$. A sequence ω in the alphabet A_k is called an *Arnoux–Rauzy sequence* if it satisfies the following three conditions:

- ω is recurrent,
- the complexity function $p_\omega(m)$ equals $(k - 1)m + 1$,
- for each m there is exactly one right special and one left special factor of ω of length m .

Recall that a factor u of ω is called *right special* (resp. *left special*) if u is a prefix (resp. suffix) of at least two words of length $|u| + 1$ which are factors of ω . A word which is both right and left special is called *bispecial*. Arnoux–Rauzy sequences are a natural generalization of Sturmian words; Sturmian words correspond to taking $k = 2$ in the above definition. For $k = 3$ the combinatorial conditions listed in Definition 1 distinguish them from other sequences of complexity $2n + 1$ such as those obtained by coding trajectories of 3-interval exchange transformations [16, 17, 18], or those of *Chacon type*, i.e., topologically isomorphic to the subshift generated by the Chacon sequence [6, 15]. Perhaps the best known example on three letters is the so-called *Tribonacci sequence* defined as the fixed point of the morphism $\tau(1) = 12$, $\tau(2) = 13$ and $\tau(3) = 1$. In [26] Rauzy showed that the subshift generated by τ is isomorphic (in measure) to an exchange of three fractal domains in \mathbb{R}^2 which generate a tiling of the plane.

Arnoux and Rauzy [2] showed that each Arnoux–Rauzy sequence may be geometrically realized by an exchange of $2k$ intervals on the circle, and is uniquely ergodic. It was further believed, as in the case of Tribonacci and the Rauzy fractal, that each Arnoux–Rauzy sequence is measure isomorphic to a rotation on the torus, i.e., is obtained by a symbolic coding of the trajectories of points under a rotation on the k -dimensional torus with respect to a natural partition. This was recently disproved by Cassaigne–Ferenczi–Zamboni in [4] where the authors exhibited an Arnoux–Rauzy sequence ω on a 3-letter alphabet $\{0, 1, 2\}$ which is *totally unbalanced* in the following sense: for each $n > 0$ there exist two factors of ω of equal length, with one having at least n more occurrences of the letter 0 than the other. It follows that the cylinder $[0]$ is not a *bounded remainder set* (in the sense of Kesten [21]) and hence via an unpublished result of Rauzy later generalized by Ferenczi [14], either ω is not a natural coding of a rotation in \mathbb{R}^n modulo a lattice, or the A–R sequence $\omega(0)$, obtained by coding ω according to first returns to 0, is not a natural coding of a rotation in \mathbb{R}^n modulo a lattice.

Arnoux–Rauzy sequences have been extensively studied from many different points of view in connection with dynamical systems (see [1, 2, 7, 8, 20]), number theory (see [7, 9, 19, 20, 27, 29, 30]) and combinatorics (see [4, 5, 7, 12, 20, 27]).

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2. Counting words. We fix $k \geq 2$ and let $\mathcal{AR}_\infty = \mathcal{AR}_\infty(k)$ be the set of all Arnoux–Rauzy sequences on the alphabet A_k . We denote by \mathcal{AR} the set of all *Arnoux–Rauzy words*, that is, the set of all words u (including the empty word) such that u is a factor of some Arnoux–Rauzy sequence $x \in \mathcal{AR}_\infty$. For each $n \geq 0$ we let \mathcal{AR}_n be the set of all $u \in \mathcal{AR}$ of length n .

For each $a \in A_k$ define the morphism τ_a on A_k by $\tau_a(a) = a$ and $\tau_a(b) = ab$ for all $b \in A_k$ different from a . Then it is proved in [2] (see also [27]) that each Arnoux–Rauzy sequence ω is in the shift orbit closure of a unique sequence of the form

$$\omega_* = \lim_{j \rightarrow \infty} \tau_{i_1} \circ \dots \circ \tau_{i_j}(1)$$

where the sequence of indices (i_j) (called the *coding sequence*) takes values in A_k . Moreover each $a \in A_k$ occurs in the coding sequence an infinite number of times. The sequence ω_* is called a *characteristic Arnoux–Rauzy sequence*.

