

## The square-free kernel of $x^{2^n} - a^{2^n}$

by

PAULO RIBENBOIM (Kingston, Ont.)

*Dedicated to my long-time friend and collaborator Wayne McDaniel,  
at the occasion of his retirement*

### 1. Introduction

*A. Statement of the results.* We investigate the number  $\nu(x^{2^n} - a^{2^n})$  of odd prime factors of the square-free kernel of numbers  $x^{2^n} - a^{2^n}$ , where  $x > a \geq 1$  and  $n \geq 2$ . The main theorem states that (under a certain assumption) for each  $a \geq 1$  the set  $T_a = \{(x, n) \mid n \geq 2, x > a \text{ and the square-free kernel of } x^{2^n} - a^{2^n} \text{ has } n - 1 \text{ odd prime factors}\}$  is finite and effectively computable.

In the final section, we show with several examples how to determine explicitly the sets  $T_a$ , namely  $T_1, T_2, T_3, T_4, T_5, T_6, T_{10}$ . As an illustration of the results obtained,

$$\nu(3^{2^n} - 1) \geq n \quad \text{for all } n \geq 4,$$

$$\nu(7^{2^n} - 1) \geq n \quad \text{for all } n \geq 4,$$

$$\nu(99^{2^n} - 1) \geq n \quad \text{for all } n \geq 3,$$

and if  $x \neq 3, 7, 99$  then

$$\nu(x^{2^n} - 1) \geq n \quad \text{for all } n \geq 2.$$

The proofs rely on properties of binary linearly recurring sequences and more specifically on a special case of the main theorem in Ribenboim [7].

Now we gather the concepts and facts used in this paper.

*B. Binary linearly recurring sequences.* Let  $P > 0$ ,  $Q \neq 0$  be integers such that  $\gcd(P, Q) = 1$  and  $D = P^2 - 4Q \neq 0$ .

Let  $U_0 = 0$ ,  $U_1 = 1$ ,  $V_0 = 2$ ,  $V_1 = P$  and for  $n \geq 2$ :

$$U_n = PU_{n-1} - QU_{n-2}, \quad V_n = PV_{n-1} - QV_{n-2}.$$

We also define  $U_{-n} = -U_n/Q^n$ ,  $V_{-n} = V_n/Q^n$  (for  $n > 0$ ); then the above formulas still hold.

$\mathcal{U} = (U_n)_n$ ,  $\mathcal{V} = (V_n)_n$  are called *binary linearly recurring sequences of first kind*, respectively of *second kind*, with parameters  $(P, Q)$ . We also use the notation  $U_n(P, Q)$ ,  $V_n(P, Q)$ .

For an expository account of the theory of binary linearly recurring sequences, see Chapter 1 of Ribenboim [6]. Here we limit ourselves to mention explicitly the facts which are used in what follows.

If  $P = 2$ ,  $Q = -1$  the numbers  $U_n$ ,  $V_n$  are the Pell numbers of first kind, respectively of second kind. These numbers are (for  $n \geq 0$ ):

$$\begin{aligned} U_n &: 0 \ 1 \ 2 \ 5 \ 12 \ 29 \ 70 \ 169 \dots, \\ V_n &: 2 \ 2 \ 6 \ 14 \ 34 \ 82 \ 198 \ 478 \dots \end{aligned}$$

Then  $U_n$  is even if and only if  $n$  is even,  $2 \mid V_n$  but  $4 \nmid V_n$  for all  $n$ .

The symbol  $\square$  denotes any non-zero integer which is a square.

Concerning square and double square Pell numbers, we quote the following important result of Ljunggren [1] (see also Ribenboim [5]); in particular, the proof of (a) is difficult.

(1.1) *For Pell numbers:*

- (a)  $U_n = \square$  if and only if  $n = 1, 7$ ;
- (b)  $U_n = 2\square$  if and only if  $n = 2$ ;
- (c)  $V_n \neq \square$  for all  $n$ ;
- (d)  $V_n = 2\square$  if and only if  $n = 0, 1$ .

