

# Some PA<sup>-</sup> exercises

24 października 2007

## 1 2007/10/09

1. Show that

- $\text{PA}^- \vdash \forall x \forall y (x < y \Rightarrow x + 1 \leq y)$ ,
- $\text{PA}^- \vdash \forall x \exists y (y + 1 = x)$ ,
- for all  $n \in \omega$ ,  $\text{PA}^- \vdash \forall x (x \leq \underline{n} \iff \bigvee_{i \leq n} x = \underline{i})$ .

2. Define ordering on  $\mathbb{Z}[X]$ , the set of polynomials with coefficients from  $\mathbb{Z}$  such that the set of polynomials greater or equal the zero polynomial satisfies  $\text{PA}^-$ .

Show that  $\text{PA}^- \not\vdash \forall x \exists y (2y = x \vee 2y + 1 = x)$ .

3. If we would define this ordering on  $\mathbb{N}[X]$  instead of  $\mathbb{Z}[X]$ , would  $\mathbb{N}[X] \models \text{PA}^-$ ?

Let  $T = \text{Th}(\mathcal{N})$ .

**Definition 1** An element  $a \in M \models \text{PA}$  codes a set

$$X = \{i \in \omega : M \models \underline{p}_i | x[a/x]\},$$

where  $p_i$  is the  $i$ -th prime number. We denote this set by  $c_M(a)$ .

If, for  $X \subseteq \omega$ , there exists  $a \in M$  such that  $X = c_M(a)$  we say that  $X$  is coded in  $M$ .

4. Show that each set  $X \subseteq \omega$  is coded in some countable model of  $T$ . Conclude that there are continuum many nonisomorphic countable models of  $T$ .

5. Show that all subsets of  $\omega$  are coded in  $\prod_{i \in \omega} \mathcal{N}/\mathcal{U}$ . What is the cardinality of  $\prod_{i \in \omega} \mathcal{N}/\mathcal{U}$ ?
6. Describe all, up to isomorphism, models of  $\text{Th}((\omega, S, 0, 1))$ , where  $S$  is the successor function.  
You may need to use Ehrenfeuch–Fraïsse games or elimination of quantifiers for  $\text{Th}((\omega, S, 0, 1))$ .
7. What about models of  $\text{Th}((\omega, \leq, 0, 1))$ ?

## 2 2007 /10 /16 and 2007/10 /23

**Definition 2** Let  $M \models \text{PA}^-$ .  $\emptyset \neq I \subsetneq_e M$  is a cut if it is closed downward and closed on successor.

1. Show that for  $M \models \text{PA}^-$ ,  $M$  satisfies overspill for all cuts if and only if  $M \models \text{PA}$ .
2. (\*) Show that for each model  $M \models \text{PA}^-$  there is a model  $N$  such that  $M \prec N$  and  $N$  satisfies overspill for  $\mathbb{N}$ . Conclude that overspill only for  $\mathbb{N}$  is weaker than induction.
3. Show that for each nonstandard  $M \models \text{PA}$  there is continuum many cuts of  $M$ . (It suffices to restrict only to countable models. Then, use the characterization of the order type of countable  $M \models \text{PA}$  as  $\omega + (\omega^* + \omega)\eta$ , where  $\eta$  is the order type of rationals.)
4. Repeat the above but with cuts closed on addition and multiplication.
5. Show that  $I\Sigma_n, L\Sigma_n, I\Pi_n, L\Pi_n$  are equivalent for all  $n$ .

**Definition 3** A model  $M \models \text{PA}^-$  is  $\omega_1$ -like if for all  $a \in |M|$ ,

$$\text{card}(\{b \in |M| : M \models b \leq a\}) = \omega_0 \text{ and } \text{card}(|M|) = \omega_1.$$

6. Show that if  $M \models \text{PA}^-$  is  $\omega_1$ -like then  $M \models \text{Coll}_n$ , for all  $n \in \mathbb{N}$ .
7. (\*) Show that for each model  $M \models \text{PA}^-$  there is  $N$  such that  $\text{card}(N) = \text{card}(M)$  and  $M \subseteq_e N$ . (Use the characterization of a model  $M$  of  $\text{PA}^-$  as a positive part of a discretely ordered ring  $R_M$ . Then, consider a positive part of a ring of polynomials  $R_M[X]$ .)

8. (\*) Show that for each countable model  $M \models \text{PA}^-$  there is  $\omega_1$ -like  $N \models \text{PA}^-$  such that  $M \subseteq_e N$ .
9. (\*) Show that  $\text{PA}^- \cup \bigcup_{n \in \mathbb{N}} \text{Coll}_n$  does not prove  $\forall x \exists y (2y = x \vee 2y + 1 = x)$  (use the fact that  $\text{PA}^-$  does not prove this).
10. Show that  $I\Sigma_n \vdash B\Sigma_n$ .
11. Show that  $B\Sigma_{n+1} \equiv B\Pi_n$ .
12. Show that  $B\Sigma_{n+1} \vdash I\Sigma_n$ . Also show that  $\Sigma_n(B\Sigma_n)$  is closed on bounded quantification (all it was done during the lecture so there is no need to do this).
13. (\*) Show that induction for parameter free formulas is equivalent to PA (Note. It is not the case that induction for parameter free  $\Sigma_n$  formulas is equivalent to  $I\Sigma_n$ .)