LEMMA 1. *Let $u \in \mathcal{AR}$ and suppose that for some $b, c \in A_k$ distinct, ub and uc are in \mathcal{AR} . Then there exists an Arnoux–Rauzy sequence $\omega \in \mathcal{AR}_\infty$ which contains as factors the k words $u1, u2, \dots, uk$. In other words if u is a right special factor of \mathcal{AR} then u is a right special factor of some Arnoux–Rauzy sequence ω .*

Proof. We proceed by induction on the length of u . The result is clearly true if u is empty, or if $|u| = 1$. Writing $u = av$ with $a \in A_k$ and $|v| \geq 1$, we make the inductive hypothesis that the result of the lemma holds for all words of length smaller than $|u|$. Thus avb and avc are each in \mathcal{AR} . Without loss of generality we can assume 1 is the last letter of v .

CASE 1: $a = 1$. If av is of the form $av = 1^n$, then for each Arnoux–Rauzy sequence ω , the Arnoux–Rauzy sequence $\tau_1^n(\omega)$ contains the k words $1^{n+1}, 1^n2, 1^n3, \dots, 1^nk$. Next suppose that av is not of the form 1^n . If b and c are each different from 1, then we can write $avb = 1vb = \tau_1(v'b)$ and $avc = 1vc = \tau_1(v'c)$ for some v' with $v'b, v'c \in \mathcal{AR}$ and $|v'| < |u|$. By the inductive hypothesis there exists an Arnoux–Rauzy sequence ω which contains both $v'b$ and $v'c$ as factors. It follows that the Arnoux–Rauzy sequence $\tau_1(\omega)$ contains both avb and avc as factors. Hence av is a right special factor of $\tau_1(\omega)$.

Next assume that one of b or c (say b) is equal to 1. Then $avb = 1v1 = 1v'11$ and $avc = 1vc = 1v'1c$ for some v' in \mathcal{AR} . Thus we can write $1v'1 = \tau_1(v''1)$ and $1v'1c = \tau_1(v''c)$ for some v'' with $v''1, v''c \in \mathcal{AR}$ and $|v''| < |u|$. By the inductive hypothesis there exists an Arnoux–Rauzy sequence

ω which contains both $v''1$ and $v''c$ as factors. Thus $\tau_1(\omega)$ contains both $1v'1 = \tau_1(v''1)$ and $1v'1c = \tau_1(v''c)$. As τ_1 of each letter begins with 1, it follows that $1v'11$ is also a factor of $\tau_1(\omega)$. Hence $1v'1 = av$ is a right special factor of $\tau_1(\omega)$.

CASE 2: $a \neq 1$. In this case it is easy to see that $1avb$ and $1avc$ are both in \mathcal{AR} . Applying the arguments of Case 1 we deduce that $1av$ is a right special factor of some Arnoux–Rauzy sequence ω . Hence so is av . ■

As an immediate consequence of Lemma 1 we have

COROLLARY 1. *Let $r(n)$ denote the number of right special factors of \mathcal{AR} of length n . Then*

$$\text{Card}(\mathcal{AR}_n) = \text{Card}(\mathcal{AR}_{n-1}) + (k-1)r(n-1).$$

Proof. In fact each right special factor of length $n-1$ is a prefix of k factors of length n . ■

LEMMA 2. *Suppose $u \in \mathcal{AR}$ is a bispecial factor of \mathcal{AR} , that is, there exist letters $a \neq b$ and $c \neq d$ such that au, bu, uc, ud are in \mathcal{AR} . Then there exists an Arnoux–Rauzy sequence $\omega \in \mathcal{AR}_\infty$ such that u is a bispecial factor of ω .*

Proof. The proof of Lemma 2 is similar to the proof of Lemma 1: using the τ_i the result follows by induction on the length of the words. ■

LEMMA 3. *If u is a bispecial factor of an Arnoux–Rauzy sequence, then for each $a \in A_k$ there exists an Arnoux–Rauzy sequence ω such that u is a bispecial factor of ω and au is a right special factor of ω .*

Proof. Let $\nu = \tau_{n_1} \circ \tau_{n_2} \circ \dots$ be a characteristic Arnoux–Rauzy sequence containing u as its r th bispecial factor, where we order the bispecial factors of ν according to increasing length. Fix $a \in A_k$ and let ω be any characteristic Arnoux–Rauzy sequence whose S -adic expansion begins with $\tau_{n_1} \circ \tau_{n_2} \circ \dots \tau_{n_r} \circ \tau_a$. Then in [27] it is proved that ω has the same first r bispecial factors of ν (the r th bispecial factor of a characteristic Arnoux–Rauzy sequence is completely determined by the first r terms of its S -adic expansion), and that au is a right special factor of ω (the $r+1$ st term of the S -adic expansion of ω determines which of the k factors $1u, 2u, \dots, ku$ is right special in ω). ■