*C. Pell equations.* Let  $F > 1$  be a square-free integer, and  $\varepsilon = c + d\sqrt{F}$  be the fundamental unit of the ring  $\mathbb{Z}[\sqrt{F}]$ , so  $1 < \varepsilon$ . Let  $Q = N(\varepsilon) = c^2 - d^2F = \pm 1$  be the norm of  $\varepsilon$ . We consider the equations

$$x^2 - Fy^2 = \pm 1.$$

(1.2) *Solutions of  $x^2 - Fy^2 = 1$ .* The solutions  $(x, y)$  with  $x + y\sqrt{F} > 0$  are given by  $(x_n, y_n)$ , where

$$x_n + y_n\sqrt{F} = \varepsilon^n \begin{cases} \text{for all } n \text{ if } Q = 1, \\ \text{for all even } n \text{ if } Q = -1. \end{cases}$$

(1.3) *Solutions of  $x^2 - Fy^2 = -1$ .* The solutions  $(x, y)$  with  $x + y\sqrt{F} > 0$  are given by  $(x_n, y_n)$ , where  $x_n + y_n\sqrt{F} = \varepsilon^n$  and  $Q = -1$ ,  $n$  odd. If  $Q = 1$  there are no solutions.

It is possible to express  $(x_n, y_n)$  by means of terms of a binary linearly recurring sequence.

Let  $\varepsilon = c + d\sqrt{F}$  as before, let  $P = 2c$ ,  $Q = N(\varepsilon) = \pm 1$  and consider the sequences  $\mathcal{U}$ ,  $\mathcal{V}$  with parameters  $(P, Q)$ . We note that  $V_n$  is even for all  $n$ . Then:

$$(1.4) \quad x_n = V_n/2, \quad y_n = dU_n \text{ for all } n.$$

We shall require the following result (part (a) was first proved by Ljunggren [2] and a simpler proof was given by Samuel [8]; in the same paper, Samuel proved also (b)):

(1.5) *Let  $x > 1$  and let  $p$  be any prime.*

(a) *If  $x^4 - 1 = p\Box$  then  $(x, p) = (3, 5)$  or  $(99, 29)$ .*

(b) *If  $x^4 - 1 = 2p\Box$  then  $(x, p) = (7, 3)$ .*

In Ribenboim [7] we considered families of systems of two Pell equations. Let  $F > 1$  and  $G > 0$  be square-free integers, let  $f, g$  be non-zero integers. We denote by  $(F, f | G, g)$  the family of systems—one for each prime  $p$ —of Pell equations

$$(F, f | G, g) \quad \begin{cases} x^2 - f = F\Box, \\ x^2 - g = Gp\Box. \end{cases}$$

We proved a theorem for certain families of the above kind. Here we shall only need the following special case:

(1.6) *For each  $b \geq 1$  the set of solutions  $(x, b)$  of each family below is finite and effectively computable:  $(2, b^2 | 1, -b^2)$ ,  $(2, -b^2 | 1, b^2)$ ,  $(2, b^2 | 2, -b^2)$ ,  $(2, -b^2 | 2, b^2)$ .*

**2. The main theorem.** For every  $m \geq 1$  let  $\nu(m)$  denote the number of odd prime factors of the square-free kernel of  $m$ . So  $\nu(m) = 0$  if and only if  $m = \Box$  or  $m = 2\Box$ . And  $\nu(m) = 1$  if and only if  $m = p\Box$  or  $m = 2p\Box$ , where  $p$  is any odd prime. It is immediate that if  $\gcd(m, n) = 1$  or  $2$ , then  $\nu(mn) = \nu(m) + \nu(n)$ .

For all  $a \geq 1$  and  $n \geq 1$  we define the set

$$S_{a,n} = \{x \mid x > a \text{ and } \nu(x^{2^n} - a^{2^n}) = n - 1\}.$$

In particular,  $S_{a,1} = \{x \mid x > a \text{ and } x^2 - a^2 = \Box \text{ or } 2\Box\}$ .

We introduce the following notation. Let  $x > a \geq 1$  and  $n \geq 1$ ; we define the integers  $u_n, v_n$  (which depend on  $x, a$ ) as follows:

$$u_n = x^{2^n} - a^{2^n}, \quad v_n = x^{2^n} + a^{2^n}.$$

It is easy to verify the following properties. If  $\gcd(x, a) = 1$  then  $\gcd(u_n, v_m) = 1$  or  $2$  (for all  $n, m$ ),  $\gcd(v_n, v_m) = 1$  or  $2$  (for all  $n \neq m$ ) and  $u_n = u_{n-1}v_{n-1}$  for all  $n \geq 2$ . The integers  $u_n, v_n$  may be also defined with the help of a binary linearly recurring sequence. Let  $P = x + a, Q = xa$ ; then  $\gcd(P, Q) = 1$  and

$$u_n = (x - a) \cdot U_{2^n}(P, Q), \quad v_n = V_{2^n}(P, Q).$$

We shall need the following facts.

(2.1) LEMMA. *Let  $x > a \geq 1$  and  $n \geq 2$ .*

- 1)  $\nu(x^{2^n} + a^{2^n}) \neq 0$ .
- 2)  $\nu(x^{2^n} - a^{2^n}) > n - 2$ .

*Proof.* 1) We show that  $x^{2^n} + a^{2^n} \neq \square, 2\square$ . As  $n \geq 2$ , we have  $x^{2^n} + a^{2^n} = (x^{2^{n-2}})^4 + (a^{2^{n-2}})^4 \neq \square$  by the classical result of Fermat (see for example Ribenboim [4]). Similarly, if  $x^{2^n} + a^{2^n} = (x^{2^{n-2}})^4 + (a^{2^{n-2}})^4 = 2\square$  then again  $x^{2^{n-2}} = a^{2^{n-2}}$ , so  $x = a$  (see Ribenboim [4]) and this has been excluded.

2) We may assume without loss of generality that  $\gcd(x, a) = 1$ . Indeed, if  $\gcd(x, a) = e$ , let  $x = ze$ ,  $a = be$ , hence  $x^{2^n} - a^{2^n} = e^{2^n}(z^{2^n} - b^{2^n})$  and  $\nu(x^{2^n} - a^{2^n}) = \nu(z^{2^n} - b^{2^n})$ .

We prove the statement by induction on  $n$ . Let  $n = 2$ . By the classical theorem of Fermat (see [4]),  $x^4 - a^4 \neq \square$ . Next we show that  $x^4 - a^4 \neq 2\square$ . We quote the following theorem of Euler: If  $u^4 - v^4 = 2w^2$  then  $u = v, w = 0$ . For a proof, see Ribenboim [3], Proposition A14.5. Therefore if  $x > a \geq 1$  then  $x^4 - a^4 \neq 2\square$ .

Now, let  $n \geq 3$  and assume that the statement is true for  $n - 1$ . We have  $x^{2^n} - a^{2^n} = u_n = u_{n-1}v_{n-1}$  with  $\gcd(u_{n-1}, v_{n-1}) = 1$  or  $2$ , since  $\gcd(x, a) = 1$ . So  $\nu(u_n) = \nu(u_{n-1}v_{n-1}) = \nu(u_{n-1}) + \nu(v_{n-1}) > n - 3 + 1 = n - 2$ . ■

We introduce some sets. For all  $a \geq 1, n \geq 1$  and for all  $e$  dividing  $a$ , let

$$S_{a,n}(e) = \{x \in S_{a,n} \mid \gcd(x, a) = e\}.$$

If  $e, e'$  divide  $a$  and  $e \neq e'$  then  $S_{a,n}(e) \cap S_{a,n}(e') = \emptyset$  and  $S_{a,n} = \bigcup_{e|a} S_{a,n}(e)$ . If  $x \in S_{a,n}(e)$ , let  $x = ze$  and  $a = be$ . Then  $z > b, \gcd(z, b) = 1$  and  $\nu(e^{2^n}(z^{2^n} - b^{2^n})) = n - 1$ , so  $\nu(z^{2^n} - b^{2^n}) = n - 1$ , so  $z \in S_{b,n}(1)$ . The mapping  $x \mapsto z$  is a bijection between  $S_{a,n}(e)$  and  $S_{b,n}(1)$ ; moreover the mapping is effectively computable.