COROLLARY 2. *Let $b(n)$ denote the number of bispecial factors of \mathcal{AR} of length n . Then*

$$r(n) = r(n-1) + (k-1)b(n-1).$$

Proof. In fact each bispecial factor of length $n-1$ is a suffix of k right special factors of length n . ■

Combining Corollaries 1 and 2 we have:

COROLLARY 3. Fix k and let \mathcal{AR}_n denote the set of all Arnoux–Rauzy words of length n on the alphabet $A_k = \{1, \dots, k\}$. Let $b(n)$ denote the number of bispecial words in \mathcal{AR} of length n . Then

$$\text{Card}(\mathcal{AR}_n) = k + (n - 1)k(k - 1) + (k - 1)^2 \sum_{i=1}^{n-2} (n - i - 1)b(i).$$

Proof. By Corollaries 1 and 2 we have

$$\begin{aligned} \text{Card}(\mathcal{AR}_n) &= k + (k - 1) \sum_{i=1}^{n-1} r(i) \\ &= k + (k - 1) \sum_{i=1}^{n-1} \left(k + (k - 1) \sum_{j=1}^{i-1} b(j) \right) \\ &= k + (n - 1)k(k - 1) + (k - 1)^2 \sum_{i=1}^{n-1} \sum_{j=1}^{i-1} b(j) \\ &= k + (n - 1)k(k - 1) + (k - 1)^2 \sum_{i=1}^{n-2} (n - i - 1)b(i). \blacksquare \end{aligned}$$

As a special case of Corollary 3 we recover the formula for the number of Sturmian words of length n .

COROLLARY 4. The number St_n of Sturmian words of length n is

$$1 + \sum_{i=1}^n (n - i + 1)\varphi(i)$$

where $\varphi(i)$ is the Euler phi function.

Proof. Applying Corollary 3 to the case $k = 2$ and the fact that $b(i) = \varphi(i + 2)$ (see [24] or Corollary 5 ahead) gives

$$\begin{aligned} \text{St}_n &= 2n + \sum_{i=1}^{n-2} (n - i - 1)b(i) = 2n + \sum_{j=3}^n (n - j + 1)\varphi(j) \\ &= 1 + \sum_{j=1}^n (n - j + 1)\varphi(j) \end{aligned}$$

where the last step follows from the equality $\varphi(1) = \varphi(2) = 1$. \blacksquare

3. A generalization of the Euler phi function. In the Sturmian case we have $b(n) = \varphi(n + 2)$. We now give a general arithmetic interpretation for the quantity $b(n)$ in terms of a multidimensional generalization of the Euclidean algorithm.

Fix k and let

$$E = \{(x_1, \dots, x_k) : \text{each } x_i \text{ is a nonnegative integer}\}.$$

For $z = (x_1, \dots, x_k) \in E$ set $|z| = \sum_{i=1}^k x_i$. Define a function $f : E \rightarrow E$ as follows: For $z = (x_1, \dots, x_k) \in E$ fix the least $1 \leq j \leq k$ such that $x_j \leq x_i$ for all $1 \leq i \leq k$ and set

$$f(z) = (x_1 - x_j, \dots, x_{j-1} - x_j, x_j, x_{j+1} - x_j, \dots, x_k - x_j).$$

Clearly for each $z \in E$ there exists a (unique) vector $\tilde{f}(z) \in E$ such that $f^n(z) = \tilde{f}(z)$ for all n sufficiently large. For $z \in E$ define the *generalized greatest common divisor* of z , denoted $\text{ggcd}(z)$, by

$$\text{ggcd}(z) = |\tilde{f}(z)|.$$

For instance, $f(4, 2, 5) = (2, 2, 3)$ and $f(2, 2, 3) = (2, 0, 1)$ so that $\tilde{f}(4, 2, 5) = (2, 0, 1)$ and $\text{ggcd}(4, 2, 5) = 3$. For $k = 2$ it follows immediately from the definition that $\text{ggcd}(a, b) = \text{gcd}(a, b)$.

To the best of our knowledge, this algorithm was first defined in [5] (in the special case $k = 3$) in connection with a generalization of the Fine–Wilf theorem to three periods.