Let  $a \geq 1$ . The set  $S_{a,1}$  is infinite. Indeed, let  $\varepsilon = 1 + \sqrt{2}$  be the fundamental unit of  $\mathbb{Z}[\sqrt{2}]$ , and for every even  $m \geq 1$ , let  $z_m + u_m\sqrt{2} = (1 + \sqrt{2})^m$ . Hence  $z_m^2 - 2u_m^2 = 1$ , so if  $x_m = az_m$  then  $x_m^2 - a^2 = 2\square$ . So  $x_m \in S_{a,1}$ , showing that this set is infinite.

For  $n \geq 2$  we have:

- (2.2) THEOREM. 1)  $S_{a,2} \supseteq S_{a,3} \supseteq \dots$   
 2)  $S_{a,2}$  is a finite effectively computable set.

*Proof.* 1) Let  $n \geq 3$ ; we show that  $S_{a,n} \subseteq S_{a,n-1}$ . It suffices to show that, for every  $e \mid a, S_{a,n}(e) \subseteq S_{a,n-1}(e)$ , or equivalently, for every  $b$  dividing  $a, S_{b,n}(1) \subseteq S_{b,n-1}(1)$ .

Let  $z \in S_{b,n}(1)$ , so  $z > b, \gcd(z, b) = 1$  and  $\nu(z^{2^n} - b^{2^n}) = n - 1$ . Let  $d = \gcd(z^{2^{n-1}} - b^{2^{n-1}}, z^{2^{n-1}} + b^{2^{n-1}})$ , so  $d \mid 2b^{2^{n-1}}$ ; but  $\gcd(z, b) = 1$ , hence

$d = 1$  or  $2$ . We may write

$$\begin{cases} z^{2^{n-1}} - b^{2^{n-1}} = k, \\ z^{2^{n-1}} + b^{2^{n-1}} = h, \end{cases}$$

with  $\gcd(k, h) = 1$  or  $2$ ,  $n - 1 = \nu(kh) = \nu(k) + \nu(h)$ . By (2.1),  $\nu(h) \geq 1$ , so  $\nu(k) \leq n - 2$ . By (2.1),  $\nu(k) > n - 3$ , hence  $\nu(k) = n - 2$ , showing that  $z \in S_{b,n-1}(1)$ .

2) To show that  $S_{a,2}$  is finite and effectively computable, it suffices to show that for every  $e \mid a$ ,  $S_{a,2}(e)$  is finite and effectively computable, or equivalently, for every  $b \mid a$ , the set  $S_{b,2}(1)$  is finite and effectively computable.

Now  $z \in S_{b,2}(1)$  if and only if  $z > b$ ,  $\gcd(z, b) = 1$  and  $\nu(z^4 - b^4) = 1$  and this means that  $z^4 - b^4 = p\Box$  or  $2p\Box$ , for some odd prime  $p$ . We have  $\gcd(z^2 - b^2, z^2 + b^2) = 1$  or  $2$ , because  $\gcd(z, b) = 1$ . Then the following cases may happen:

$$\begin{cases} z^2 - b^2 = \Box & \Big| & p\Box & \Big| & 2\Box & \Big| & 2p\Box \\ z^2 + b^2 = p\Box & \Big| & \Box & \Big| & 2p\Box & \Big| & 2\Box \end{cases} \quad \text{when } z^4 - b^4 = p\Box,$$

(1)    (2)    (3)    (4)

$$\begin{cases} z^2 - b^2 = \Box & \Big| & 2\Box & \Big| & p\Box & \Big| & 2p\Box \\ z^2 + b^2 = 2p\Box & \Big| & p\Box & \Big| & 2\Box & \Big| & \Box \end{cases} \quad \text{when } z^4 - b^4 = 2p\Box.$$

(5)    (6)    (7)    (8)

In cases (1), (2), (5) and (8),  $z$  belongs to a finite and effectively computable set. By (1.6), the families  $(2, \pm b^2 \mid 2, \mp b^2)$  and  $(2, \pm b^2 \mid 1, \mp b^2)$  have a finite effectively computable set of solutions  $(z, p)$ . So, in cases (3), (4), (6) and (7),  $z$  belongs to a finite and effectively computable set. This shows that  $S_{b,2}(1)$  is finite and effectively computable. ■

Consider the following statement about the pair of integers  $(b, z)$ :

$(H_{b,z})$  If  $z > b \geq 1$ ,  $\gcd(z, b) = 1$  and  $\nu(z^4 - b^4) = 1$ , there exists an effectively computable  $h \geq 2$  (depending on  $z, b$ ) such that  $\nu(z^{2^h} + b^{2^h}) > 1$ .