Set $P = \{z = (x_1, \dots, x_k) \in E : \text{ggcd}(z) = 1\}$ and $P(n) = \{z \in P : |z| = n\}$. Then we have

THEOREM 1. *Fix k and let \mathcal{AR} denote the set of all Arnoux–Rauzy words on the alphabet $A_k = \{1, \dots, k\}$. Let $b(n)$ denote the number of bispecial words $u \in \mathcal{AR}$ of length n . Then $b(n) = \text{Card } P((k - 1)n + k)$.*

Proof. Let $B(n)$ denote the set of bispecial words in \mathcal{AR} of length n , so that $b(n) = \text{Card } B(n)$. For each $n \geq 1$ we construct a bijection $\psi_n : B(n) \rightarrow P((k - 1)n + k)$ as follows: Let $u \in B(n)$; according to Lemma 2, the word u is a bispecial factor of some Arnoux–Rauzy sequence $\omega \in \mathcal{AR}_\infty$. For each $1 \leq i \leq k$ let v (possibly the empty word) denote the longest proper prefix of u so that iv is a right special factor of ω . If such a v exists, set $x_i = |u| - |v|$. If no such v exists, set $x_i = |u| + 1$. It follows from the so-called “hat algorithm” given in Section III of [27] that for each i the quantity x_i is independent of the choice of ω . Set $\psi_n(u) = (x_1, \dots, x_k)$.

We now show that $\psi_n : B(n) \rightarrow P((k - 1)n + k)$ and is a bijection for each n . Taking $n = 1$ we have $B(1) = A_k = \{1, \dots, k\}$. Fixing $i \in B(1)$ we see by definition of ψ_1 that $\psi_1(i)$ is the vector whose i th coordinate is 1 and all other coordinates are 2, so that $|\psi_1(i)| = (k - 1)2 + 1 = (k - 1)1 + k$ as required. Moreover $f(\psi_1(i)) = (1, 1, \dots, 1)$ and $f^2(\psi_1(i)) = (1, 0, 0, \dots, 0)$ so that $\text{ggcd}(\psi_1(i)) = 1$. Clearly ψ_1 is injective. To see that ψ_1 is also surjective, let $z = (x_1, \dots, x_k) \in P((k - 1)2 + 1)$. Hence $|z| = (k - 1)2 + 1$ and $\text{ggcd}(z) = 1$. These conditions clearly imply that each $x_i > 0$. If all $x_i \geq 2$ we would have $|z| \geq 2k$, a contradiction. Hence some $x_i = 1$. We claim that all other coordinates of z are 2. In fact, if $x_s \neq x_t$ for some choice of $s \neq i$ and $t \neq i$, then $\tilde{f}(z)$ would have two nonzero coordinates, contradicting the fact that $\text{ggcd}(z) = 1$. Hence z is a vector whose i th coordinate is 1 and all

other coordinates are equal to one another. As $|z| = (k - 1)2 + 1$, the other coordinates of z must all be 2, whence it follows that $z = \psi_1(i)$ as required.

Now let $n > 1$ and suppose that $\psi_m : B(m) \rightarrow P((k - 1)m + k)$ is a bijection for all $m < n$. Let u be a bispecial word of length n . We begin by showing that $\psi_n(u) = (x_1, \dots, x_k)$ as defined above is in $P((k - 1)n + k)$. Fix an Arnoux–Rauzy sequence ω in which u is a bispecial factor. Let v denote the longest proper prefix of u which is also bispecial in ω and fix $i \in A_k$ so that iv is a right special factor of ω . Hence $x_i < x_j$ for all $1 \leq j \leq k$. Then by definition of ψ_n we have $x_i = ||u| - |v||$. Set $\psi_{|v|}(v) = (y_1, \dots, y_k)$. Thus $y_i = x_i$, in fact $u = vv'$ where v' is a suffix of v of length x_i (see the hat algorithm in Section III of [27]). Moreover, for $j \neq i$ we have $y_j = x_j - x_i$. Hence $\psi_{|v|}(v) = f(\psi_n(u)) = (x_1 - x_i, \dots, x_{i-1} - x_i, x_i, x_{i+1} - x_i, \dots, x_k - x_i)$. By the inductive hypothesis we have $\text{ggcd}(\psi_{|v|}(v)) = 1$ (and hence $\text{ggcd}(\psi_n(u)) = 1$) and $|\psi_{|v|}(v)| = (k - 1)|v| + k = (k - 1)(|u| - x_i) + k$, whence $|\psi_n(u)| = \sum_{j=1}^k x_j = (k - 1)|u| + k$ as required. Hence $\psi_n(B(n)) \subset P((k - 1)n + k)$.

If for some $u' \in B(n)$ with $u' \neq u$ we had $\psi_n(u) = \psi_n(u')$ then $\psi_{|v|}$ (where v is as above) would fail to be injective on $B(|v|)$. Hence ψ_n is one-to-one. To see that ψ_n is a surjection, let $z = (x_1, \dots, x_k) \in P((k - 1)n + k)$. Thus $|z| = (k - 1)n + k$ and $\text{ggcd}(z) = 1$. As in the case $n = 1$ these conditions imply that each $x_j > 0$. Fix i such that $x_i \leq x_j$ for all $1 \leq j \leq k$. We claim that $x_i < x_j$ for all $j \neq i$. In fact, if for some $j \neq i$ we had $x_j = x_i$, then $f(z)$ would have a coordinate equal to zero. Since $\text{ggcd}(z) = 1$ this would imply that $z = (1, 1, \dots, 1)$, contradicting $|z| = (k - 1)n + k$. Consider $f(z) = (x_1 - x_i, \dots, x_{i-1} - x_i, x_i, x_{i+1} - x_i, \dots, x_k - x_i)$. Then $\text{ggcd}(f(z)) = 1$ (since $\text{ggcd}(z) = 1$) and hence $f(z) \in P((k - 1)(n - x_i) + k)$. By the inductive hypothesis, since $\psi_{n-x_i} : B(n - x_i) \rightarrow P((k - 1)(n - x_i) + k)$ is onto, we have $f(z) = \psi_{n-x_i}(v)$ for some bispecial word $v \in \mathcal{AR}$ of length $n - x_i$. Let ω be any Arnoux–Rauzy sequence containing iv as a right special factor, and let u be the shortest bispecial factor of ω beginning with vi . Then it follows from the hat algorithm that $u = vv'$ where v' is a suffix of v of length x_i (see Section III of [27]). Hence $|u| = n$ and $\psi_n(u) = (x_1, \dots, x_k) = z$ as required. ■

As a special case of Theorem 1 we have:

COROLLARY 5. *The number of bispecial Sturmian words of length n is $\varphi(n + 2)$, where φ denotes the Euler phi function.*

Proof. Applying Theorem 1 to the case $k = 2$ gives $b(n) = \text{Card } P(n+2)$. But

$$\begin{aligned} P(n + 2) &= \{(a, b) \in E \mid a + b = n + 2 \text{ and } \text{ggcd}(a, b) = 1\} \\ &= \{(a, n + 2 - a) \in E \mid \text{gcd}(a, n + 2 - a) = 1\} \\ &= \{(a, n + 2 - a) \in E \mid \text{gcd}(a, n + 2) = 1\}, \end{aligned}$$

hence $\text{Card } P(n+2) = \text{Card}\{a \mid 1 \leq a < n+2 \text{ and } \gcd(a, n+2) = 1\} = \varphi(n+2)$. ■

From Theorem 1 and Corollary 3 we deduce:

COROLLARY 6. *Fix k and let \mathcal{AR}_n denote the set of all Arnoux–Rauzy words of length n on the alphabet $A_k = \{1, \dots, k\}$. Then*

$$\text{Card}(\mathcal{AR}_n) = k + (n-1)k(k-1) + (k-1)^2 \sum_{i=1}^{n-2} (n-i-1) \text{Card } P((k-1)i+k).$$

REMARK 1. In [24], following a suggestion of G. Rauzy, Mignosi establishes a connection between the number of Sturmian words, Farey numbers, and the Riemann hypothesis. Another such connection between the Riemann hypothesis, the Euler phi function and the formula in Corollary 4 was given by Bender, Patashnik, and Rumsey [3] using a result of Codèca [10]. It would be interesting to find similar connections involving the multidimensional generalization of the Euler phi function described in this paper for $k > 2$, two-dimensional Farey numbers in the sense of [2, 29, 30], and deep results and conjectures in analytical number theory. Also, Rychlik points out a possible connection between our multidimensional generalization of the Euler phi function and Gröbner bases [28].

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