No proof is known for this statement but, of course it holds in every numerical example computed thus far.

(2.3) THEOREM. Assume that the statement  $(H_{b,z})$  holds for  $z > b \geq 1$  with  $\gcd(z, b) = 1$  and  $\nu(z^4 - b^4) = 1$ . Let  $h \geq 2$  be the smallest integer such that  $\nu(z^{2^h} + b^{2^h}) > 1$ . Then  $z \notin S_{b,j}(1)$  for all  $j \geq h + 1$ .

*Proof.* With the notation introduced, we have  $\nu(u_2) = \nu(z^4 - b^4) = 1$ , and

$$z^{2^j} - b^{2^j} = u_j = v_{j-1}v_{j-2} \dots v_{h+1}v_h v_{h-1} \dots v_2 u_2.$$

As already stated,  $\gcd(u_2, v_i) = 1$  or  $2$  (for all  $i$ ) and  $\gcd(v_i, v_l) = 1, 2$  for  $i \neq l$ . So

$$\nu(u_j) = \nu(v_{j-1}) + \dots + \nu(v_{h+1}) + \nu(v_h) + \dots + \nu(v_2) + \nu(u_2).$$

By (2.1) and the hypothesis,  $\nu(u_j) \geq (j - 1 - h) + 2 + (h - 2) + 1 = j$ , so  $z \notin S_{b,j}(1)$ . ■

If  $a \geq 1$  let

$$T_a = \{(x, n) \mid n \geq 2, x \in S_{a,n}\}.$$

For every  $e$  dividing  $a$ , let

$$T_a(e) = \{(x, n) \in T_a \mid \gcd(x, a) = e\}.$$

If  $e \mid a$ ,  $b = a/e$ ,  $z = x/e$  and  $(x, n) \in T_a(e)$  then  $(z, n) \in T_b(1)$ . The mapping  $(x, n) \mapsto (z, n)$  is a bijection between  $T_a(e)$  and  $T_b(1)$ .

(2.4) THEOREM. *Let  $a \geq 1$  and assume that  $(H_{b,z})$  holds for every  $b$  dividing  $a$  and  $z > b$ . Then  $T_a$  is a finite and effectively computable set.*

*Proof.* It suffices to show that for every  $e$  dividing  $a$ , the set  $T_a(e)$  is finite and effectively computable. By the above remark it suffices to show that for every  $b$  dividing  $a$ , the set  $T_b(1)$  is finite and effectively computable. By (2.2) the set  $S_{b,2}(1)$  is finite and effectively computable. By (2.3) and the hypothesis, for every  $z_0 \in S_{b,2}(1)$  there exists an effectively computable integer  $h \geq 2$  (depending on  $b$  and  $z_0$ ) such that if  $z_0 \in S_{b,i}(1)$  then  $i \leq h$ . So the set

$$T_b(1)|_{z_0} = \{(z, n) \in T_b(1) \mid z = z_0\}$$

is finite and effectively computable, hence

$$T_b(1) = \bigcup_{z_0 \in S_{b,2}(1)} T_b(1)|_{z_0}$$

is also finite and effectively computable. ■

**3. Explicit computations.** For specific values of  $a \geq 1$ , it is possible to determine explicitly the finite effectively computable set  $T_a$ . This determination requires the actual solution of certain families of systems of Pell equations. We recall that if  $a \geq 1$  then

$$T_a = \{(x, n) \mid n \geq 2, x > a, \nu(x^{2^n} - a^{2^n}) = n - 1\}.$$

The following easy remark will be useful: If  $(x, n) \in T_a$  then  $(mx, n) \in T_{ma}$ .

(3.1) *Let  $a = 1$ . Then  $T_1 = \{(3, 2), (3, 3), (7, 2), (7, 3), (99, 2)\}$ .*

*Proof.* We determine explicitly  $S_{1,2} = \{x \mid x > 1, \nu(x^4 - 1) = 1\}$ . If  $x^4 - 1 = p\Box$  for some odd prime  $p$ , then by (1.5),  $(x, p) = (3, 5)$  or  $(99, 29)$ .

If  $x^4 - 1 = 2p\Box$  for some odd prime  $p$ , then by (1.5),  $(x, p) = (7, 3)$ . This shows that  $S_{1,2} = \{3, 7, 99\}$ .

Now

$$\begin{aligned} 3^4 + 1 &= 82 = 2 \times 41, & \text{so } \nu(3^4 + 1) &= 1, \\ 3^8 + 1 &= 2 \times 17 \times 193, & \text{so } \nu(3^8 + 1) &= 2, \\ 7^4 + 1 &= 2 \times 1201, & \text{so } \nu(7^4 + 1) &= 1, \\ 7^8 + 1 &= 2 \times 17 \times 169553, & \text{so } \nu(7^8 + 1) &= 2, \\ 99^4 + 1 &= 2 \times 2617 \times 18353, & \text{so } \nu(99^4 + 1) &= 2. \end{aligned}$$

Thus  $(3, 2), (3, 3) \in T_1$ ,  $(3, j) \notin T_1$  for all  $j \geq 4$ ;  $(7, 2), (7, 3) \in T_1$ ,  $(7, j) \notin T_1$  for all  $j \geq 4$ ;  $(99, 2) \in T_1$ ,  $(99, j) \notin T_1$  for all  $j \geq 3$ . ■

$$(3.2) \quad T_2 = \{(6, 2), (6, 3), (14, 2), (14, 3), (198, 2)\}.$$

*Proof.* Let  $x > 2$  be such that  $x^4 - 2^4 = p\Box$  or  $2p\Box$ , for some odd prime  $p$ .

*First case:*  $x$  is even. Let  $x = 2z$ . Then  $2^4(z^4 - 1) = p\Box$  or  $2p\Box$ , hence  $z^4 - 1 = p\Box$  or  $2p\Box$ . As stated in (3.1),  $z = 3, 99$  or  $7$ , hence  $x = 6, 198$  or  $14$ . We have  $6^4 + 2^4 = 2^4(3^4 + 1)$  so

$$\nu(6^4 + 2^4) = \nu(3^4 + 1) = 1;$$

similarly

$$\nu(6^8 + 2^8) = \nu(3^8 + 1) = 2.$$

In the same manner

$$\begin{aligned} \nu(14^4 + 2^4) &= \nu(7^4 + 1) = 1, & \nu(14^8 + 2^8) &= \nu(7^8 + 1) = 2, \\ \nu(198^4 + 2^4) &= \nu(99^4 + 1) = 2. \end{aligned}$$

Altogether, only  $(6, 2), (6, 3), (14, 2), (14, 3), (198, 2) \in T_2$ .

*Second case:*  $x$  is odd. So  $\gcd(x^2 - 4, x^2 + 4) = 1$ . Since  $x^4 - 2^4$  is odd we have  $x^4 - 2^4 \neq 2p\Box$  and there are only the following cases:

$$\begin{cases} x^2 - 4 = \Box & | & p\Box \\ x^2 + 4 = p\Box & | & \Box \end{cases} \begin{matrix} (1) \\ (2) \end{matrix}$$

Subcase (1): there exists  $t \neq 0$  such that  $x^2 - t^2 = 4$ , which is clearly impossible.

Subcase (2): there exists  $t$  such that  $t^2 - x^2 = 4$ , which is again impossible. ■

$$(3.3) \quad T_3 = \{(9, 2), (9, 3), (21, 2), (21, 3), (297, 2), (4, 2), (4, 3), (5, 2), (5, 3), (5, 4)\}.$$

*Proof.* Let  $x > 3$  be such that  $x^4 - 3^4 = p\Box$  or  $2p\Box$ , for some odd prime  $p$ .

*First case:*  $3 \mid x$ . Let  $x = 3z$ . Then  $z^4 - 1 = p\Box$  or  $2p\Box$ . As already seen,  $z = 3, 99, 7$  so  $x = 9, 297, 21$ . We have, as computed in (3.1),

$$\nu(9^4 + 3^4) = \nu(3^4(3^4 + 1)) = 1, \quad \nu(9^8 + 3^8) = \nu(3^8(3^8 + 1)) = 2$$

and similarly

$$\nu(21^4 + 3^4) = 1, \quad \nu(21^8 + 3^8) = 2, \quad \nu(297^4 + 3^4) = 2.$$

Thus, only  $(9, 2), (9, 3), (21, 2), (21, 3)$ , and  $(297, 2)$  are in  $T_3$ .

*Second case:*  $\gcd(x, 3) = 1$ . Then  $d = \gcd(x^2 - 3^2, x^2 + 3^2) = 1$  or  $2$ , because  $d \mid 18$  but  $3 \nmid d$ .

*Case A:*  $d = 1$ . If  $x^4 - 3^4 = p\Box$  then

$$\begin{cases} x^2 - 3^2 = \Box & \Big| & p\Box \\ x^2 + 3^2 = p\Box & \Big| & \Box \end{cases} \begin{matrix} (1) \\ (2) \end{matrix}$$

(1) is not possible, while (2) gives  $(x, p) = (4, 7)$ .

We have  $4^4 + 3^4 = 337$ , prime,  $\nu(4^8 + 3^8) = \nu(17 \times 4241) = 2$ . Then only  $(4, 2)$  and  $(4, 3)$  are in  $T_3$ .

If  $x^4 - 3^4 = 2p\Box$  then  $x$  is odd. On the other hand, since  $d = 1$ , it follows that  $x$  is even, a contradiction.

*Case B:*  $d = 2$ . If  $x^4 - 3^4 = p\Box$  then

$$\begin{cases} x^2 - 3^2 = 2\Box & \Big| & 2p\Box \\ x^2 + 3^2 = 2p\Box & \Big| & 2\Box \end{cases} \begin{matrix} (1) \\ (2) \end{matrix}$$

Both cases are impossible; this is seen modulo 3:

$$1 \equiv x^2 \mp 3^2 = 2\Box \pmod{3}.$$

If  $x^4 - 3^4 = 2p\Box$  we have one of the following cases:

$$\begin{cases} x^2 - 3^2 = \Box & \Big| & 2\Box & \Big| & p\Box & \Big| & 2p\Box \\ x^2 + 3^2 = 2p\Box & \Big| & p\Box & \Big| & 2\Box & \Big| & \Box \end{cases} \begin{matrix} (1) \\ (2) \\ (3) \\ (4) \end{matrix}$$

In (1) we have  $(x, p) = (5, 17)$ . Since  $\nu(5^4 + 3^4) = \nu(2 \times 353) = 1$ ,  $\nu(5^8 + 3^8) = \nu(2 \times 198593) = 1$  and  $\nu(5^{16} + 3^{16}) = \nu(2 \times 97 \times 786757409) = 2$ , we have only  $(5, 2), (5, 3), (5, 4) \in T_3$ .

In (2),  $x$  is odd, so  $2 \equiv x^2 + 3^2 = p\Box \pmod{4}$ , which is impossible.

In (3), since  $3 \nmid x$  we have  $1 \equiv x^2 + 3^2 \equiv 2\Box \pmod{3}$  and this is impossible.

(4) is also impossible.

The reader may wish to show, with the same method:

$$(3.4) \quad T_5 = \{(15, 2), (15, 3), (35, 2), (35, 3), (495, 2), (13, 2), (13, 3)\}.$$



$$(3.5) \quad T_4 = \{(12, 2), (12, 3), (28, 2), (28, 3), (396, 2), (5, 2), (5, 3)\}.$$

$$(3.6) \quad T_6 = \{(18, 2), (18, 3), (42, 2), (42, 3), (594, 2), (8, 2), (8, 3), \\ (10, 2), (10, 3), (10, 4)\}.$$

$$(3.7) \quad T_{10} = \{(30, 2), (30, 3), (70, 2), (70, 3), (990, 2), (26, 2), (26, 3)\}.$$

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Department of Mathematics and Statistics  
 Queen's University  
 Kingston, Ontario  
 Canada K7L 3N6

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